

The Libgcrypt Reference Manual

Version 1.6.0-git5e14de0
29 June 2011

Werner Koch (wk@gnupg.org)
Moritz Schulte (mo@g10code.com)

This manual is for Libgcrypt (version 1.6.0-git5e14de0, 29 June 2011), which is GNU's library of cryptographic building blocks.

Copyright © 2000, 2002, 2003, 2004, 2006, 2007, 2008, 2009, 2011 Free Software Foundation, Inc.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU General Public License as published by the Free Software Foundation; either version 2 of the License, or (at your option) any later version. The text of the license can be found in the section entitled “GNU General Public License”.

Short Contents

1	Introduction	1
2	Preparation	3
3	Generalities	9
4	Handler Functions	19
5	Symmetric cryptography	23
6	Public Key cryptography	33
7	Hashing	55
8	Key Derivation	63
9	Random Numbers	65
10	S-expressions	67
11	MPI library	71
12	Prime numbers	77
13	Utilities	79
14	Architecture	81
A	Description of the Self-Tests	89
B	Description of the FIPS Mode	95
	GNU Lesser General Public License	101
	GNU General Public License	111
	List of Figures and Tables	117
	Concept Index	119
	Function and Data Index	121

Table of Contents

1	Introduction	1
1.1	Getting Started	1
1.2	Features	1
1.3	Overview	1
2	Preparation	3
2.1	Header	3
2.2	Building sources	3
2.3	Building sources using Automake	4
2.4	Initializing the library	4
2.5	Multi-Threading	6
2.6	How to enable the FIPS mode	7
3	Generalities	9
3.1	Controlling the library	9
3.2	Modules	13
3.3	Error Handling	13
3.3.1	Error Values	14
3.3.2	Error Sources	15
3.3.3	Error Codes	16
3.3.4	Error Strings	18
4	Handler Functions	19
4.1	Progress handler	19
4.2	Allocation handler	20
4.3	Error handler	20
4.4	Logging handler	21
5	Symmetric cryptography	23
5.1	Available ciphers	23
5.2	Cipher modules	24
5.3	Available cipher modes	26
5.4	Working with cipher handles	27
5.5	General cipher functions	30
6	Public Key cryptography	33
6.1	Available algorithms	33
6.2	Used S-expressions	33
6.2.1	RSA key parameters	33
6.2.2	DSA key parameters	34
6.2.3	ECC key parameters	34

6.3	Public key modules	36
6.4	Cryptographic Functions	38
6.5	General public-key related Functions	41
6.6	Alternative Public Key Interface	46
6.6.1	Available asymmetric algorithms	46
6.6.2	Working with sets of data	46
6.6.3	Working with IO objects	48
6.6.4	Working with handles	49
6.6.5	Working with keys	50
6.6.6	Using cryptographic functions	51
6.6.7	Handle-independent functions	54
7	Hashing	55
7.1	Available hash algorithms	55
7.2	Hash algorithm modules	56
7.3	Working with hash algorithms	58
8	Key Derivation	63
9	Random Numbers	65
9.1	Quality of random numbers	65
9.2	Retrieving random numbers	65
10	S-expressions	67
10.1	Data types for S-expressions	67
10.2	Working with S-expressions	67
11	MPI library	71
11.1	Data types	71
11.2	Basic functions	71
11.3	MPI formats	72
11.4	Calculations	73
11.5	Comparisons	74
11.6	Bit manipulations	74
11.7	Miscellaneous	75
12	Prime numbers	77
12.1	Generation	77
12.2	Checking	77
13	Utilities	79
13.1	Memory allocation	79

14	Architecture	81
14.1	Public-Key Architecture	82
14.2	Symmetric Encryption Subsystem Architecture	83
14.3	Hashing and MACing Subsystem Architecture	83
14.4	Multi-Precision-Integer Subsystem Architecture	84
14.5	Prime-Number-Generator Subsystem Architecture	84
14.6	Random-Number Subsystem Architecture	85
14.6.1	Description of the CSPRNG	86
14.6.2	Description of the FIPS X9.31 PRNG	86
Appendix A	Description of the Self-Tests	89
A.1	Power-Up Tests	89
A.1.1	Symmetric Cipher Algorithm Power-Up Tests	89
A.1.2	Hash Algorithm Power-Up Tests	89
A.1.3	MAC Algorithm Power-Up Tests	90
A.1.4	Random Number Power-Up Test	90
A.1.5	Public Key Algorithm Power-Up Tests	90
A.1.6	Integrity Power-Up Tests	91
A.1.7	Critical Functions Power-Up Tests	91
A.2	Conditional Tests	91
A.2.1	Key-Pair Generation Tests	91
A.2.2	Software Load Tests	92
A.2.3	Manual Key Entry Tests	92
A.2.4	Continuous RNG Tests	92
A.3	Application Requested Tests	92
A.3.1	Symmetric Cipher Algorithm Tests	92
A.3.2	Hash Algorithm Tests	92
A.3.3	MAC Algorithm Tests	93
Appendix B	Description of the FIPS Mode	95
B.1	Restrictions in FIPS Mode	95
B.2	FIPS Finite State Machine	96
B.3	FIPS Miscellaneous Information	100
	GNU Lesser General Public License	101
	GNU General Public License	111
	List of Figures and Tables	117
	Concept Index	119
	Function and Data Index	121

1 Introduction

Libgcrypt is a library providing cryptographic building blocks.

1.1 Getting Started

This manual documents the Libgcrypt library application programming interface (API). All functions and data types provided by the library are explained.

The reader is assumed to possess basic knowledge about applied cryptography.

This manual can be used in several ways. If read from the beginning to the end, it gives a good introduction into the library and how it can be used in an application. Forward references are included where necessary. Later on, the manual can be used as a reference manual to get just the information needed about any particular interface of the library. Experienced programmers might want to start looking at the examples at the end of the manual, and then only read up those parts of the interface which are unclear.

1.2 Features

Libgcrypt might have a couple of advantages over other libraries doing a similar job.

It's Free Software

Anybody can use, modify, and redistribute it under the terms of the GNU Lesser General Public License (see [\[Library Copying\]](#), page 101). Note, that some parts (which are in general not needed by applications) are subject to the terms of the GNU General Public License (see [\[Copying\]](#), page 111); please see the README file of the distribution for of list of these parts.

It encapsulates the low level cryptography

Libgcrypt provides a high level interface to cryptographic building blocks using an extensible and flexible API.

1.3 Overview

The Libgcrypt library is fully thread-safe, where it makes sense to be thread-safe. Not thread-safe are some cryptographic functions that modify a certain context stored in handles. If the user really intends to use such functions from different threads on the same handle, he has to take care of the serialization of such functions himself. If not described otherwise, every function is thread-safe.

Libgcrypt depends on the library 'libgpg-error', which contains common error handling related code for GnuPG components.

2 Preparation

To use Libgcrypt, you have to perform some changes to your sources and the build system. The necessary changes are small and explained in the following sections. At the end of this chapter, it is described how the library is initialized, and how the requirements of the library are verified.

2.1 Header

All interfaces (data types and functions) of the library are defined in the header file `'gcrypt.h'`. You must include this in all source files using the library, either directly or through some other header file, like this:

```
#include <gcrypt.h>
```

The name space of Libgcrypt is `gcry_*` for function and type names and `GCRY*` for other symbols. In addition the same name prefixes with one prepended underscore are reserved for internal use and should never be used by an application. Note that Libgcrypt uses `libgpg-error`, which uses `gpg_*` as name space for function and type names and `GPG_*` for other symbols, including all the error codes.

Certain parts of `gcrypt.h` may be excluded by defining these macros:

`GCRYPT_NO_MPI_MACROS`

Do not define the shorthand macros `mpi_*` for `gcry_mpi_*`.

`GCRYPT_NO_DEPRECATED`

Do not include definitions for deprecated features. This is useful to make sure that no deprecated features are used.

2.2 Building sources

If you want to compile a source file including the `'gcrypt.h'` header file, you must make sure that the compiler can find it in the directory hierarchy. This is accomplished by adding the path to the directory in which the header file is located to the compilers include file search path (via the `'-I'` option).

However, the path to the include file is determined at the time the source is configured. To solve this problem, Libgcrypt ships with a small helper program `libgcrypt-config` that knows the path to the include file and other configuration options. The options that need to be added to the compiler invocation at compile time are output by the `'--cflags'` option to `libgcrypt-config`. The following example shows how it can be used at the command line:

```
gcc -c foo.c 'libgcrypt-config --cflags'
```

Adding the output of `'libgcrypt-config --cflags'` to the compilers command line will ensure that the compiler can find the Libgcrypt header file.

A similar problem occurs when linking the program with the library. Again, the compiler has to find the library files. For this to work, the path to the library files has to be added to the library search path (via the `'-L'` option). For this, the option `'--libs'` to `libgcrypt-config` can be used. For convenience, this option also outputs all other options that are required to link the program with the Libgcrypt libraries (in particular, the `'-lgcrypt'`

option). The example shows how to link ‘foo.o’ with the Libgcrypt library to a program foo.

```
gcc -o foo foo.o 'libgcrypt-config --libs'
```

Of course you can also combine both examples to a single command by specifying both options to `libgcrypt-config`:

```
gcc -o foo foo.c 'libgcrypt-config --cflags --libs'
```

2.3 Building sources using Automake

It is much easier if you use GNU Automake instead of writing your own Makefiles. If you do that, you do not have to worry about finding and invoking the `libgcrypt-config` script at all. Libgcrypt provides an extension to Automake that does all the work for you.

`AM_PATH_LIBGCRYPT` (*[minimum-version]*, *[action-if-found]*, [Macro]
[action-if-not-found])

Check whether Libgcrypt (at least version *minimum-version*, if given) exists on the host system. If it is found, execute *action-if-found*, otherwise do *action-if-not-found*, if given.

Additionally, the function defines `LIBGCRYPT_CFLAGS` to the flags needed for compilation of the program to find the ‘`gcrypt.h`’ header file, and `LIBGCRYPT_LIBS` to the linker flags needed to link the program to the Libgcrypt library.

You can use the defined Autoconf variables like this in your ‘`Makefile.am`’:

```
AM_CPPFLAGS = $(LIBGCRYPT_CFLAGS)
LDADD = $(LIBGCRYPT_LIBS)
```

2.4 Initializing the library

Before the library can be used, it must initialize itself. This is achieved by invoking the function `gcry_check_version` described below.

Also, it is often desirable to check that the version of Libgcrypt used is indeed one which fits all requirements. Even with binary compatibility, new features may have been introduced, but due to problem with the dynamic linker an old version may actually be used. So you may want to check that the version is okay right after program startup.

`const char * gcry_check_version` (*const char *req_version*) [Function]

The function `gcry_check_version` initializes some subsystems used by Libgcrypt and must be invoked before any other function in the library, with the exception of the `GCRYCTL_SET_THREAD_CBS` command (called via the `gcry_control` function). See [Section 2.5 \[Multi-Threading\]](#), page 6.

Furthermore, this function returns the version number of the library. It can also verify that the version number is higher than a certain required version number *req_version*, if this value is not a null pointer.

Libgcrypt uses a concept known as secure memory, which is a region of memory set aside for storing sensitive data. Because such memory is a scarce resource, it needs to be setup in advanced to a fixed size. Further, most operating systems have special requirements on

how that secure memory can be used. For example, it might be required to install an application as “setuid(root)” to allow allocating such memory. Libgcrypt requires a sequence of initialization steps to make sure that this works correctly. The following examples show the necessary steps.

If you don't have a need for secure memory, for example if your application does not use secret keys or other confidential data or it runs in a controlled environment where key material floating around in memory is not a problem, you should initialize Libgcrypt this way:

```
/* Version check should be the very first call because it
   makes sure that important subsystems are intialized. */
if (!gcry_check_version (GCRYPT_VERSION))
{
    fputs ("libgcrypt version mismatch\n", stderr);
    exit (2);
}

/* Disable secure memory. */
gcry_control (GCRYCTL_DISABLE_SECMEM, 0);

/* ... If required, other initialization goes here. */

/* Tell Libgcrypt that initialization has completed. */
gcry_control (GCRYCTL_INITIALIZATION_FINISHED, 0);
```

If you have to protect your keys or other information in memory against being swapped out to disk and to enable an automatic overwrite of used and freed memory, you need to initialize Libgcrypt this way:

```
/* Version check should be the very first call because it
   makes sure that important subsystems are intialized. */
if (!gcry_check_version (GCRYPT_VERSION))
{
    fputs ("libgcrypt version mismatch\n", stderr);
    exit (2);
}

/* We don't want to see any warnings, e.g. because we have not yet
   parsed program options which might be used to suppress such
   warnings. */
gcry_control (GCRYCTL_SUSPEND_SECMEM_WARN);

/* ... If required, other initialization goes here. Note that the
   process might still be running with increased privileges and that
   the secure memory has not been intialized. */

/* Allocate a pool of 16k secure memory. This make the secure memory
   available and also drops privileges where needed. */
```

```

gcry_control (GCRYCTL_INIT_SECMEM, 16384, 0);

/* It is now okay to let Libgcrypt complain when there was/is
   a problem with the secure memory. */
gcry_control (GCRYCTL_RESUME_SECMEM_WARN);

/* ... If required, other initialization goes here. */

/* Tell Libgcrypt that initialization has completed. */
gcry_control (GCRYCTL_INITIALIZATION_FINISHED, 0);

```

It is important that these initialization steps are not done by a library but by the actual application. A library using Libgcrypt might want to check for finished initialization using:

```

if (!gcry_control (GCRYCTL_INITIALIZATION_FINISHED_P))
{
    fputs ("libgcrypt has not been initialized\n", stderr);
    abort ();
}

```

Instead of terminating the process, the library may instead print a warning and try to initialize Libgcrypt itself. See also the section on multi-threading below for more pitfalls.

2.5 Multi-Threading

As mentioned earlier, the Libgcrypt library is thread-safe if you adhere to the following requirements:

- If your application is multi-threaded, you must set the thread support callbacks with the `GCRYCTL_SET_THREAD_CBS` command **before** any other function in the library.

This is easy enough if you are indeed writing an application using Libgcrypt. It is rather problematic if you are writing a library instead. Here are some tips what to do if you are writing a library:

If your library requires a certain thread package, just initialize Libgcrypt to use this thread package. If your library supports multiple thread packages, but needs to be configured, you will have to implement a way to determine which thread package the application wants to use with your library anyway. Then configure Libgcrypt to use this thread package.

If your library is fully reentrant without any special support by a thread package, then you are lucky indeed. Unfortunately, this does not relieve you from doing either of the two above, or use a third option. The third option is to let the application initialize Libgcrypt for you. Then you are not using Libgcrypt transparently, though.

As if this was not difficult enough, a conflict may arise if two libraries try to initialize Libgcrypt independently of each others, and both such libraries are then linked into the same application. To make it a bit simpler for you, this will probably work, but only if both libraries have the same requirement for the thread package. This is currently only supported for the non-threaded case, GNU Pth and pthread.

If you use pthread and your applications forks and does not directly call exec (even calling stdio functions), all kind of problems may occur. Future versions of Libgcrypt

will try to cleanup using `pthread_atfork` but even that may lead to problems. This is a common problem with almost all applications using `pthread` and `fork`.

Note that future versions of Libgcrypt will drop this flexible thread support and instead only support the platforms standard thread implementation.

- The function `gcry_check_version` must be called before any other function in the library, except the `GCRYCTL_SET_THREAD_CBS` command (called via the `gcry_control` function), because it initializes the thread support subsystem in Libgcrypt. To achieve this in multi-threaded programs, you must synchronize the memory with respect to other threads that also want to use Libgcrypt. For this, it is sufficient to call `gcry_check_version` before creating the other threads using Libgcrypt¹.
- Just like the function `gpg_strerror`, the function `gcry_strerror` is not thread safe. You have to use `gpg_strerror_r` instead.

Libgcrypt contains convenient macros, which define the necessary thread callbacks for PThread and for GNU Pth:

`GCRY_THREAD_OPTION_PTH_IMPL`

This macro defines the following (static) symbols: `gcry_pth_init`, `gcry_pth_mutex_init`, `gcry_pth_mutex_destroy`, `gcry_pth_mutex_lock`, `gcry_pth_mutex_unlock`, `gcry_pth_read`, `gcry_pth_write`, `gcry_pth_select`, `gcry_pth_waitpid`, `gcry_pth_accept`, `gcry_pth_connect`, `gcry_threads_pth`.

After including this macro, `gcry_control()` shall be used with a command of `GCRYCTL_SET_THREAD_CBS` in order to register the thread callback structure named “`gcry_threads_pth`”. Example:

```
ret = gcry_control (GCRYCTL_SET_THREAD_CBS, &gcry_threads_pth);
```

`GCRY_THREAD_OPTION_PTHREAD_IMPL`

This macro defines the following (static) symbols: `gcry_pthread_mutex_init`, `gcry_pthread_mutex_destroy`, `gcry_pthread_mutex_lock`, `gcry_pthread_mutex_unlock`, `gcry_threads_pthread`.

After including this macro, `gcry_control()` shall be used with a command of `GCRYCTL_SET_THREAD_CBS` in order to register the thread callback structure named “`gcry_threads_pthread`”. Example:

```
ret = gcry_control (GCRYCTL_SET_THREAD_CBS, &gcry_threads_pthread);
```

Note that these macros need to be terminated with a semicolon. Keep in mind that these are convenient macros for C programmers; C++ programmers might have to wrap these macros in an “extern C” body.

2.6 How to enable the FIPS mode

Libgcrypt may be used in a FIPS 140-2 mode. Note, that this does not necessary mean that Libgcrypt is an approved FIPS 140-2 module. Check the NIST database

¹ At least this is true for POSIX threads, as `pthread_create` is a function that synchronizes memory with respects to other threads. There are many functions which have this property, a complete list can be found in POSIX, IEEE Std 1003.1-2003, Base Definitions, Issue 6, in the definition of the term “Memory Synchronization”. For other thread packages, more relaxed or more strict rules may apply.

at <http://csrc.nist.gov/groups/STM/cmvp/> to see what versions of Libgcrypt are approved.

Because FIPS 140 has certain restrictions on the use of cryptography which are not always wanted, Libgcrypt needs to be put into FIPS mode explicitly. Three alternative mechanisms are provided to switch Libgcrypt into this mode:

- If the file `‘/proc/sys/crypto/fips_enabled’` exists and contains a numeric value other than 0, Libgcrypt is put into FIPS mode at initialization time. Obviously this works only on systems with a `proc` file system (i.e. GNU/Linux).
- If the file `‘/etc/gcrypt/fips_enabled’` exists, Libgcrypt is put into FIPS mode at initialization time. Note that this filename is hardwired and does not depend on any configuration options.
- If the application requests FIPS mode using the control command `GCRYCTL_FORCE_FIPS_MODE`. This must be done prior to any initialization (i.e. before `gcry_check_version`).

In addition to the standard FIPS mode, Libgcrypt may also be put into an Enforced FIPS mode by writing a non-zero value into the file `‘/etc/gcrypt/fips_enabled’`. The Enforced FIPS mode helps to detect applications which don’t fulfill all requirements for using Libgcrypt in FIPS mode (see [Appendix B \[FIPS Mode\]](#), page 95).

Once Libgcrypt has been put into FIPS mode, it is not possible to switch back to standard mode without terminating the process first. If the logging verbosity level of Libgcrypt has been set to at least 2, the state transitions and the self-tests are logged.

3 Generalities

3.1 Controlling the library

`gcry_error_t gcry_control (enum gcry_ctl_cmds cmd, ...)` [Function]

This function can be used to influence the general behavior of Libgcrypt in several ways. Depending on *cmd*, more arguments can or have to be provided.

GCRYCTL_ENABLE_M_GUARD; Arguments: none

This command enables the built-in memory guard. It must not be used to activate the memory guard after the memory management has already been used; therefore it can **ONLY** be used before `gcry_check_version`. Note that the memory guard is **NOT** used when the user of the library has set his own memory management callbacks.

GCRYCTL_ENABLE_QUICK_RANDOM; Arguments: none

This command inhibits the use the very secure random quality level (`GCRY_VERY_STRONG_RANDOM`) and degrades all request down to `GCRY_STRONG_RANDOM`. In general this is not recommended. However, for some applications the extra quality random Libgcrypt tries to create is not justified and this option may help to get better performance. Please check with a crypto expert whether this option can be used for your application. This option can only be used at initialization time.

GCRYCTL_DUMP_RANDOM_STATS; Arguments: none

This command dumps random number generator related statistics to the library's logging stream.

GCRYCTL_DUMP_MEMORY_STATS; Arguments: none

This command dumps memory management related statistics to the library's logging stream.

GCRYCTL_DUMP_SECMEM_STATS; Arguments: none

This command dumps secure memory management related statistics to the library's logging stream.

GCRYCTL_DROP_PRIVS; Arguments: none

This command disables the use of secure memory and drops the privileges of the current process. This command has not much use; the suggested way to disable secure memory is to use `GCRYCTL_DISABLE_SECMEM` right after initialization.

GCRYCTL_DISABLE_SECMEM; Arguments: none

This command disables the use of secure memory. If this command is used in FIPS mode, FIPS mode will be disabled and the function `gcry_fips_mode_active` returns false. However, in Enforced FIPS mode this command has no effect at all.

Many applications do not require secure memory, so they should disable it right away. This command should be executed right after `gcry_check_version`.

GCRYCTL_INIT_SECMEM; Arguments: int nbytes

This command is used to allocate a pool of secure memory and thus enabling the use of secure memory. It also drops all extra privileges the process has (i.e. if it is run as `setuid (root)`). If the argument *nbytes* is 0, secure memory will be disabled. The minimum amount of secure memory allocated is currently 16384 bytes; you may thus use a value of 1 to request that default size.

GCRYCTL_TERM_SECMEM; Arguments: none

This command zeroes the secure memory and destroys the handler. The secure memory pool may not be used anymore after running this command. If the secure memory pool has already been destroyed, this command has no effect. Applications might want to run this command from their exit handler to make sure that the secure memory gets properly destroyed. This command is not necessarily thread-safe but that should not be needed in cleanup code. It may be called from a signal handler.

GCRYCTL_DISABLE_SECMEM_WARN; Arguments: none

Disable warning messages about problems with the secure memory subsystem. This command should be run right after `gcry_check_version`.

GCRYCTL_SUSPEND_SECMEM_WARN; Arguments: none

Postpone warning messages from the secure memory subsystem. See [\[the initialization example\]](#), page 5, on how to use it.

GCRYCTL_RESUME_SECMEM_WARN; Arguments: none

Resume warning messages from the secure memory subsystem. See [\[the initialization example\]](#), page 6, on how to use it.

GCRYCTL_USE_SECURE_RNDPOOL; Arguments: none

This command tells the PRNG to store random numbers in secure memory. This command should be run right after `gcry_check_version` and not later than the command `GCRYCTL_INIT_SECMEM`. Note that in FIPS mode the secure memory is always used.

GCRYCTL_SET_RANDOM_SEED_FILE; Arguments: const char *filename

This command specifies the file, which is to be used as seed file for the PRNG. If the seed file is registered prior to initialization of the PRNG, the seed file's content (if it exists and seems to be valid) is fed into the PRNG pool. After the seed file has been registered, the PRNG can be signalled to write out the PRNG pool's content into the seed file with the following command.

GCRYCTL_UPDATE_RANDOM_SEED_FILE; Arguments: none

Write out the PRNG pool's content into the registered seed file.

Multiple instances of the applications sharing the same random seed file can be started in parallel, in which case they will read out the same pool and then race for updating it (the last update overwrites earlier updates). They will differentiate only by the weak entropy that is added in `read_seed_file` based on the PID and clock, and up to 16 bytes of weak random non-blockingly. The consequence is that the output of

these different instances is correlated to some extent. In a perfect attack scenario, the attacker can control (or at least guess) the PID and clock of the application, and drain the system's entropy pool to reduce the "up to 16 bytes" above to 0. Then the dependencies of the initial states of the pools are completely known. Note that this is not an issue if random of `GCRY_VERY_STRONG_RANDOM` quality is requested as in this case enough extra entropy gets mixed. It is also not an issue when using Linux (rndlinux driver), because this one guarantees to read full 16 bytes from `/dev/urandom` and thus there is no way for an attacker without kernel access to control these 16 bytes.

`GCRYCTL_SET_VERBOSITY`; Arguments: int level

This command sets the verbosity of the logging. A level of 0 disables all extra logging whereas positive numbers enable more verbose logging. The level may be changed at any time but be aware that no memory synchronization is done so the effect of this command might not immediately show up in other threads. This command may even be used prior to `gcry_check_version`.

`GCRYCTL_SET_DEBUG_FLAGS`; Arguments: unsigned int flags

Set the debug flag bits as given by the argument. Be aware that that no memory synchronization is done so the effect of this command might not immediately show up in other threads. The debug flags are not considered part of the API and thus may change without notice. As of now bit 0 enables debugging of cipher functions and bit 1 debugging of multi-precision-integers. This command may even be used prior to `gcry_check_version`.

`GCRYCTL_CLEAR_DEBUG_FLAGS`; Arguments: unsigned int flags

Set the debug flag bits as given by the argument. Be aware that that no memory synchronization is done so the effect of this command might not immediately show up in other threads. This command may even be used prior to `gcry_check_version`.

`GCRYCTL_DISABLE_INTERNAL_LOCKING`; Arguments: none

This command does nothing. It exists only for backward compatibility.

`GCRYCTL_ANY_INITIALIZATION_P`; Arguments: none

This command returns true if the library has been basically initialized. Such a basic initialization happens implicitly with many commands to get certain internal subsystems running. The common and suggested way to do this basic initialization is by calling `gcry_check_version`.

`GCRYCTL_INITIALIZATION_FINISHED`; Arguments: none

This command tells the library that the application has finished the initialization.

`GCRYCTL_INITIALIZATION_FINISHED_P`; Arguments: none

This command returns true if the command `GCRYCTL_INITIALIZATION_FINISHED` has already been run.

- GCRYCTL_SET_THREAD_CBS**; Arguments: `struct ath_ops *ath_ops`
 This command registers a thread-callback structure. See [Section 2.5 \[Multi-Threading\]](#), page 6.
- GCRYCTL_FAST_POLL**; Arguments: none
 Run a fast random poll.
- GCRYCTL_SET_RNDEGD_SOCKET**; Arguments: `const char *filename`
 This command may be used to override the default name of the EGD socket to connect to. It may be used only during initialization as it is not thread safe. Changing the socket name again is not supported. The function may return an error if the given filename is too long for a local socket name.
 EGD is an alternative random gatherer, used only on systems lacking a proper random device.
- GCRYCTL_PRINT_CONFIG**; Arguments: `FILE *stream`
 This command dumps information pertaining to the configuration of the library to the given stream. If NULL is given for *stream*, the log system is used. This command may be used before the initialization has been finished but not before a `gcry_version_check`.
- GCRYCTL_OPERATIONAL_P**; Arguments: none
 This command returns true if the library is in an operational state. This information makes only sense in FIPS mode. In contrast to other functions, this is a pure test function and won't put the library into FIPS mode or change the internal state. This command may be used before the initialization has been finished but not before a `gcry_version_check`.
- GCRYCTL_FIPS_MODE_P**; Arguments: none
 This command returns true if the library is in FIPS mode. Note, that this is no indication about the current state of the library. This command may be used before the initialization has been finished but not before a `gcry_version_check`. An application may use this command or the convenience macro below to check whether FIPS mode is actually active.
- ```
int gcry_fips_mode_active (void) [Function]
```
- Returns true if the FIPS mode is active. Note that this is implemented as a macro.
- GCRYCTL\_FORCE\_FIPS\_MODE**; Arguments: none  
 Running this command puts the library into FIPS mode. If the library is already in FIPS mode, a self-test is triggered and thus the library will be put into operational state. This command may be used before a call to `gcry_check_version` and that is actually the recommended way to let an application switch the library into FIPS mode. Note that Libcrypt will reject an attempt to switch to fips mode during or after the initialization.
- GCRYCTL\_SELFTEST**; Arguments: none  
 This may be used at anytime to have the library run all implemented self-tests. It works in standard and in FIPS mode. Returns 0 on success or an error code on failure.

`GCRYCTL_DISABLE_HWF`; Arguments: `const char *name`

Libgcrypt detects certain features of the CPU at startup time. For performance tests it is sometimes required not to use such a feature. This option may be used to disable a certain feature; i.e. Libgcrypt behaves as if this feature has not been detected. Note that the detection code might be run if the feature has been disabled. This command must be used at initialization time; i.e. before calling `gcry_check_version`.

## 3.2 Modules

Libgcrypt supports the use of ‘extension modules’, which implement algorithms in addition to those already built into the library directly.

`gcry_module_t`

[Data type]

This data type represents a ‘module’.

Functions registering modules provided by the user take a ‘module specification structure’ as input and return a value of `gcry_module_t` and an ID that is unique in the modules’ category. This ID can be used to reference the newly registered module. After registering a module successfully, the new functionality should be able to be used through the normal functions provided by Libgcrypt until it is unregistered again.

## 3.3 Error Handling

Many functions in Libgcrypt can return an error if they fail. For this reason, the application should always catch the error condition and take appropriate measures, for example by releasing the resources and passing the error up to the caller, or by displaying a descriptive message to the user and cancelling the operation.

Some error values do not indicate a system error or an error in the operation, but the result of an operation that failed properly. For example, if you try to decrypt a tempered message, the decryption will fail. Another error value actually means that the end of a data buffer or list has been reached. The following descriptions explain for many error codes what they mean usually. Some error values have specific meanings if returned by a certain functions. Such cases are described in the documentation of those functions.

Libgcrypt uses the `libgpg-error` library. This allows to share the error codes with other components of the GnuPG system, and to pass error values transparently from the crypto engine, or some helper application of the crypto engine, to the user. This way no information is lost. As a consequence, Libgcrypt does not use its own identifiers for error codes, but uses those provided by `libgpg-error`. They usually start with `GPG_ERR_`.

However, Libgcrypt does provide aliases for the functions defined in `libgpg-error`, which might be preferred for name space consistency.

Most functions in Libgcrypt return an error code in the case of failure. For this reason, the application should always catch the error condition and take appropriate measures, for example by releasing the resources and passing the error up to the caller, or by displaying a descriptive message to the user and canceling the operation.

Some error values do not indicate a system error or an error in the operation, but the result of an operation that failed properly.

GnuPG components, including Libgcrypt, use an extra library named libgpg-error to provide a common error handling scheme. For more information on libgpg-error, see the according manual.

### 3.3.1 Error Values

`gcry_err_code_t` [Data type]

The `gcry_err_code_t` type is an alias for the libgpg-error type `gpg_err_code_t`. The error code indicates the type of an error, or the reason why an operation failed. A list of important error codes can be found in the next section.

`gcry_err_source_t` [Data type]

The `gcry_err_source_t` type is an alias for the libgpg-error type `gpg_err_source_t`. The error source has not a precisely defined meaning. Sometimes it is the place where the error happened, sometimes it is the place where an error was encoded into an error value. Usually the error source will give an indication to where to look for the problem. This is not always true, but it is attempted to achieve this goal.

A list of important error sources can be found in the next section.

`gcry_error_t` [Data type]

The `gcry_error_t` type is an alias for the libgpg-error type `gpg_error_t`. An error value like this has always two components, an error code and an error source. Both together form the error value.

Thus, the error value can not be directly compared against an error code, but the accessor functions described below must be used. However, it is guaranteed that only 0 is used to indicate success (`GPG_ERR_NO_ERROR`), and that in this case all other parts of the error value are set to 0, too.

Note that in Libgcrypt, the error source is used purely for diagnostic purposes. Only the error code should be checked to test for a certain outcome of a function. The manual only documents the error code part of an error value. The error source is left unspecified and might be anything.

`gcry_err_code_t gcry_err_code (gcry_error_t err)` [Function]

The static inline function `gcry_err_code` returns the `gcry_err_code_t` component of the error value `err`. This function must be used to extract the error code from an error value in order to compare it with the `GPG_ERR_*` error code macros.

`gcry_err_source_t gcry_err_source (gcry_error_t err)` [Function]

The static inline function `gcry_err_source` returns the `gcry_err_source_t` component of the error value `err`. This function must be used to extract the error source from an error value in order to compare it with the `GPG_ERR_SOURCE_*` error source macros.

`gcry_error_t gcry_err_make (gcry_err_source_t source, gcry_err_code_t code)` [Function]

The static inline function `gcry_err_make` returns the error value consisting of the error source `source` and the error code `code`.

This function can be used in callback functions to construct an error value to return it to the library.

`gcry_error_t gcry_error (gcry_err_code_t code)` [Function]

The static inline function `gcry_error` returns the error value consisting of the default error source and the error code `code`.

For GCRY applications, the default error source is `GPG_ERR_SOURCE_USER_1`. You can define `GCRY_ERR_SOURCE_DEFAULT` before including ‘`gcrypt.h`’ to change this default.

This function can be used in callback functions to construct an error value to return it to the library.

The `libpgp-error` library provides error codes for all system error numbers it knows about. If `err` is an unknown error number, the error code `GPG_ERR_UNKNOWN_ERRNO` is used. The following functions can be used to construct error values from system `errno` numbers.

`gcry_error_t gcry_err_make_from_errno` [Function]  
(`gcry_err_source_t source`, `int err`)

The function `gcry_err_make_from_errno` is like `gcry_err_make`, but it takes a system error like `errno` instead of a `gcry_err_code_t` error code.

`gcry_error_t gcry_error_from_errno (int err)` [Function]

The function `gcry_error_from_errno` is like `gcry_error`, but it takes a system error like `errno` instead of a `gcry_err_code_t` error code.

Sometimes you might want to map system error numbers to error codes directly, or map an error code representing a system error back to the system error number. The following functions can be used to do that.

`gcry_err_code_t gcry_err_code_from_errno (int err)` [Function]

The function `gcry_err_code_from_errno` returns the error code for the system error `err`. If `err` is not a known system error, the function returns `GPG_ERR_UNKNOWN_ERRNO`.

`int gcry_err_code_to_errno (gcry_err_code_t err)` [Function]

The function `gcry_err_code_to_errno` returns the system error for the error code `err`. If `err` is not an error code representing a system error, or if this system error is not defined on this system, the function returns 0.

### 3.3.2 Error Sources

The library `libpgp-error` defines an error source for every component of the GnuPG system. The error source part of an error value is not well defined. As such it is mainly useful to improve the diagnostic error message for the user.

If the error code part of an error value is 0, the whole error value will be 0. In this case the error source part is of course `GPG_ERR_SOURCE_UNKNOWN`.

The list of error sources that might occur in applications using Libgcrypt is:

`GPG_ERR_SOURCE_UNKNOWN`

The error source is not known. The value of this error source is 0.

`GPG_ERR_SOURCE_GPGME`

The error source is GPGME itself.

**GPG\_ERR\_SOURCE\_GPG**

The error source is GnuPG, which is the crypto engine used for the OpenPGP protocol.

**GPG\_ERR\_SOURCE\_GPGSM**

The error source is GPGSM, which is the crypto engine used for the OpenPGP protocol.

**GPG\_ERR\_SOURCE\_GCRYPT**

The error source is `libgcrypt`, which is used by crypto engines to perform cryptographic operations.

**GPG\_ERR\_SOURCE\_GPGAGENT**

The error source is `gpg-agent`, which is used by crypto engines to perform operations with the secret key.

**GPG\_ERR\_SOURCE\_PINENTRY**

The error source is `pinentry`, which is used by `gpg-agent` to query the passphrase to unlock a secret key.

**GPG\_ERR\_SOURCE\_SCD**

The error source is the SmartCard Daemon, which is used by `gpg-agent` to delegate operations with the secret key to a SmartCard.

**GPG\_ERR\_SOURCE\_KEYBOX**

The error source is `libkbx`, a library used by the crypto engines to manage local keyrings.

**GPG\_ERR\_SOURCE\_USER\_1****GPG\_ERR\_SOURCE\_USER\_2****GPG\_ERR\_SOURCE\_USER\_3****GPG\_ERR\_SOURCE\_USER\_4**

These error sources are not used by any GnuPG component and can be used by other software. For example, applications using Libgcrypt can use them to mark error values coming from callback handlers. Thus `GPG_ERR_SOURCE_USER_1` is the default for errors created with `gcry_error` and `gcry_error_from_errno`, unless you define `GCRY_ERR_SOURCE_DEFAULT` before including `'gcrypt.h'`.

### 3.3.3 Error Codes

The library `libgpg-error` defines many error values. The following list includes the most important error codes.

**GPG\_ERR\_EOF**

This value indicates the end of a list, buffer or file.

**GPG\_ERR\_NO\_ERROR**

This value indicates success. The value of this error code is 0. Also, it is guaranteed that an error value made from the error code 0 will be 0 itself (as a whole). This means that the error source information is lost for this error code, however, as this error code indicates that no error occurred, this is generally not a problem.



**GPG\_ERR\_GENERAL**

This value means that something went wrong, but either there is not enough information about the problem to return a more useful error value, or there is no separate error value for this type of problem.

**GPG\_ERR\_ENOMEM**

This value means that an out-of-memory condition occurred.

**GPG\_ERR\_E...**

System errors are mapped to GPG\_ERR\_EFOO where FOO is the symbol for the system error.

**GPG\_ERR\_INV\_VALUE**

This value means that some user provided data was out of range.

**GPG\_ERR\_UNUSABLE\_PUBKEY**

This value means that some recipients for a message were invalid.

**GPG\_ERR\_UNUSABLE\_SECKEY**

This value means that some signers were invalid.

**GPG\_ERR\_NO\_DATA**

This value means that data was expected where no data was found.

**GPG\_ERR\_CONFLICT**

This value means that a conflict of some sort occurred.

**GPG\_ERR\_NOT\_IMPLEMENTED**

This value indicates that the specific function (or operation) is not implemented. This error should never happen. It can only occur if you use certain values or configuration options which do not work, but for which we think that they should work at some later time.

**GPG\_ERR\_DECRYPT\_FAILED**

This value indicates that a decryption operation was unsuccessful.

**GPG\_ERR\_WRONG\_KEY\_USAGE**

This value indicates that a key is not used appropriately.

**GPG\_ERR\_NO\_SECKEY**

This value indicates that no secret key for the user ID is available.

**GPG\_ERR\_UNSUPPORTED\_ALGORITHM**

This value means a verification failed because the cryptographic algorithm is not supported by the crypto backend.

**GPG\_ERR\_BAD\_SIGNATURE**

This value means a verification failed because the signature is bad.

**GPG\_ERR\_NO\_PUBKEY**

This value means a verification failed because the public key is not available.

**GPG\_ERR\_NOT\_OPERATIONAL**

This value means that the library is not yet in state which allows to use this function. This error code is in particular returned if Libgcrypt is operated in

FIPS mode and the internal state of the library does not yet or not anymore allow the use of a service.

This error code is only available with newer libgpg-error versions, thus you might see “invalid error code” when passing this to `gpg_strerror`. The numeric value of this error code is 176.

GPG\_ERR\_USER\_1

GPG\_ERR\_USER\_2

...

GPG\_ERR\_USER\_16

These error codes are not used by any GnuPG component and can be freely used by other software. Applications using Libgcrypt might use them to mark specific errors returned by callback handlers if no suitable error codes (including the system errors) for these errors exist already.

### 3.3.4 Error Strings

`const char * gcry_strerror (gcry_error_t err)` [Function]

The function `gcry_strerror` returns a pointer to a statically allocated string containing a description of the error code contained in the error value `err`. This string can be used to output a diagnostic message to the user.

`const char * gcry_strerror (gcry_error_t err)` [Function]

The function `gcry_strerror` returns a pointer to a statically allocated string containing a description of the error source contained in the error value `err`. This string can be used to output a diagnostic message to the user.

The following example illustrates the use of the functions described above:

```
{
 gcry_cipher_hd_t handle;
 gcry_error_t err = 0;

 err = gcry_cipher_open (&handle, GCRY_CIPHER_AES,
 GCRY_CIPHER_MODE_CBC, 0);

 if (err)
 {
 fprintf (stderr, "Failure: %s/%s\n",
 gcry_strerror (err),
 gcry_strerror (err));
 }
}
```

## 4 Handler Functions

Libgcrypt makes it possible to install so called ‘handler functions’, which get called by Libgcrypt in case of certain events.

### 4.1 Progress handler

It is often useful to retrieve some feedback while long running operations are performed.

`gcry_handler_progress_t` [Data type]  
 Progress handler functions have to be of the type `gcry_handler_progress_t`, which is defined as:

```
void (*gcry_handler_progress_t) (void *, const char *, int, int, int)
```

The following function may be used to register a handler function for this purpose.

`void gcry_set_progress_handler (gcry_handler_progress_t cb, void *cb_data)` [Function]

This function installs `cb` as the ‘Progress handler’ function. It may be used only during initialization. `cb` must be defined as follows:

```
void
my_progress_handler (void *cb_data, const char *what,
 int printchar, int current, int total)
{
 /* Do something. */
}
```

A description of the arguments of the progress handler function follows.

`cb_data` The argument provided in the call to `gcry_set_progress_handler`.

`what` A string identifying the type of the progress output. The following values for `what` are defined:

|                           |                                                                                         |
|---------------------------|-----------------------------------------------------------------------------------------|
| <code>need_entropy</code> | Not enough entropy is available. <code>total</code> holds the number of required bytes. |
| <code>primegen</code>     | Values for <code>printchar</code> :                                                     |
| <code>\n</code>           | Prime generated.                                                                        |
| <code>!</code>            | Need to refresh the pool of prime numbers.                                              |
| <code>&lt;, &gt;</code>   | Number of bits adjusted.                                                                |
| <code>^</code>            | Searching for a generator.                                                              |
| <code>.</code>            | Fermat test on 10 candidates failed.                                                    |
| <code>:</code>            | Restart with a new random value.                                                        |
| <code>+</code>            | Rabin Miller test passed.                                                               |

## 4.2 Allocation handler

It is possible to make Libgcrypt use special memory allocation functions instead of the built-in ones.

Memory allocation functions are of the following types:

`gcry_handler_alloc_t` [Data type]

This type is defined as: `void *(*gcry_handler_alloc_t) (size_t n)`.

`gcry_handler_secure_check_t` [Data type]

This type is defined as: `int *(*gcry_handler_secure_check_t) (const void *)`.

`gcry_handler_realloc_t` [Data type]

This type is defined as: `void *(*gcry_handler_realloc_t) (void *p, size_t n)`.

`gcry_handler_free_t` [Data type]

This type is defined as: `void *(*gcry_handler_free_t) (void *)`.

Special memory allocation functions can be installed with the following function:

`void gcry_set_allocation_handler` (*gcry\_handler\_alloc\_t* [Function]

*func\_alloc*, *gcry\_handler\_alloc\_t func\_alloc\_secure*,  
*gcry\_handler\_secure\_check\_t func\_secure\_check*, *gcry\_handler\_realloc\_t*  
*func\_realloc*, *gcry\_handler\_free\_t func\_free*)

Install the provided functions and use them instead of the built-in functions for doing memory allocation. Using this function is in general not recommended because the standard Libgcrypt allocation functions are guaranteed to zeroize memory if needed.

This function may be used only during initialization and may not be used in fips mode.

## 4.3 Error handler

The following functions may be used to register handler functions that are called by Libgcrypt in case certain error conditions occur. They may and should be registered prior to calling `gcry_check_version`.

`gcry_handler_no_mem_t` [Data type]

This type is defined as: `int (*gcry_handler_no_mem_t) (void *, size_t, unsigned int)`

`void gcry_set_outofcore_handler` (*gcry\_handler\_no\_mem\_t* [Function]

*func\_no\_mem*, *void \*cb\_data*)

This function registers *func\_no\_mem* as ‘out-of-core handler’, which means that it will be called in the case of not having enough memory available. The handler is called with 3 arguments: The first one is the pointer *cb\_data* as set with this function, the second is the requested memory size and the last being a flag. If bit 0 of the flag is set, secure memory has been requested. The handler should either return true to indicate that Libgcrypt should try again allocating memory or return false to let Libgcrypt use its default fatal error handler.

`gcry_handler_error_t` [Data type]  
This type is defined as: `void (*gcry_handler_error_t) (void *, int, const char *)`

`void gcry_set_fatalerror_handler (gcry_handler_error_t func_error, void *cb_data)` [Function]  
This function registers *func\_error* as ‘error handler’, which means that it will be called in error conditions.

## 4.4 Logging handler

`gcry_handler_log_t` [Data type]  
This type is defined as: `void (*gcry_handler_log_t) (void *, int, const char *, va_list)`

`void gcry_set_log_handler (gcry_handler_log_t func_log, void *cb_data)` [Function]  
This function registers *func\_log* as ‘logging handler’, which means that it will be called in case Libgcrypt wants to log a message. This function may and should be used prior to calling `gcry_check_version`.



## 5 Symmetric cryptography

The cipher functions are used for symmetrical cryptography, i.e. cryptography using a shared key. The programming model follows an open/process/close paradigm and is in that similar to other building blocks provided by Libgcrypt.

### 5.1 Available ciphers

#### GCRY\_CIPHER\_NONE

This is not a real algorithm but used by some functions as error return. The value always evaluates to false.

#### GCRY\_CIPHER\_IDEA

This is the IDEA algorithm. The constant is provided but there is currently no implementation for it because the algorithm is patented.

#### GCRY\_CIPHER\_3DES

Triple-DES with 3 Keys as EDE. The key size of this algorithm is 168 but you have to pass 192 bits because the most significant bits of each byte are ignored.

#### GCRY\_CIPHER\_CAST5

CAST128-5 block cipher algorithm. The key size is 128 bits.

#### GCRY\_CIPHER\_BLOWFISH

The blowfish algorithm. The current implementation allows only for a key size of 128 bits.

#### GCRY\_CIPHER\_SAFER\_SK128

Reserved and not currently implemented.

#### GCRY\_CIPHER\_DES\_SK

Reserved and not currently implemented.

#### GCRY\_CIPHER\_AES

#### GCRY\_CIPHER\_AES128

#### GCRY\_CIPHER\_RIJNDAEL

#### GCRY\_CIPHER\_RIJNDAEL128

AES (Rijndael) with a 128 bit key.

#### GCRY\_CIPHER\_AES192

#### GCRY\_CIPHER\_RIJNDAEL192

AES (Rijndael) with a 192 bit key.

#### GCRY\_CIPHER\_AES256

#### GCRY\_CIPHER\_RIJNDAEL256

AES (Rijndael) with a 256 bit key.

#### GCRY\_CIPHER\_TWOFISH

The Twofish algorithm with a 256 bit key.

#### GCRY\_CIPHER\_TWOFISH128

The Twofish algorithm with a 128 bit key.

**GCRY\_CIPHER\_ARCFOUR**

An algorithm which is 100% compatible with RSA Inc.'s RC4 algorithm. Note that this is a stream cipher and must be used very carefully to avoid a couple of weaknesses.

**GCRY\_CIPHER\_DES**

Standard DES with a 56 bit key. You need to pass 64 bit but the high bits of each byte are ignored. Note, that this is a weak algorithm which can be broken in reasonable time using a brute force approach.

**GCRY\_CIPHER\_SERPENT128****GCRY\_CIPHER\_SERPENT192****GCRY\_CIPHER\_SERPENT256**

The Serpent cipher from the AES contest.

**GCRY\_CIPHER\_RFC2268\_40****GCRY\_CIPHER\_RFC2268\_128**

Ron's Cipher 2 in the 40 and 128 bit variants. Note, that we currently only support the 40 bit variant. The identifier for 128 is reserved for future use.

**GCRY\_CIPHER\_SEED**

A 128 bit cipher as described by RFC4269.

**GCRY\_CIPHER\_CAMELLIA128****GCRY\_CIPHER\_CAMELLIA192****GCRY\_CIPHER\_CAMELLIA256**

The Camellia cipher by NTT. See <http://info.isl.ntt.co.jp/crypt/eng/camellia/specifications.html>.

## 5.2 Cipher modules

Libcrypt makes it possible to load additional 'cipher modules'; these ciphers can be used just like the cipher algorithms that are built into the library directly. For an introduction into extension modules, see See [Section 3.2 \[Modules\]](#), page 13.

**gcry\_cipher\_spec\_t**

[Data type]

This is the 'module specification structure' needed for registering cipher modules, which has to be filled in by the user before it can be used to register a module. It contains the following members:

**const char \*name**

The primary name of the algorithm.

**const char \*\*aliases**

A list of strings that are 'aliases' for the algorithm. The list must be terminated with a NULL element.

**gcry\_cipher\_oid\_spec\_t \*oids**

A list of OIDs that are to be associated with the algorithm. The list's last element must have its 'oid' member set to NULL. See below for an explanation of this type.



- size\_t** `blocksize`  
The block size of the algorithm, in bytes.
- size\_t** `keylen`  
The length of the key, in bits.
- size\_t** `contextsize`  
The size of the algorithm-specific ‘context’, that should be allocated for each handle.
- gcry\_cipher\_setkey\_t** `setkey`  
The function responsible for initializing a handle with a provided key. See below for a description of this type.
- gcry\_cipher\_encrypt\_t** `encrypt`  
The function responsible for encrypting a single block. See below for a description of this type.
- gcry\_cipher\_decrypt\_t** `decrypt`  
The function responsible for decrypting a single block. See below for a description of this type.
- gcry\_cipher\_stencrypt\_t** `stencrypt`  
Like ‘encrypt’, for stream ciphers. See below for a description of this type.
- gcry\_cipher\_stdecrypt\_t** `stdecrypt`  
Like ‘decrypt’, for stream ciphers. See below for a description of this type.
- gcry\_cipher\_oid\_spec\_t** [Data type]  
This type is used for associating a user-provided algorithm implementation with certain OIDs. It contains the following members:
- const char \***`oid`  
Textual representation of the OID.
- int** `mode` Cipher mode for which this OID is valid.
- gcry\_cipher\_setkey\_t** [Data type]  
Type for the ‘setkey’ function, defined as: `gcry_err_code_t (*gcry_cipher_setkey_t)(void *c, const unsigned char *key, unsigned keylen)`
- gcry\_cipher\_encrypt\_t** [Data type]  
Type for the ‘encrypt’ function, defined as: `gcry_err_code_t (*gcry_cipher_encrypt_t)(void *c, const unsigned char *outbuf, const unsigned char *inbuf)`
- gcry\_cipher\_decrypt\_t** [Data type]  
Type for the ‘decrypt’ function, defined as: `gcry_err_code_t (*gcry_cipher_decrypt_t)(void *c, const unsigned char *outbuf, const unsigned char *inbuf)`
- gcry\_cipher\_stencrypt\_t** [Data type]  
Type for the ‘stencrypt’ function, defined as: `gcry_err_code_t (*gcry_cipher_stencrypt_t)(void *c, const unsigned char *outbuf, const unsigned char *, unsigned int n)`

- gcry\_cipher\_stdecrypt\_t** [Data type]  
 Type for the ‘stdecrypt’ function, defined as: `gcry_err_code_t (*gcry_cipher_stdecrypt_t)(void *c, const unsigned char *outbuf, const unsigned char *, unsigned int n)`
- gcry\_error\_t gcry\_cipher\_register** (*gcry\_cipher\_spec\_t \*cipher*, [Function]  
*unsigned int \*algorithm\_id*, *gcry\_module\_t \*module*)  
 Register a new cipher module whose specification can be found in *cipher*. On success, a new algorithm ID is stored in *algorithm\_id* and a pointer representing this module is stored in *module*. Deprecated; the module register interface will be removed in a future version.
- void gcry\_cipher\_unregister** (*gcry\_module\_t module*) [Function]  
 Unregister the cipher identified by *module*, which must have been registered with `gcry_cipher_register`.
- gcry\_error\_t gcry\_cipher\_list** (*int \*list*, *int \*list\_length*) [Function]  
 Get a list consisting of the IDs of the loaded cipher modules. If *list* is zero, write the number of loaded cipher modules to *list\_length* and return. If *list* is non-zero, the first *\*list\_length* algorithm IDs are stored in *list*, which must be of according size. In case there are less cipher modules than *\*list\_length*, *\*list\_length* is updated to the correct number.

### 5.3 Available cipher modes

#### GCRY\_CIPHER\_MODE\_NONE

No mode specified. This should not be used. The only exception is that if Libcrypt is not used in FIPS mode and if any debug flag has been set, this mode may be used to bypass the actual encryption.

#### GCRY\_CIPHER\_MODE\_ECB

Electronic Codebook mode.

#### GCRY\_CIPHER\_MODE\_CFB

Cipher Feedback mode. The shift size equals the block size of the cipher (e.g. for AES it is CFB-128).

#### GCRY\_CIPHER\_MODE\_CBC

Cipher Block Chaining mode.

#### GCRY\_CIPHER\_MODE\_STREAM

Stream mode, only to be used with stream cipher algorithms.

#### GCRY\_CIPHER\_MODE\_OFB

Output Feedback mode.

#### GCRY\_CIPHER\_MODE\_CTR

Counter mode.

#### GCRY\_CIPHER\_MODE\_AESWRAP

This mode is used to implement the AES-Wrap algorithm according to RFC-3394. It may be used with any 128 bit block length algorithm, however the

specs require one of the 3 AES algorithms. These special conditions apply: If `gcry_cipher_setiv` has not been used the standard IV is used; if it has been used the lower 64 bit of the IV are used as the Alternative Initial Value. On encryption the provided output buffer must be 64 bit (8 byte) larger than the input buffer; in-place encryption is still allowed. On decryption the output buffer may be specified 64 bit (8 byte) shorter than then input buffer. As per specs the input length must be at least 128 bits and the length must be a multiple of 64 bits.

## 5.4 Working with cipher handles

To use a cipher algorithm, you must first allocate an according handle. This is to be done using the open function:

```
gcry_error_t gcry_cipher_open (gcry_cipher_hd_t *hd, int algo, int mode, unsigned int flags) [Function]
```

This function creates the context handle required for most of the other cipher functions and returns a handle to it in 'hd'. In case of an error, an according error code is returned.

The ID of algorithm to use must be specified via *algo*. See See [Section 5.1 \[Available ciphers\]](#), [page 23](#), for a list of supported ciphers and the according constants.

Besides using the constants directly, the function `gcry_cipher_map_name` may be used to convert the textual name of an algorithm into the according numeric ID.

The cipher mode to use must be specified via *mode*. See See [Section 5.3 \[Available cipher modes\]](#), [page 26](#), for a list of supported cipher modes and the according constants. Note that some modes are incompatible with some algorithms - in particular, stream mode (`GCRY_CIPHER_MODE_STREAM`) only works with stream ciphers. Any block cipher mode (`GCRY_CIPHER_MODE_ECB`, `GCRY_CIPHER_MODE_CBC`, `GCRY_CIPHER_MODE_CFB`, `GCRY_CIPHER_MODE_OFB` or `GCRY_CIPHER_MODE_CTR`) will work with any block cipher algorithm.

The third argument *flags* can either be passed as 0 or as the bit-wise OR of the following constants.

### `GCRY_CIPHER_SECURE`

Make sure that all operations are allocated in secure memory. This is useful when the key material is highly confidential.

### `GCRY_CIPHER_ENABLE_SYNC`

This flag enables the CFB sync mode, which is a special feature of Libgcrypt's CFB mode implementation to allow for OpenPGP's CFB variant. See `gcry_cipher_sync`.

### `GCRY_CIPHER_CBC_CTS`

Enable cipher text stealing (CTS) for the CBC mode. Cannot be used simultaneous as `GCRY_CIPHER_CBC_MAC`. CTS mode makes it possible to transform data of almost arbitrary size (only limitation is that it must be greater than the algorithm's block size).

**GCRY\_CIPHER\_CBC\_MAC**

Compute CBC-MAC keyed checksums. This is the same as CBC mode, but only output the last block. Cannot be used simultaneous as GCRY\_CIPHER\_CBC\_CTS.

Use the following function to release an existing handle:

**void** `gcry_cipher_close` (*gcry\_cipher\_hd\_t* **h**) [Function]

This function releases the context created by `gcry_cipher_open`. It also zeroes all sensitive information associated with this cipher handle.

In order to use a handle for performing cryptographic operations, a ‘key’ has to be set first:

**gcry\_error\_t** `gcry_cipher_setkey` (*gcry\_cipher\_hd\_t* **h**, *const void* **\*k**, *size\_t* **l**) [Function]

Set the key *k* used for encryption or decryption in the context denoted by the handle *h*. The length *l* (in bytes) of the key *k* must match the required length of the algorithm set for this context or be in the allowed range for algorithms with variable key size. The function checks this and returns an error if there is a problem. A caller should always check for an error.

Most crypto modes requires an initialization vector (IV), which usually is a non-secret random string acting as a kind of salt value. The CTR mode requires a counter, which is also similar to a salt value. To set the IV or CTR, use these functions:

**gcry\_error\_t** `gcry_cipher_setiv` (*gcry\_cipher\_hd\_t* **h**, *const void* **\*k**, *size\_t* **l**) [Function]

Set the initialization vector used for encryption or decryption. The vector is passed as the buffer *K* of length *l* bytes and copied to internal data structures. The function checks that the IV matches the requirement of the selected algorithm and mode.

**gcry\_error\_t** `gcry_cipher_setctr` (*gcry\_cipher\_hd\_t* **h**, *const void* **\*c**, *size\_t* **l**) [Function]

Set the counter vector used for encryption or decryption. The counter is passed as the buffer *c* of length *l* bytes and copied to internal data structures. The function checks that the counter matches the requirement of the selected algorithm (i.e., it must be the same size as the block size).

**gcry\_error\_t** `gcry_cipher_reset` (*gcry\_cipher\_hd\_t* **h**) [Function]

Set the given handle’s context back to the state it had after the last call to `gcry_cipher_setkey` and clear the initialization vector.

Note that `gcry_cipher_reset` is implemented as a macro.

The actual encryption and decryption is done by using one of the following functions. They may be used as often as required to process all the data.

**gcry\_error\_t** `gcry_cipher_encrypt` (*gcry\_cipher\_hd\_t* **h**, *unsigned char* **\*out**, *size\_t* **outside**, *const unsigned char* **\*in**, *size\_t* **inlen**) [Function]

`gcry_cipher_encrypt` is used to encrypt the data. This function can either work in place or with two buffers. It uses the cipher context already setup and described by

the handle *h*. There are 2 ways to use the function: If *in* is passed as `NULL` and *inlen* is 0, in-place encryption of the data in *out* or length *outsize* takes place. With *in* being not `NULL`, *inlen* bytes are encrypted to the buffer *out* which must have at least a size of *inlen*. *outsize* must be set to the allocated size of *out*, so that the function can check that there is sufficient space. Note that overlapping buffers are not allowed.

Depending on the selected algorithms and encryption mode, the length of the buffers must be a multiple of the block size.

The function returns 0 on success or an error code.

```
gcry_error_t gcry_cipher_decrypt (gcry_cipher_hd_t h, unsigned [Function]
 char *out, size_t outsize, const unsigned char *in, size_t inlen)
```

`gcry_cipher_decrypt` is used to decrypt the data. This function can either work in place or with two buffers. It uses the cipher context already setup and described by the handle *h*. There are 2 ways to use the function: If *in* is passed as `NULL` and *inlen* is 0, in-place decryption of the data in *out* or length *outsize* takes place. With *in* being not `NULL`, *inlen* bytes are decrypted to the buffer *out* which must have at least a size of *inlen*. *outsize* must be set to the allocated size of *out*, so that the function can check that there is sufficient space. Note that overlapping buffers are not allowed.

Depending on the selected algorithms and encryption mode, the length of the buffers must be a multiple of the block size.

The function returns 0 on success or an error code.

OpenPGP (as defined in RFC-2440) requires a special sync operation in some places. The following function is used for this:

```
gcry_error_t gcry_cipher_sync (gcry_cipher_hd_t h) [Function]
 Perform the OpenPGP sync operation on context h. Note that this is a no-op unless
 the context was created with the flag GCRY_CIPHER_ENABLE_SYNC
```

Some of the described functions are implemented as macros utilizing a catch-all control function. This control function is rarely used directly but there is nothing which would inhibit it:

```
gcry_error_t gcry_cipher_ctl (gcry_cipher_hd_t h, int cmd, void [Function]
 *buffer, size_t buflen)
```

`gcry_cipher_ctl` controls various aspects of the cipher module and specific cipher contexts. Usually some more specialized functions or macros are used for this purpose. The semantics of the function and its parameters depends on the the command *cmd* and the passed context handle *h*. Please see the comments in the source code (`src/global.c`) for details.

```
gcry_error_t gcry_cipher_info (gcry_cipher_hd_t h, int what, void [Function]
 *buffer, size_t *nbytes)
```

`gcry_cipher_info` is used to retrieve various information about a cipher context or the cipher module in general.

Currently no information is available.

## 5.5 General cipher functions

To work with the algorithms, several functions are available to map algorithm names to the internal identifiers, as well as ways to retrieve information about an algorithm or the current cipher context.

`gcry_error_t gcry_cipher_algo_info (int algo, int what, void [Function]  
*buffer, size_t *nbytes)`

This function is used to retrieve information on a specific algorithm. You pass the cipher algorithm ID as *algo* and the type of information requested as *what*. The result is either returned as the return code of the function or copied to the provided *buffer* whose allocated length must be available in an integer variable with the address passed in *nbytes*. This variable will also receive the actual used length of the buffer.

Here is a list of supported codes for *what*:

**GCRYCTL\_GET\_KEYLEN:**

Return the length of the key. If the algorithm supports multiple key lengths, the maximum supported value is returned. The length is returned as number of octets (bytes) and not as number of bits in *nbytes*; *buffer* must be zero. Note that it is usually better to use the convenience function `gcry_cipher_get_algo_keylen`.

**GCRYCTL\_GET\_BLKLEN:**

Return the block length of the algorithm. The length is returned as a number of octets in *nbytes*; *buffer* must be zero. Note that it is usually better to use the convenience function `gcry_cipher_get_algo_blklen`.

**GCRYCTL\_TEST\_ALGO:**

Returns 0 when the specified algorithm is available for use. *buffer* and *nbytes* must be zero.

`size_t gcry_cipher_get_algo_keylen (algo) [Function]`

This function returns length of the key for algorithm *algo*. If the algorithm supports multiple key lengths, the maximum supported key length is returned. On error 0 is returned. The key length is returned as number of octets.

This is a convenience functions which should be preferred over `gcry_cipher_algo_info` because it allows for proper type checking.

`size_t gcry_cipher_get_algo_blklen (int algo) [Function]`

This functions returns the blocklength of the algorithm *algo* counted in octets. On error 0 is returned.

This is a convenience functions which should be preferred over `gcry_cipher_algo_info` because it allows for proper type checking.

`const char * gcry_cipher_algo_name (int algo) [Function]`

`gcry_cipher_algo_name` returns a string with the name of the cipher algorithm *algo*. If the algorithm is not known or another error occurred, the string "?" is returned. This function should not be used to test for the availability of an algorithm.

`int gcry_cipher_map_name (const char *name)` [Function]  
gcry\_cipher\_map\_name returns the algorithm identifier for the cipher algorithm described by the string *name*. If this algorithm is not available 0 is returned.

`int gcry_cipher_mode_from_oid (const char *string)` [Function]  
Return the cipher mode associated with an ASN.1 object identifier. The object identifier is expected to be in the IETF-style dotted decimal notation. The function returns 0 for an unknown object identifier or when no mode is associated with it.





## 6 Public Key cryptography

Public key cryptography, also known as asymmetric cryptography, is an easy way for key management and to provide digital signatures. Libcrypt provides two completely different interfaces to public key cryptography, this chapter explains the one based on S-expressions.

### 6.1 Available algorithms

Libcrypt supports the RSA (Rivest-Shamir-Adleman) algorithms as well as DSA (Digital Signature Algorithm) and Elgamal. The versatile interface allows to add more algorithms in the future.

### 6.2 Used S-expressions

Libcrypt's API for asymmetric cryptography is based on data structures called S-expressions (see <http://people.csail.mit.edu/rivest/sexp.html>) and does not work with contexts as most of the other building blocks of Libcrypt do.

The following information are stored in S-expressions:

- keys
- plain text data
- encrypted data
- signatures

To describe how Libcrypt expect keys, we use examples. Note that words in italics indicate parameters whereas lowercase words are literals.

Note that all MPI (multi-precision-integers) values are expected to be in GCRYMPI\_FMT\_USG format. An easy way to create S-expressions is by using `gcry_sexp_build` which allows to pass a string with printf-like escapes to insert MPI values.

#### 6.2.1 RSA key parameters

An RSA private key is described by this S-expression:

```
(private-key
 (rsa
 (n n-mpi)
 (e e-mpi)
 (d d-mpi)
 (p p-mpi)
 (q q-mpi)
 (u u-mpi)))
```

An RSA public key is described by this S-expression:

```
(public-key
 (rsa
 (n n-mpi)
 (e e-mpi)))
```

*n-mpi*      RSA public modulus *n*.

|              |                                                     |
|--------------|-----------------------------------------------------|
| <i>e-mpi</i> | RSA public exponent $e$ .                           |
| <i>d-mpi</i> | RSA secret exponent $d = e^{-1} \bmod (p-1)(q-1)$ . |
| <i>p-mpi</i> | RSA secret prime $p$ .                              |
| <i>q-mpi</i> | RSA secret prime $q$ with $p < q$ .                 |
| <i>u-mpi</i> | Multiplicative inverse $u = p^{-1} \bmod q$ .       |

For signing and decryption the parameters  $(p, q, u)$  are optional but greatly improve the performance. Either all of these optional parameters must be given or none of them. They are mandatory for `gcry_pk_testkey`.

Note that OpenSSL uses slightly different parameters:  $q < p$  and  $u = q^{-1} \bmod p$ . To use these parameters you will need to swap the values and recompute  $u$ . Here is example code to do this:

```
if (gcry_mpi_cmp (p, q) > 0)
{
 gcry_mpi_swap (p, q);
 gcry_mpi_invm (u, p, q);
}
```

### 6.2.2 DSA key parameters

A DSA private key is described by this S-expression:

```
(private-key
 (dsa
 (p p-mpi)
 (q q-mpi)
 (g g-mpi)
 (y y-mpi)
 (x x-mpi)))
```

|              |                                                           |
|--------------|-----------------------------------------------------------|
| <i>p-mpi</i> | DSA prime $p$ .                                           |
| <i>q-mpi</i> | DSA group order $q$ (which is a prime divisor of $p-1$ ). |
| <i>g-mpi</i> | DSA group generator $g$ .                                 |
| <i>y-mpi</i> | DSA public key value $y = g^x \bmod p$ .                  |
| <i>x-mpi</i> | DSA secret exponent $x$ .                                 |

The public key is similar with "private-key" replaced by "public-key" and no *x-mpi*.

### 6.2.3 ECC key parameters

An ECC private key is described by this S-expression:

```
(private-key
 (ecc
 (p p-mpi)
 (a a-mpi)
 (b b-mpi)
 (g g-point))
```

```

 (n n-mpi)
 (q q-point)
 (d d-mpi))

```

*p-mpi* Prime specifying the field  $GF(p)$ .

*a-mpi*

*b-mpi* The two coefficients of the Weierstrass equation  $y^2 = x^3 + ax + b$

*g-point* Base point  $g$ .

*n-mpi* Order of  $g$

*q-point* The point representing the public key  $Q = dP$ .

*d-mpi* The private key  $d$

All point values are encoded in standard format; Libgcrypt does currently only support uncompressed points, thus the first byte needs to be 0x04.

The public key is similar with "private-key" replaced by "public-key" and no *d-mpi*.

If the domain parameters are well-known, the name of this curve may be used. For example

```

(private-key
 (ecc
 (curve "NIST P-192")
 (q q-point)
 (d d-mpi)))

```

The curve parameter may be given in any case and is used to replace missing parameters.

Currently implemented curves are:

NIST P-192

1.2.840.10045.3.1.1

prime192v1

secp192r1

The NIST 192 bit curve, its OID, X9.62 and SECP aliases.

NIST P-224

secp224r1

The NIST 224 bit curve and its SECP alias.

NIST P-256

1.2.840.10045.3.1.7

prime256v1

secp256r1

The NIST 256 bit curve, its OID, X9.62 and SECP aliases.

NIST P-384

secp384r1

The NIST 384 bit curve and its SECP alias.

NIST P-521

secp521r1

The NIST 521 bit curve and its SECP alias.

As usual the OIDs may optionally be prefixed with the string `OID.` or `oid..`

### 6.3 Public key modules

Libcrypt makes it possible to load additional ‘public key modules’; these public key algorithms can be used just like the algorithms that are built into the library directly. For an introduction into extension modules, see See [Section 3.2 \[Modules\]](#), page 13.

`gcry_pk_spec_t` [Data type]

This is the ‘module specification structure’ needed for registering public key modules, which has to be filled in by the user before it can be used to register a module. It contains the following members:

`const char *name`  
The primary name of this algorithm.

`char **aliases`  
A list of strings that are ‘aliases’ for the algorithm. The list must be terminated with a NULL element.

`const char *elements_pkey`  
String containing the one-letter names of the MPI values contained in a public key.

`const char *element_skey`  
String containing the one-letter names of the MPI values contained in a secret key.

`const char *elements_enc`  
String containing the one-letter names of the MPI values that are the result of an encryption operation using this algorithm.

`const char *elements_sig`  
String containing the one-letter names of the MPI values that are the result of a sign operation using this algorithm.

`const char *elements_grip`  
String containing the one-letter names of the MPI values that are to be included in the ‘key grip’.

`int use` The bitwise-OR of the following flags, depending on the abilities of the algorithm:

`GCRY_PK_USAGE_SIGN`  
The algorithm supports signing and verifying of data.

`GCRY_PK_USAGE_ENCR`  
The algorithm supports the encryption and decryption of data.

`gcry_pk_generate_t generate`  
The function responsible for generating a new key pair. See below for a description of this type.

`gcry_pk_check_secret_key_t check_secret_key`  
The function responsible for checking the sanity of a provided secret key. See below for a description of this type.

`gcry_pk_encrypt_t encrypt`  
 The function responsible for encrypting data. See below for a description of this type.

`gcry_pk_decrypt_t decrypt`  
 The function responsible for decrypting data. See below for a description of this type.

`gcry_pk_sign_t sign`  
 The function responsible for signing data. See below for a description of this type.

`gcry_pk_verify_t verify`  
 The function responsible for verifying that the provided signature matches the provided data. See below for a description of this type.

`gcry_pk_get_nbits_t get_nbits`  
 The function responsible for returning the number of bits of a provided key. See below for a description of this type.

`gcry_pk_generate_t` [Data type]  
 Type for the ‘generate’ function, defined as: `gcry_err_code_t (*gcry_pk_generate_t) (int algo, unsigned int nbits, unsigned long use_e, gcry_mpi_t *skey, gcry_mpi_t **ret-factors)`

`gcry_pk_check_secret_key_t` [Data type]  
 Type for the ‘check\_secret\_key’ function, defined as: `gcry_err_code_t (*gcry_pk_check_secret_key_t) (int algo, gcry_mpi_t *skey)`

`gcry_pk_encrypt_t` [Data type]  
 Type for the ‘encrypt’ function, defined as: `gcry_err_code_t (*gcry_pk_encrypt_t) (int algo, gcry_mpi_t *resarr, gcry_mpi_t data, gcry_mpi_t *pkey, int flags)`

`gcry_pk_decrypt_t` [Data type]  
 Type for the ‘decrypt’ function, defined as: `gcry_err_code_t (*gcry_pk_decrypt_t) (int algo, gcry_mpi_t *result, gcry_mpi_t *data, gcry_mpi_t *skey, int flags)`

`gcry_pk_sign_t` [Data type]  
 Type for the ‘sign’ function, defined as: `gcry_err_code_t (*gcry_pk_sign_t) (int algo, gcry_mpi_t *resarr, gcry_mpi_t data, gcry_mpi_t *skey)`

`gcry_pk_verify_t` [Data type]  
 Type for the ‘verify’ function, defined as: `gcry_err_code_t (*gcry_pk_verify_t) (int algo, gcry_mpi_t hash, gcry_mpi_t *data, gcry_mpi_t *pkey, int (*cmp) (void *, gcry_mpi_t), void *opaquev)`

`gcry_pk_get_nbits_t` [Data type]  
 Type for the ‘get\_nbits’ function, defined as: `unsigned (*gcry_pk_get_nbits_t) (int algo, gcry_mpi_t *pkey)`

`gcry_error_t gcry_pk_register (gcry_pk_spec_t *pubkey, unsigned int *algorithm_id, gcry_module_t *module)` [Function]

Register a new public key module whose specification can be found in *pubkey*. On success, a new algorithm ID is stored in *algorithm\_id* and a pointer representing this module is stored in *module*. Deprecated; the module register interface will be removed in a future version.

`void gcry_pk_unregister (gcry_module_t module)` [Function]

Unregister the public key module identified by *module*, which must have been registered with `gcry_pk_register`.

`gcry_error_t gcry_pk_list (int *list, int *list_length)` [Function]

Get a list consisting of the IDs of the loaded pubkey modules. If *list* is zero, write the number of loaded pubkey modules to *list\_length* and return. If *list* is non-zero, the first *\*list\_length* algorithm IDs are stored in *list*, which must be of according size. In case there are less pubkey modules than *\*list\_length*, *\*list\_length* is updated to the correct number.

## 6.4 Cryptographic Functions

Note that we will in future allow to use keys without p,q and u specified and may also support other parameters for performance reasons.

Some functions operating on S-expressions support ‘flags’, that influence the operation. These flags have to be listed in a sub-S-expression named ‘flags’; the following flags are known:

**pkcs1** Use PKCS#1 block type 2 padding for encryption, block type 1 padding for signing.

**oaep** Use RSA-OAEP padding for encryption.

**pss** Use RSA-PSS padding for signing.

**no-blinding**

Do not use a technique called ‘blinding’, which is used by default in order to prevent leaking of secret information. Blinding is only implemented by RSA, but it might be implemented by other algorithms in the future as well, when necessary.

Now that we know the key basics, we can carry on and explain how to encrypt and decrypt data. In almost all cases the data is a random session key which is in turn used for the actual encryption of the real data. There are 2 functions to do this:

`gcry_error_t gcry_pk_encrypt (gcry_sexp_t *r_ciph, gcry_sexp_t data, gcry_sexp_t pkey)` [Function]

Obviously a public key must be provided for encryption. It is expected as an appropriate S-expression (see above) in *pkey*. The data to be encrypted can either be in the simple old format, which is a very simple S-expression consisting only of one MPI, or it may be a more complex S-expression which also allows to specify flags for operation, like e.g. padding rules.

If you don't want to let Libcrypt handle the padding, you must pass an appropriate MPI using this expression for *data*:

```
(data
 (flags raw)
 (value mpi))
```

This has the same semantics as the old style MPI only way. *MPI* is the actual data, already padded appropriate for your protocol. Most RSA based systems however use PKCS#1 padding and so you can use this S-expression for *data*:

```
(data
 (flags pkcs1)
 (value block))
```

Here, the "flags" list has the "pkcs1" flag which let the function know that it should provide PKCS#1 block type 2 padding. The actual data to be encrypted is passed as a string of octets in *block*. The function checks that this data actually can be used with the given key, does the padding and encrypts it.

If the function could successfully perform the encryption, the return value will be 0 and a new S-expression with the encrypted result is allocated and assigned to the variable at the address of *r\_ciph*. The caller is responsible to release this value using `gcry_sexp_release`. In case of an error, an error code is returned and *r\_ciph* will be set to NULL.

The returned S-expression has this format when used with RSA:

```
(enc-val
 (rsa
 (a a-mpi)))
```

Where *a-mpi* is an MPI with the result of the RSA operation. When using the Elgamal algorithm, the return value will have this format:

```
(enc-val
 (elg
 (a a-mpi)
 (b b-mpi)))
```

Where *a-mpi* and *b-mpi* are MPIs with the result of the Elgamal encryption operation.

`gcry_error_t gcry_pk_decrypt (gcry_sexp_t *r_plain, [Function]  
gcry_sexp_t data, gcry_sexp_t skey)`

Obviously a private key must be provided for decryption. It is expected as an appropriate S-expression (see above) in *skey*. The data to be decrypted must match the format of the result as returned by `gcry_pk_encrypt`, but should be enlarged with a `flags` element:

```
(enc-val
 (flags)
 (elg
 (a a-mpi)
 (b b-mpi)))
```

This function does not remove padding from the data by default. To let Libgcrypt remove padding, give a hint in ‘flags’ telling which padding method was used when encrypting:

```
(flags padding-method)
```

Currently *padding-method* is either `pkcs1` for PKCS#1 block type 2 padding, or `oaep` for RSA-OAEP padding.

The function returns 0 on success or an error code. The variable at the address of *r\_plain* will be set to NULL on error or receive the decrypted value on success. The format of *r\_plain* is a simple S-expression part (i.e. not a valid one) with just one MPI if there was no `flags` element in *data*; if at least an empty `flags` is passed in *data*, the format is:

```
(value plaintext)
```

Another operation commonly performed using public key cryptography is signing data. In some sense this is even more important than encryption because digital signatures are an important instrument for key management. Libgcrypt supports digital signatures using 2 functions, similar to the encryption functions:

```
gcry_error_t gcry_pk_sign (gcry_sexp_t *r_sig, gcry_sexp_t data, [Function]
gcry_sexp_t skey)
```

This function creates a digital signature for *data* using the private key *skey* and place it into the variable at the address of *r\_sig*. *data* may either be the simple old style S-expression with just one MPI or a modern and more versatile S-expression which allows to let Libgcrypt handle padding:

```
(data
 (flags pkcs1)
 (hash hash-algo block))
```

This example requests to sign the data in *block* after applying PKCS#1 block type 1 style padding. *hash-algo* is a string with the hash algorithm to be encoded into the signature, this may be any hash algorithm name as supported by Libgcrypt. Most likely, this will be "sha256" or "sha1". It is obvious that the length of *block* must match the size of that message digests; the function checks that this and other constraints are valid.

If PKCS#1 padding is not required (because the caller does already provide a padded value), either the old format or better the following format should be used:

```
(data
 (flags raw)
 (value mpi))
```

Here, the data to be signed is directly given as an *MPI*.

The signature is returned as a newly allocated S-expression in *r\_sig* using this format for RSA:

```
(sig-val
 (rsa
 (s s-mpi)))
```

Where *s-mpi* is the result of the RSA sign operation. For DSA the S-expression returned is:



```
(sig-val
 (dsa
 (r r-mpi)
 (s s-mpi)))
```

Where *r-mpi* and *s-mpi* are the result of the DSA sign operation. For Elgamal signing (which is slow, yields large numbers and probably is not as secure as the other algorithms), the same format is used with "elg" replacing "dsa".

The operation most commonly used is definitely the verification of a signature. Libgcrypt provides this function:

```
gcry_error_t gcry_pk_verify (gcry_sexp_t sig, gcry_sexp_t data, [Function]
 gcry_sexp_t pkey)
```

This is used to check whether the signature *sig* matches the *data*. The public key *pkey* must be provided to perform this verification. This function is similar in its parameters to `gcry_pk_sign` with the exceptions that the public key is used instead of the private key and that no signature is created but a signature, in a format as created by `gcry_pk_sign`, is passed to the function in *sig*.

The result is 0 for success (i.e. the data matches the signature), or an error code where the most relevant code is `GCRYERR_BAD_SIGNATURE` to indicate that the signature does not match the provided data.

## 6.5 General public-key related Functions

A couple of utility functions are available to retrieve the length of the key, map algorithm identifiers and perform sanity checks:

```
const char * gcry_pk_algo_name (int algo) [Function]
```

Map the public key algorithm id *algo* to a string representation of the algorithm name. For unknown algorithms this function returns the string "?". This function should not be used to test for the availability of an algorithm.

```
int gcry_pk_map_name (const char *name) [Function]
```

Map the algorithm *name* to a public key algorithm Id. Returns 0 if the algorithm name is not known.

```
int gcry_pk_test_algo (int algo) [Function]
```

Return 0 if the public key algorithm *algo* is available for use. Note that this is implemented as a macro.

```
unsigned int gcry_pk_get_nbits (gcry_sexp_t key) [Function]
```

Return what is commonly referred as the key length for the given public or private in *key*.

```
unsigned char * gcry_pk_get_keygrip (gcry_sexp_t key, [Function]
 unsigned char *array)
```

Return the so called "keygrip" which is the SHA-1 hash of the public key parameters expressed in a way depended on the algorithm. *array* must either provide space for 20 bytes or be NULL. In the latter case a newly allocated array of that size is returned.

On success a pointer to the newly allocated space or to *array* is returned. `NULL` is returned to indicate an error which is most likely an unknown algorithm or one where a "keygrip" has not yet been defined. The function accepts public or secret keys in *key*.

`gcry_error_t gcry_pk_testkey (gcry_sexp_t key)` [Function]

Return zero if the private key *key* is 'sane', an error code otherwise. Note that it is not possible to check the 'saneness' of a public key.

`gcry_error_t gcry_pk_algo_info (int algo, int what, void *buffer, size_t *nbytes)` [Function]

Depending on the value of *what* return various information about the public key algorithm with the id *algo*. Note that the function returns -1 on error and the actual error code must be retrieved using the function `gcry_errno`. The currently defined values for *what* are:

`GCRYCTL_TEST_ALGO:`

Return 0 if the specified algorithm is available for use. *buffer* must be `NULL`, *nbytes* may be passed as `NULL` or point to a variable with the required usage of the algorithm. This may be 0 for "don't care" or the bit-wise OR of these flags:

`GCRY_PK_USAGE_SIGN`

Algorithm is usable for signing.

`GCRY_PK_USAGE_ENCR`

Algorithm is usable for encryption.

Unless you need to test for the allowed usage, it is in general better to use the macro `gcry_pk_test_algo` instead.

`GCRYCTL_GET_ALGO_USAGE:`

Return the usage flags for the given algorithm. An invalid algorithm return 0. Disabled algorithms are ignored here because we want to know whether the algorithm is at all capable of a certain usage.

`GCRYCTL_GET_ALGO_NPKEY`

Return the number of elements the public key for algorithm *algo* consist of. Return 0 for an unknown algorithm.

`GCRYCTL_GET_ALGO_NSKEY`

Return the number of elements the private key for algorithm *algo* consist of. Note that this value is always larger than that of the public key. Return 0 for an unknown algorithm.

`GCRYCTL_GET_ALGO_NSIGN`

Return the number of elements a signature created with the algorithm *algo* consists of. Return 0 for an unknown algorithm or for an algorithm not capable of creating signatures.

`GCRYCTL_GET_ALGO_NENC`

Return the number of elements a encrypted message created with the algorithm *algo* consists of. Return 0 for an unknown algorithm or for an algorithm not capable of encryption.

Please note that parameters not required should be passed as NULL.

`gcry_error_t gcry_pk_ctl (int cmd, void *buffer, size_t buflen)` [Function]

This is a general purpose function to perform certain control operations. *cmd* controls what is to be done. The return value is 0 for success or an error code. Currently supported values for *cmd* are:

`GCRYCTL_DISABLE_ALGO`

Disable the algorithm given as an algorithm id in *buffer*. *buffer* must point to an `int` variable with the algorithm id and *buflen* must have the value `sizeof (int)`.

Libgcrypt also provides a function to generate public key pairs:

`gcry_error_t gcry_pk_genkey (gcry_sexp_t *r_key,` [Function]  
`gcry_sexp_t parms)`

This function create a new public key pair using information given in the S-expression *parms* and stores the private and the public key in one new S-expression at the address given by *r\_key*. In case of an error, *r\_key* is set to NULL. The return code is 0 for success or an error code otherwise.

Here is an example for *parms* to create an 2048 bit RSA key:

```
(genkey
 (rsa
 (nbits 4:2048)))
```

To create an Elgamal key, substitute "elg" for "rsa" and to create a DSA key use "dsa". Valid ranges for the key length depend on the algorithms; all commonly used key lengths are supported. Currently supported parameters are:

**nbits** This is always required to specify the length of the key. The argument is a string with a number in C-notation. The value should be a multiple of 8.

**curve name**

For ECC a named curve may be used instead of giving the number of requested bits. This allows to request a specific curve to override a default selection Libgcrypt would have taken if **nbits** has been given. The available names are listed with the description of the ECC public key parameters.

**rsa-use-e**

This is only used with RSA to give a hint for the public exponent. The value will be used as a base to test for a usable exponent. Some values are special:

|       |                                                                                      |
|-------|--------------------------------------------------------------------------------------|
| '0'   | Use a secure and fast value. This is currently the number 41.                        |
| '1'   | Use a value as required by some crypto policies. This is currently the number 65537. |
| '2'   | Reserved                                                                             |
| '> 2' | Use the given value.                                                                 |

If this parameter is not used, Libcrypt uses for historic reasons 65537.

**qbits** This is only meaningful for DSA keys. If it is given the DSA key is generated with a Q parameter of this size. If it is not given or zero Q is deduced from NBITS in this way:

```
'512 <= N <= 1024'
 Q = 160
'N = 2048' Q = 224
'N = 3072' Q = 256
'N = 7680' Q = 384
'N = 15360'
 Q = 512
```

Note that in this case only the values for N, as given in the table, are allowed. When specifying Q all values of N in the range 512 to 15680 are valid as long as they are multiples of 8.

**transient-key**

This is only meaningful for RSA, DSA, ECDSA, and ECDH keys. This is a flag with no value. If given the key is created using a faster and a somewhat less secure random number generator. This flag may be used for keys which are only used for a short time or per-message and do not require full cryptographic strength.

**domain**

This is only meaningful for DLP algorithms. If specified keys are generated with domain parameters taken from this list. The exact format of this parameter depends on the actual algorithm. It is currently only implemented for DSA using this format:

```
(genkey
 (dsa
 (domain
 (p p-mpi)
 (q q-mpi)
 (g q-mpi))))
```

**nbits** and **qbits** may not be specified because they are derived from the domain parameters.

**derive-parms**

This is currently only implemented for RSA and DSA keys. It is not allowed to use this together with a **domain** specification. If given, it is used to derive the keys using the given parameters.

If given for an RSA key the X9.31 key generation algorithm is used even if libcrypt is not in FIPS mode. If given for a DSA key, the FIPS 186 algorithm is used even if libcrypt is not in FIPS mode.

```
(genkey
 (rsa
 (nbits 4:1024)
```

```

(rsa-use-e 1:3)
(derive-parms
 (Xp1 #1A1916DDB29B4EB7EB6732E128#)
 (Xp2 #192E8AAC41C576C822D93EA433#)
 (Xp #D8CD81F035EC57EFE822955149D3BFF70C53520D
 769D6D76646C7A792E16EBD89FE6FC5B605A6493
 39DFC925A86A4C6D150B71B9EEA02D68885F5009
 B98BD984#)
 (Xq1 #1A5CF72EE770DE50CB09ACCEA9#)
 (Xq2 #134E4CAA16D2350A21D775C404#)
 (Xq #CC1092495D867E64065DEE3E7955F2EBC7D47A2D
 7C9953388F97DDDC3E1CA19C35CA659EDC2FC325
 6D29C2627479C086A699A49C4C9CEE7EF7BD1B34
 321DE34A#))))

(genkey
 (dsa
 (nbits 4:1024)
 (derive-parms
 (seed seed-mpi))))

```

**use-x931** Force the use of the ANSI X9.31 key generation algorithm instead of the default algorithm. This flag is only meaningful for RSA and usually not required. Note that this algorithm is implicitly used if either **derive-parms** is given or Libgcrypt is in FIPS mode.

**use-fips186**

Force the use of the FIPS 186 key generation algorithm instead of the default algorithm. This flag is only meaningful for DSA and usually not required. Note that this algorithm is implicitly used if either **derive-parms** is given or Libgcrypt is in FIPS mode. As of now FIPS 186-2 is implemented; after the approval of FIPS 186-3 the code will be changed to implement 186-3.

**use-fips186-2**

Force the use of the FIPS 186-2 key generation algorithm instead of the default algorithm. This algorithm is slightly different from FIPS 186-3 and allows only 1024 bit keys. This flag is only meaningful for DSA and only required for FIPS testing backward compatibility.

The key pair is returned in a format depending on the algorithm. Both private and public keys are returned in one container and may be accompanied by some miscellaneous information.

As an example, here is what the Elgamal key generation returns:

```

(key-data
 (public-key
 (elg
 (p p-mpi)
 (g g-mpi))

```

```

 (y y-mpi)))
(private-key
 (elg
 (p p-mpi)
 (g g-mpi)
 (y y-mpi)
 (x x-mpi)))
(misc-key-info
 (pm1-factors n1 n2 ... nn))

```

As you can see, some of the information is duplicated, but this provides an easy way to extract either the public or the private key. Note that the order of the elements is not defined, e.g. the private key may be stored before the public key. *n1 n2 ... nn* is a list of prime numbers used to composite *p-mpi*; this is in general not a very useful information and only available if the key generation algorithm provides them.

## 6.6 Alternative Public Key Interface

This section documents the alternative interface to asymmetric cryptography (*ac*) that is not based on S-expressions, but on native C data structures. As opposed to the *pk* interface described in the former chapter, this one follows an open/use/close paradigm like other building blocks of the library.

**This interface has a few known problems; most noteworthy an inherent tendency to leak memory. It might not be available in forthcoming versions of Libgcrypt.**

### 6.6.1 Available asymmetric algorithms

Libgcrypt supports the RSA (Rivest-Shamir-Adleman) algorithms as well as DSA (Digital Signature Algorithm) and Elgamal. The versatile interface allows to add more algorithms in the future.

`gcry_ac_id_t` [Data type]

The following constants are defined for this type:

```

GCRY_AC_RSA
 Rivest-Shamir-Adleman

GCRY_AC_DSA
 Digital Signature Algorithm

GCRY_AC_ELG
 Elgamal

GCRY_AC_ELG_E
 Elgamal, encryption only.

```

### 6.6.2 Working with sets of data

In the context of this interface the term ‘data set’ refers to a list of ‘named MPI values’ that is used by functions performing cryptographic operations; a named MPI value is an MPI value, associated with a label.

Such data sets are used for representing keys, since keys simply consist of a variable amount of numbers. Furthermore some functions return data sets to the caller that are to be provided to other functions.

This section documents the data types, symbols and functions that are relevant for working with data sets.

**gcry\_ac\_data\_t** [Data type]  
A single data set.

The following flags are supported:

**GCRY\_AC\_FLAG\_DEALLOC**

Used for storing data in a data set. If given, the data will be released by the library. Note that whenever one of the ac functions is about to release objects because of this flag, the objects are expected to be stored in memory allocated through the Libgcrypt memory management. In other words: `gcry_free()` is used instead of `free()`.

**GCRY\_AC\_FLAG\_COPY**

Used for storing/retrieving data in/from a data set. If given, the library will create copies of the provided/contained data, which will then be given to the user/associated with the data set.

**gcry\_error\_t gcry\_ac\_data\_new** (*gcry\_ac\_data\_t \*data*) [Function]  
Creates a new, empty data set and stores it in *data*.

**void gcry\_ac\_data\_destroy** (*gcry\_ac\_data\_t data*) [Function]  
Destroys the data set *data*.

**gcry\_error\_t gcry\_ac\_data\_set** (*gcry\_ac\_data\_t data, unsigned int flags, char \*name, gcry\_mpi\_t mpi*) [Function]  
Add the value *mpi* to *data* with the label *name*. If *flags* contains **GCRY\_AC\_FLAG\_COPY**, the data set will contain copies of *name* and *mpi*. If *flags* contains **GCRY\_AC\_FLAG\_DEALLOC** or **GCRY\_AC\_FLAG\_COPY**, the values contained in the data set will be deallocated when they are to be removed from the data set.

**gcry\_error\_t gcry\_ac\_data\_copy** (*gcry\_ac\_data\_t \*data\_cp, gcry\_ac\_data\_t data*) [Function]  
Create a copy of the data set *data* and store it in *data\_cp*. FIXME: exact semantics undefined.

**unsigned int gcry\_ac\_data\_length** (*gcry\_ac\_data\_t data*) [Function]  
Returns the number of named MPI values inside of the data set *data*.

**gcry\_error\_t gcry\_ac\_data\_get\_name** (*gcry\_ac\_data\_t data, unsigned int flags, char \*name, gcry\_mpi\_t \*mpi*) [Function]  
Store the value labelled with *name* found in *data* in *mpi*. If *flags* contains **GCRY\_AC\_FLAG\_COPY**, store a copy of the *mpi* value contained in the data set. *mpi* may be NULL (this might be useful for checking the existence of an MPI with extracting it).

`gcry_error_t gcry_ac_data_get_index` (`gcry_ac_data_t data`, [Function]  
*unsigned int flags, unsigned int index, const char \*\*name, gcry\_mpi\_t \*mpi*)

Stores in *name* and *mpi* the named *mpi* value contained in the data set *data* with the index *idx*. If *flags* contains `GCRY_AC_FLAG_COPY`, store copies of the values contained in the data set. *name* or *mpi* may be `NULL`.

`void gcry_ac_data_clear` (`gcry_ac_data_t data`) [Function]

Destroys any values contained in the data set *data*.

`gcry_error_t gcry_ac_data_to_sexp` (`gcry_ac_data_t data`, [Function]  
*gcry\_sexp\_t \*sexp, const char \*\*identifiers*)

This function converts the data set *data* into a newly created S-Expression, which is to be stored in *sexp*; *identifiers* is a `NULL` terminated list of C strings, which specifies the structure of the S-Expression.

Example:

If *identifiers* is a list of pointers to the strings “foo” and “bar” and if *data* is a data set containing the values “val1 = 0x01” and “val2 = 0x02”, then the resulting S-Expression will look like this: (foo (bar ((val1 0x01) (val2 0x02))).

`gcry_error_t gcry_ac_data_from_sexp` (`gcry_ac_data_t *data`, [Function]  
*gcry\_sexp\_t sexp, const char \*\*identifiers*)

This function converts the S-Expression *sexp* into a newly created data set, which is to be stored in *data*; *identifiers* is a `NULL` terminated list of C strings, which specifies the structure of the S-Expression. If the list of identifiers does not match the structure of the S-Expression, the function fails.

### 6.6.3 Working with IO objects

Note: IO objects are currently only used in the context of message encoding/decoding and encryption/signature schemes.

`gcry_ac_io_t` [Data type]

`gcry_ac_io_t` is the type to be used for IO objects.

IO objects provide an uniform IO layer on top of different underlying IO mechanisms; either they can be used for providing data to the library (mode is `GCRY_AC_IO_READABLE`) or they can be used for retrieving data from the library (mode is `GCRY_AC_IO_WRITABLE`).

IO object need to be initialized by calling on of the following functions:

`void gcry_ac_io_init` (`gcry_ac_io_t *ac_io, gcry_ac_io_mode_t mode`, [Function]  
*gcry\_ac\_io\_type\_t type, ...*);

Initialize *ac\_io* according to *mode*, *type* and the variable list of arguments. The list of variable arguments to specify depends on the given *type*.

`void gcry_ac_io_init_va` (`gcry_ac_io_t *ac_io, gcry_ac_io_mode_t` [Function]  
*mode, gcry\_ac\_io\_type\_t type, va\_list ap*);

Initialize *ac\_io* according to *mode*, *type* and the variable list of arguments *ap*. The list of variable arguments to specify depends on the given *type*.



The following types of IO objects exist:

#### GCRY\_AC\_IO\_STRING

In case of GCRY\_AC\_IO\_READABLE the IO object will provide data from a memory string. Arguments to specify at initialization time:

`unsigned char *`  
 Pointer to the beginning of the memory string

`size_t`     Size of the memory string

In case of GCRY\_AC\_IO\_WRITABLE the object will store retrieved data in a newly allocated memory string. Arguments to specify at initialization time:

`unsigned char **`  
 Pointer to address, at which the pointer to the newly created memory string is to be stored

`size_t *`    Pointer to address, at which the size of the newly created memory string is to be stored

#### GCRY\_AC\_IO\_CALLBACK

In case of GCRY\_AC\_IO\_READABLE the object will forward read requests to a provided callback function. Arguments to specify at initialization time:

`gcry_ac_data_read_cb_t`  
 Callback function to use

`void *`     Opaque argument to provide to the callback function

In case of GCRY\_AC\_IO\_WRITABLE the object will forward write requests to a provided callback function. Arguments to specify at initialization time:

`gcry_ac_data_write_cb_t`  
 Callback function to use

`void *`     Opaque argument to provide to the callback function

### 6.6.4 Working with handles

In order to use an algorithm, an according handle must be created. This is done using the following function:

`gcry_error_t gcry_ac_open (gcry_ac_handle_t *handle, int algorithm, int flags)` [Function]

Creates a new handle for the algorithm *algorithm* and stores it in *handle*. *flags* is not used currently.

*algorithm* must be a valid algorithm ID, see See [Section 6.6.1 \[Available asymmetric algorithms\]](#), page 46, for a list of supported algorithms and the according constants. Besides using the listed constants directly, the functions `gcry_pk_name_to_id` may be used to convert the textual name of an algorithm into the according numeric ID.

`void gcry_ac_close (gcry_ac_handle_t handle)` [Function]

Destroys the handle *handle*.

### 6.6.5 Working with keys

`gcry_ac_key_type_t` [Data type]

Defined constants:

`GCRY_AC_KEY_SECRET`

Specifies a secret key.

`GCRY_AC_KEY_PUBLIC`

Specifies a public key.

`gcry_ac_key_t` [Data type]

This type represents a single ‘key’, either a secret one or a public one.

`gcry_ac_key_pair_t` [Data type]

This type represents a ‘key pair’ containing a secret and a public key.

Key data structures can be created in two different ways; a new key pair can be generated, resulting in ready-to-use key. Alternatively a key can be initialized from a given data set.

`gcry_error_t gcry_ac_key_init (gcry_ac_key_t *key, [Function]  
gcry_ac_handle_t handle, gcry_ac_key_type_t type, gcry_ac_data_t data)`

Creates a new key of type *type*, consisting of the MPI values contained in the data set *data* and stores it in *key*.

`gcry_error_t gcry_ac_key_pair_generate (gcry_ac_handle_t [Function]  
handle, unsigned int nbits, void *key_spec, gcry_ac_key_pair_t *key_pair,  
gcry_mpi_t **misc_data)`

Generates a new key pair via the handle *handle* of *NBITS* bits and stores it in *key\_pair*.

In case non-standard settings are wanted, a pointer to a structure of type `gcry_ac_key_spec_<algorithm>_t`, matching the selected algorithm, can be given as *key\_spec*. *misc\_data* is not used yet. Such a structure does only exist for RSA. A description of the members of the supported structures follows.

`gcry_ac_key_spec_rsa_t`

`gcry_mpi_t e`

Generate the key pair using a special *e*. The value of *e* has the following meanings:

= 0      Let Libgcrypt decide what exponent should be used.

= 1      Request the use of a “secure” exponent; this is required by some specification to be 65537.

> 2      Try starting at this value until a working exponent is found. Note that the current implementation leaks some information about the private key because the incrementation used is not randomized. Thus, this function will be changed in the future to return a random exponent of the given size.

Example code:

```

{
 gcry_ac_key_pair_t key_pair;
 gcry_ac_key_spec_rsa_t rsa_spec;

 rsa_spec.e = gcry_mpi_new (0);
 gcry_mpi_set_ui (rsa_spec.e, 1);

 err = gcry_ac_open (&handle, GCRY_AC_RSA, 0);
 assert (! err);

 err = gcry_ac_key_pair_generate (handle, 1024, &rsa_spec,
 &key_pair, NULL);
 assert (! err);
}

```

`gcry_ac_key_t gcry_ac_key_pair_extract (gcry_ac_key_pair_t` [Function]

`key_pair, gcry_ac_key_type_t which)`

Returns the key of type *which* out of the key pair *key\_pair*.

`void gcry_ac_key_destroy (gcry_ac_key_t key)` [Function]

Destroys the key *key*.

`void gcry_ac_key_pair_destroy (gcry_ac_key_pair_t key_pair)` [Function]

Destroys the key pair *key\_pair*.

`gcry_ac_data_t gcry_ac_key_data_get (gcry_ac_key_t key)` [Function]

Returns the data set contained in the key *key*.

`gcry_error_t gcry_ac_key_test (gcry_ac_handle_t handle,` [Function]  
`gcry_ac_key_t key)`

Verifies that the private key *key* is sane via *handle*.

`gcry_error_t gcry_ac_key_get_nbits (gcry_ac_handle_t handle,` [Function]  
`gcry_ac_key_t key, unsigned int *nbits)`

Stores the number of bits of the key *key* in *nbits* via *handle*.

`gcry_error_t gcry_ac_key_get_grip (gcry_ac_handle_t handle,` [Function]  
`gcry_ac_key_t key, unsigned char *key_grip)`

Writes the 20 byte long key grip of the key *key* to *key\_grip* via *handle*.

### 6.6.6 Using cryptographic functions

The following flags might be relevant:

`GCRY_AC_FLAG_NO_BLINDING`

Disable any blinding, which might be supported by the chosen algorithm; blinding is the default.

There exist two kinds of cryptographic functions available through the ac interface: primitives, and high-level functions.

Primitives deal with MPIs (data sets) directly; what they provide is direct access to the cryptographic operations provided by an algorithm implementation.

High-level functions deal with octet strings, according to a specified “scheme”. Schemes make use of “encoding methods”, which are responsible for converting the provided octet strings into MPIs, which are then forwarded to the cryptographic primitives. Since schemes are to be used for a special purpose in order to achieve a particular security goal, there exist “encryption schemes” and “signature schemes”. Encoding methods can be used separately or implicitly through schemes.

What follows is a description of the cryptographic primitives.

`gcry_error_t gcry_ac_data_encrypt (gcry_ac_handle_t handle, [Function]  
                   unsigned int flags, gcry_ac_key_t key, gcry_mpi_t data_plain,  
                   gcry_ac_data_t *data_encrypted)`

Encrypts the plain text MPI value *data\_plain* with the key *key* under the control of the flags *flags* and stores the resulting data set into *data\_encrypted*.

`gcry_error_t gcry_ac_data_decrypt (gcry_ac_handle_t handle, [Function]  
                   unsigned int flags, gcry_ac_key_t key, gcry_mpi_t *data_plain,  
                   gcry_ac_data_t data_encrypted)`

Decrypts the encrypted data contained in the data set *data\_encrypted* with the secret key *key* under the control of the flags *flags* and stores the resulting plain text MPI value in *DATA\_PLAIN*.

`gcry_error_t gcry_ac_data_sign (gcry_ac_handle_t handle, [Function]  
                   gcry_ac_key_t key, gcry_mpi_t data, gcry_ac_data_t *data_signature)`

Signs the data contained in *data* with the secret key *key* and stores the resulting signature in the data set *data\_signature*.

`gcry_error_t gcry_ac_data_verify (gcry_ac_handle_t handle, [Function]  
                   gcry_ac_key_t key, gcry_mpi_t data, gcry_ac_data_t data_signature)`

Verifies that the signature contained in the data set *data\_signature* is indeed the result of signing the data contained in *data* with the secret key belonging to the public key *key*.

What follows is a description of the high-level functions.

The type “`gcry_ac_em_t`” is used for specifying encoding methods; the following methods are supported:

`GCRY_AC_EME_PKCS_V1_5`

PKCS-V1.5 Encoding Method for Encryption. Options must be provided through a pointer to a correctly initialized object of type `gcry_ac_eme_pkcs_v1_5_t`.

`GCRY_AC EMSA_PKCS_V1_5`

PKCS-V1.5 Encoding Method for Signatures with Appendix. Options must be provided through a pointer to a correctly initialized object of type `gcry_ac_emsa_pkcs_v1_5_t`.

Option structure types:

```

gcry_ac_eme_pkcs_v1_5_t
 gcry_ac_key_t key
 gcry_ac_handle_t handle
gcry_ac_emsa_pkcs_v1_5_t
 gcry_md_algo_t md
 size_t em_n

```

Encoding methods can be used directly through the following functions:

```

gcry_error_t gcry_ac_data_encode (gcry_ac_em_t method, unsigned [Function]
 int flags, void *options, unsigned char *m, size_t m_n, unsigned char **em,
 size_t *em_n)

```

Encodes the message contained in *m* of size *m\_n* according to *method*, *flags* and *options*. The newly created encoded message is stored in *em* and *em\_n*.

```

gcry_error_t gcry_ac_data_decode (gcry_ac_em_t method, unsigned [Function]
 int flags, void *options, unsigned char *em, size_t em_n, unsigned char **m,
 size_t *m_n)

```

Decodes the message contained in *em* of size *em\_n* according to *method*, *flags* and *options*. The newly created decoded message is stored in *m* and *m\_n*.

The type “*gcry\_ac\_scheme\_t*” is used for specifying schemes; the following schemes are supported:

**GCRY\_AC\_ES\_PKCS\_V1\_5**

PKCS-V1.5 Encryption Scheme. No options can be provided.

**GCRY\_AC\_SSA\_PKCS\_V1\_5**

PKCS-V1.5 Signature Scheme (with Appendix). Options can be provided through a pointer to a correctly initialized object of type *gcry\_ac\_ssa\_pkcs\_v1\_5\_t*.

Option structure types:

```

gcry_ac_ssa_pkcs_v1_5_t
 gcry_md_algo_t md

```

The functions implementing schemes:

```

gcry_error_t gcry_ac_data_encrypt_scheme (gcry_ac_handle_t [Function]
 handle, gcry_ac_scheme_t scheme, unsigned int flags, void *opts,
 gcry_ac_key_t key, gcry_ac_io_t *io_message, gcry_ac_io_t *io_cipher)

```

Encrypts the plain text readable from *io\_message* through *handle* with the public key *key* according to *scheme*, *flags* and *opts*. If *opts* is not NULL, it has to be a pointer to a structure specific to the chosen scheme (*gcry\_ac\_es\_\*\_t*). The encrypted message is written to *io\_cipher*.

```

gcry_error_t gcry_ac_data_decrypt_scheme (gcry_ac_handle_t [Function]
 handle, gcry_ac_scheme_t scheme, unsigned int flags, void *opts,
 gcry_ac_key_t key, gcry_ac_io_t *io_cipher, gcry_ac_io_t *io_message)

```

Decrypts the cipher text readable from *io\_cipher* through *handle* with the secret key *key* according to *scheme*, *flags* and *opts*. If *opts* is not NULL, it has to be a pointer

to a structure specific to the chosen scheme (`gcry_ac_es_*_t`). The decrypted message is written to `io_message`.

`gcry_error_t gcry_ac_data_sign_scheme` (`gcry_ac_handle_t` [Function]  
`handle`, `gcry_ac_scheme_t scheme`, `unsigned int flags`, `void *opts`,  
`gcry_ac_key_t key`, `gcry_ac_io_t *io_message`, `gcry_ac_io_t *io_signature`)  
 Signs the message readable from `io_message` through `handle` with the secret key `key` according to `scheme`, `flags` and `opts`. If `opts` is not NULL, it has to be a pointer to a structure specific to the chosen scheme (`gcry_ac_ssa_*_t`). The signature is written to `io_signature`.

`gcry_error_t gcry_ac_data_verify_scheme` (`gcry_ac_handle_t` [Function]  
`handle`, `gcry_ac_scheme_t scheme`, `unsigned int flags`, `void *opts`,  
`gcry_ac_key_t key`, `gcry_ac_io_t *io_message`, `gcry_ac_io_t *io_signature`)  
 Verifies through `handle` that the signature readable from `io_signature` is indeed the result of signing the message readable from `io_message` with the secret key belonging to the public key `key` according to `scheme` and `opts`. If `opts` is not NULL, it has to be an anonymous structure (`gcry_ac_ssa_*_t`) specific to the chosen scheme.

### 6.6.7 Handle-independent functions

These two functions are deprecated; do not use them for new code.

`gcry_error_t gcry_ac_id_to_name` (`gcry_ac_id_t algorithm`, `const` [Function]  
`char **name`)  
 Stores the textual representation of the algorithm whose id is given in `algorithm` in `name`. Deprecated; use `gcry_pk_algo_name`.

`gcry_error_t gcry_ac_name_to_id` (`const char *name`, `gcry_ac_id_t` [Function]  
`*algorithm`)  
 Stores the numeric ID of the algorithm whose textual representation is contained in `name` in `algorithm`. Deprecated; use `gcry_pk_map_name`.

## 7 Hashing

Libcrypt provides an easy and consistent to use interface for hashing. Hashing is buffered and several hash algorithms can be updated at once. It is possible to compute a MAC using the same routines. The programming model follows an open/process/close paradigm and is in that similar to other building blocks provided by Libcrypt.

For convenience reasons, a few cyclic redundancy check value operations are also supported.

### 7.1 Available hash algorithms

#### GCRY\_MD\_NONE

This is not a real algorithm but used by some functions as an error return value. This constant is guaranteed to have the value 0.

#### GCRY\_MD\_SHA1

This is the SHA-1 algorithm which yields a message digest of 20 bytes. Note that SHA-1 begins to show some weaknesses and it is suggested to fade out its use if strong cryptographic properties are required.

#### GCRY\_MD\_RMD160

This is the 160 bit version of the RIPE message digest (RIPE-MD-160). Like SHA-1 it also yields a digest of 20 bytes. This algorithm share a lot of design properties with SHA-1 and thus it is advisable not to use it for new protocols.

#### GCRY\_MD\_MD5

This is the well known MD5 algorithm, which yields a message digest of 16 bytes. Note that the MD5 algorithm has severe weaknesses, for example it is easy to compute two messages yielding the same hash (collision attack). The use of this algorithm is only justified for non-cryptographic application.

#### GCRY\_MD\_MD4

This is the MD4 algorithm, which yields a message digest of 16 bytes. This algorithms ha severe weaknesses and should not be used.

#### GCRY\_MD\_MD2

This is an reserved identifier for MD-2; there is no implementation yet. This algorithm has severe weaknesses and should not be used.

#### GCRY\_MD\_TIGER

This is the TIGER/192 algorithm which yields a message digest of 24 bytes. Actually this is a variant of TIGER with a different output print order as used by GnuPG up to version 1.3.2.

#### GCRY\_MD\_TIGER1

This is the TIGER variant as used by the NESSIE project. It uses the most commonly used output print order.

#### GCRY\_MD\_TIGER2

This is another variant of TIGER with a different padding scheme.

**GCRY\_MD\_HAVAL**

This is an reserved value for the HAVAL algorithm with 5 passes and 160 bit. It yields a message digest of 20 bytes. Note that there is no implementation yet available.

**GCRY\_MD\_SHA224**

This is the SHA-224 algorithm which yields a message digest of 28 bytes. See Change Notice 1 for FIPS 180-2 for the specification.

**GCRY\_MD\_SHA256**

This is the SHA-256 algorithm which yields a message digest of 32 bytes. See FIPS 180-2 for the specification.

**GCRY\_MD\_SHA384**

This is the SHA-384 algorithm which yields a message digest of 48 bytes. See FIPS 180-2 for the specification.

**GCRY\_MD\_SHA512**

This is the SHA-384 algorithm which yields a message digest of 64 bytes. See FIPS 180-2 for the specification.

**GCRY\_MD\_CRC32**

This is the ISO 3309 and ITU-T V.42 cyclic redundancy check. It yields an output of 4 bytes. Note that this is not a hash algorithm in the cryptographic sense.

**GCRY\_MD\_CRC32\_RFC1510**

This is the above cyclic redundancy check function, as modified by RFC 1510. It yields an output of 4 bytes. Note that this is not a hash algorithm in the cryptographic sense.

**GCRY\_MD\_CRC24\_RFC2440**

This is the OpenPGP cyclic redundancy check function. It yields an output of 3 bytes. Note that this is not a hash algorithm in the cryptographic sense.

**GCRY\_MD\_WHIRLPOOL**

This is the Whirlpool algorithm which yields a message digest of 64 bytes.

## 7.2 Hash algorithm modules

Libcrypt makes it possible to load additional ‘message digest modules’; these digests can be used just like the message digest algorithms that are built into the library directly. For an introduction into extension modules, see See [Section 3.2 \[Modules\]](#), page 13.

**gcry\_md\_spec\_t**

[Data type]

This is the ‘module specification structure’ needed for registering message digest modules, which has to be filled in by the user before it can be used to register a module. It contains the following members:

**const char \*name**

The primary name of this algorithm.

**unsigned char \*asnoid**

Array of bytes that form the ASN OID.



`int asnlen`  
Length of bytes in ‘asnoid’.

`gcry_md_oid_spec_t *oids`  
A list of OIDs that are to be associated with the algorithm. The list’s last element must have its ‘oid’ member set to NULL. See below for an explanation of this type. See below for an explanation of this type.

`int mdlen` Length of the message digest algorithm. See below for an explanation of this type.

`gcry_md_init_t init`  
The function responsible for initializing a handle. See below for an explanation of this type.

`gcry_md_write_t write`  
The function responsible for writing data into a message digest context. See below for an explanation of this type.

`gcry_md_final_t final`  
The function responsible for ‘finalizing’ a message digest context. See below for an explanation of this type.

`gcry_md_read_t read`  
The function responsible for reading out a message digest result. See below for an explanation of this type.

`size_t contextsize`  
The size of the algorithm-specific ‘context’, that should be allocated for each handle.

`gcry_md_oid_spec_t` [Data type]  
This type is used for associating a user-provided algorithm implementation with certain OIDs. It contains the following members:

`const char *oidstring`  
Textual representation of the OID.

`gcry_md_init_t` [Data type]  
Type for the ‘init’ function, defined as: `void (*gcry_md_init_t) (void *c)`

`gcry_md_write_t` [Data type]  
Type for the ‘write’ function, defined as: `void (*gcry_md_write_t) (void *c, unsigned char *buf, size_t nbytes)`

`gcry_md_final_t` [Data type]  
Type for the ‘final’ function, defined as: `void (*gcry_md_final_t) (void *c)`

`gcry_md_read_t` [Data type]  
Type for the ‘read’ function, defined as: `unsigned char *(*gcry_md_read_t) (void *c)`

`gcry_error_t gcry_md_register (gcry_md_spec_t *digest, unsigned int *algorithm_id, gcry_module_t *module)` [Function]

Register a new digest module whose specification can be found in *digest*. On success, a new algorithm ID is stored in *algorithm\_id* and a pointer representing this module is stored in *module*. Deprecated; the module register interface will be removed in a future version.

`void gcry_md_unregister (gcry_module_t module)` [Function]

Unregister the digest identified by *module*, which must have been registered with `gcry_md_register`.

`gcry_error_t gcry_md_list (int *list, int *list_length)` [Function]

Get a list consisting of the IDs of the loaded message digest modules. If *list* is zero, write the number of loaded message digest modules to *list\_length* and return. If *list* is non-zero, the first *\*list\_length* algorithm IDs are stored in *list*, which must be of according size. In case there are less message digests modules than *\*list\_length*, *\*list\_length* is updated to the correct number.

### 7.3 Working with hash algorithms

To use most of these function it is necessary to create a context; this is done using:

`gcry_error_t gcry_md_open (gcry_md_hd_t *hd, int algo, unsigned int flags)` [Function]

Create a message digest object for algorithm *algo*. *flags* may be given as an bitwise OR of constants described below. *algo* may be given as 0 if the algorithms to use are later set using `gcry_md_enable`. *hd* is guaranteed to either receive a valid handle or NULL.

For a list of supported algorithms, see See [Section 7.1 \[Available hash algorithms\]](#), [page 55](#).

The flags allowed for *mode* are:

**GCRY\_MD\_FLAG\_SECURE**

Allocate all buffers and the resulting digest in "secure memory". Use this is the hashed data is highly confidential.

**GCRY\_MD\_FLAG\_HMAC**

Turn the algorithm into a HMAC message authentication algorithm. This only works if just one algorithm is enabled for the handle. Note that the function `gcry_md_setkey` must be used to set the MAC key. The size of the MAC is equal to the message digest of the underlying hash algorithm. If you want CBC message authentication codes based on a cipher, see See [Section 5.4 \[Working with cipher handles\]](#), [page 27](#).

You may use the function `gcry_md_is_enabled` to later check whether an algorithm has been enabled.

If you want to calculate several hash algorithms at the same time, you have to use the following function right after the `gcry_md_open`:

`gcry_error_t gcry_md_enable (gcry_md_hd_t h, int algo)` [Function]  
 Add the message digest algorithm *algo* to the digest object described by handle *h*. Duplicated enabling of algorithms is detected and ignored.

If the flag `GCRY_MD_FLAG_HMAC` was used, the key for the MAC must be set using the function:

`gcry_error_t gcry_md_setkey (gcry_md_hd_t h, const void *key, size_t keylen)` [Function]  
 For use with the HMAC feature, set the MAC key to the value of *key* of length *keylen* bytes. There is no restriction on the length of the key.

After you are done with the hash calculation, you should release the resources by using:

`void gcry_md_close (gcry_md_hd_t h)` [Function]  
 Release all resources of hash context *h*. *h* should not be used after a call to this function. A NULL passed as *h* is ignored. The function also zeroes all sensitive information associated with this handle.

Often you have to do several hash operations using the same algorithm. To avoid the overhead of creating and releasing context, a reset function is provided:

`void gcry_md_reset (gcry_md_hd_t h)` [Function]  
 Reset the current context to its initial state. This is effectively identical to a close followed by an open and enabling all currently active algorithms.

Often it is necessary to start hashing some data and then continue to hash different data. To avoid hashing the same data several times (which might not even be possible if the data is received from a pipe), a snapshot of the current hash context can be taken and turned into a new context:

`gcry_error_t gcry_md_copy (gcry_md_hd_t *handle_dst, gcry_md_hd_t handle_src)` [Function]  
 Create a new digest object as an exact copy of the object described by handle *handle\_src* and store it in *handle\_dst*. The context is not reset and you can continue to hash data using this context and independently using the original context.

Now that we have prepared everything to calculate hashes, it is time to see how it is actually done. There are two ways for this, one to update the hash with a block of memory and one macro to update the hash by just one character. Both methods can be used on the same hash context.

`void gcry_md_write (gcry_md_hd_t h, const void *buffer, size_t length)` [Function]  
 Pass *length* bytes of the data in *buffer* to the digest object with handle *h* to update the digest values. This function should be used for large blocks of data.

`void gcry_md_putc (gcry_md_hd_t h, int c)` [Function]  
 Pass the byte in *c* to the digest object with handle *h* to update the digest value. This is an efficient function, implemented as a macro to buffer the data before an actual update.

The semantics of the hash functions do not provide for reading out intermediate message digests because the calculation must be finalized first. This finalization may for example include the number of bytes hashed in the message digest or some padding.

```
void gcry_md_final (gcry_md_hd_t h) [Function]
```

Finalize the message digest calculation. This is not really needed because `gcry_md_read` does this implicitly. After this has been done no further updates (by means of `gcry_md_write` or `gcry_md_putc` are allowed. Only the first call to this function has an effect. It is implemented as a macro.

The way to read out the calculated message digest is by using the function:

```
unsigned char * gcry_md_read (gcry_md_hd_t h, int algo) [Function]
```

`gcry_md_read` returns the message digest after finalizing the calculation. This function may be used as often as required but it will always return the same value for one handle. The returned message digest is allocated within the message context and therefore valid until the handle is released or reseted (using `gcry_md_close` or `gcry_md_reset`. `algo` may be given as 0 to return the only enabled message digest or it may specify one of the enabled algorithms. The function does return NULL if the requested algorithm has not been enabled.

Because it is often necessary to get the message digest of one block of memory, a fast convenience function is available for this task:

```
void gcry_md_hash_buffer (int algo, void *digest, const void *buffer, size_t length); [Function]
```

`gcry_md_hash_buffer` is a shortcut function to calculate a message digest of a buffer. This function does not require a context and immediately returns the message digest of the `length` bytes at `buffer`. `digest` must be allocated by the caller, large enough to hold the message digest yielded by the the specified algorithm `algo`. This required size may be obtained by using the function `gcry_md_get_algo_dlen`.

Note that this function will abort the process if an unavailable algorithm is used.

Hash algorithms are identified by internal algorithm numbers (see `gcry_md_open` for a list). However, in most applications they are used by names, so two functions are available to map between string representations and hash algorithm identifiers.

```
const char * gcry_md_algo_name (int algo) [Function]
```

Map the digest algorithm id `algo` to a string representation of the algorithm name. For unknown algorithms this function returns the string "?". This function should not be used to test for the availability of an algorithm.

```
int gcry_md_map_name (const char *name) [Function]
```

Map the algorithm with `name` to a digest algorithm identifier. Returns 0 if the algorithm name is not known. Names representing ASN.1 object identifiers are recognized if the IETF dotted format is used and the OID is prefixed with either "oid." or "OID.". For a list of supported OIDs, see the source code at 'cipher/md.c'. This function should not be used to test for the availability of an algorithm.

`gcry_error_t gcry_md_get_asnoid (int algo, void *buffer, size_t *length)` [Function]

Return an DER encoded ASN.1 OID for the algorithm *algo* in the user allocated *buffer*. *length* must point to variable with the available size of *buffer* and receives after return the actual size of the returned OID. The returned error code may be `GPG_ERR_TOO_SHORT` if the provided buffer is too short to receive the OID; it is possible to call the function with `NULL` for *buffer* to have it only return the required size. The function returns 0 on success.

To test whether an algorithm is actually available for use, the following macro should be used:

`gcry_error_t gcry_md_test_algo (int algo)` [Function]  
The macro returns 0 if the algorithm *algo* is available for use.

If the length of a message digest is not known, it can be retrieved using the following function:

`unsigned int gcry_md_get_algo_dlen (int algo)` [Function]  
Retrieve the length in bytes of the digest yielded by algorithm *algo*. This is often used prior to `gcry_md_read` to allocate sufficient memory for the digest.

In some situations it might be hard to remember the algorithm used for the ongoing hashing. The following function might be used to get that information:

`int gcry_md_get_algo (gcry_md_hd_t h)` [Function]  
Retrieve the algorithm used with the handle *h*. Note that this does not work reliably if more than one algorithm is enabled in *h*.

The following macro might also be useful:

`int gcry_md_is_secure (gcry_md_hd_t h)` [Function]  
This function returns true when the digest object *h* is allocated in "secure memory"; i.e. *h* was created with the `GCRY_MD_FLAG_SECURE`.

`int gcry_md_is_enabled (gcry_md_hd_t h, int algo)` [Function]  
This function returns true when the algorithm *algo* has been enabled for the digest object *h*.

Tracking bugs related to hashing is often a cumbersome task which requires to add a lot of `printf` statements into the code. Libgcrypt provides an easy way to avoid this. The actual data hashed can be written to files on request.

`void gcry_md_debug (gcry_md_hd_t h, const char *suffix)` [Function]  
Enable debugging for the digest object with handle *h*. This creates files named `'dbgmd-<n>.<string>'` while doing the actual hashing. *suffix* is the string part in the filename. The number is a counter incremented for each new hashing. The data in the file is the raw data as passed to `gcry_md_write` or `gcry_md_putc`. If `NULL` is used for *suffix*, the debugging is stopped and the file closed. This is only rarely required because `gcry_md_close` implicitly stops debugging.

The following two deprecated macros are used for debugging by old code. They should be replaced by `gcry_md_debug`.

`void gcry_md_start_debug (gcry_md_hd_t h, const char *suffix)` [Function]  
Enable debugging for the digest object with handle *h*. This creates files named ‘`dbgmd-<n>.<string>`’ while doing the actual hashing. *suffix* is the string part in the filename. The number is a counter incremented for each new hashing. The data in the file is the raw data as passed to `gcry_md_write` or `gcry_md_putc`.

`void gcry_md_stop_debug (gcry_md_hd_t h, int reserved)` [Function]  
Stop debugging on handle *h*. *reserved* should be specified as 0. This function is usually not required because `gcry_md_close` does implicitly stop debugging.

## 8 Key Derivation

Libgcrypt provides a general purpose function to derive keys from strings.

```
gpg_error_t gcry_kdf_derive (const void *passphrase, [Function]
 size_t passphraselen, int algo, int subalgo, const void *salt,
 size_t saltlen, unsigned long iterations, size_t keysize,
 void *keybuffer)
```

Derive a key from a passphrase. *keysize* gives the requested size of the keys in octets. *keybuffer* is a caller provided buffer filled on success with the derived key. The input passphrase is taken from *passphrase* which is an arbitrary memory buffer of *passphraselen* octets. *algo* specifies the KDF algorithm to use; see below. *subalgo* specifies an algorithm used internally by the KDF algorithms; this is usually a hash algorithm but certain KDF algorithms may use it differently. *salt* is a salt of length *saltlen* octets, as needed by most KDF algorithms. *iterations* is a positive integer parameter to most KDFs.

On success 0 is returned; on failure an error code.

Currently supported KDFs (parameter *algo*):

### GCRY\_KDF\_SIMPLE\_S2K

The OpenPGP simple S2K algorithm (cf. RFC4880). Its use is strongly deprecated. *salt* and *iterations* are not needed and may be passed as NULL/0.

### GCRY\_KDF\_SALTED\_S2K

The OpenPGP salted S2K algorithm (cf. RFC4880). Usually not used. *iterations* is not needed and may be passed as 0. *saltlen* must be given as 8.

### GCRY\_KDF\_ITERSALTED\_S2K

The OpenPGP iterated+salted S2K algorithm (cf. RFC4880). This is the default for most OpenPGP applications. *saltlen* must be given as 8. Note that OpenPGP defines a special encoding of the *iterations*; however this function takes the plain decoded iteration count.

### GCRY\_KDF\_PBKDF2

The PKCS#5 Passphrase Based Key Derivation Function number 2.





## 9 Random Numbers

### 9.1 Quality of random numbers

Libgcrypt offers random numbers of different quality levels:

`gcry_random_level_t` [Data type]

The constants for the random quality levels are of this enum type.

`GCRY_WEAK_RANDOM`

For all functions, except for `gcry_mpi_randomize`, this level maps to `GCRY_STRONG_RANDOM`. If you do not want this, consider using `gcry_create_nonce`.

`GCRY_STRONG_RANDOM`

Use this level for session keys and similar purposes.

`GCRY_VERY_STRONG_RANDOM`

Use this level for long term key material.

### 9.2 Retrieving random numbers

`void gcry_randomize (unsigned char *buffer, size_t length, enum gcry_random_level level)` [Function]

Fill *buffer* with *length* random bytes using a random quality as defined by *level*.

`void * gcry_random_bytes (size_t nbytes, enum gcry_random_level level)` [Function]

Convenience function to allocate a memory block consisting of *nbytes* fresh random bytes using a random quality as defined by *level*.

`void * gcry_random_bytes_secure (size_t nbytes, enum gcry_random_level level)` [Function]

Convenience function to allocate a memory block consisting of *nbytes* fresh random bytes using a random quality as defined by *level*. This function differs from `gcry_random_bytes` in that the returned buffer is allocated in a “secure” area of the memory.

`void gcry_create_nonce (unsigned char *buffer, size_t length)` [Function]

Fill *buffer* with *length* unpredictable bytes. This is commonly called a nonce and may also be used for initialization vectors and padding. This is an extra function nearly independent of the other random function for 3 reasons: It better protects the regular random generator’s internal state, provides better performance and does not drain the precious entropy pool.



## 10 S-expressions

S-expressions are used by the public key functions to pass complex data structures around. These LISP like objects are used by some cryptographic protocols (cf. RFC-2692) and Libgcrypt provides functions to parse and construct them. For detailed information, see *Ron Rivest, code and description of S-expressions*, <http://theory.lcs.mit.edu/~rivest/sexp.html>.

### 10.1 Data types for S-expressions

`gcry_sexp_t` [Data type]

The `gcry_sexp_t` type describes an object with the Libgcrypt internal representation of an S-expression.

### 10.2 Working with S-expressions

There are several functions to create an Libgcrypt S-expression object from its external representation or from a string template. There is also a function to convert the internal representation back into one of the external formats:

`gcry_error_t gcry_sexp_new (gcry_sexp_t *r_sexp, [Function]  
const void *buffer, size_t length, int autodetect)`

This is the generic function to create a new S-expression object from its external representation in *buffer* of *length* bytes. On success the result is stored at the address given by *r\_sexp*. With *autodetect* set to 0, the data in *buffer* is expected to be in canonized format, with *autodetect* set to 1 the parses any of the defined external formats. If *buffer* does not hold a valid S-expression an error code is returned and *r\_sexp* set to NULL. Note that the caller is responsible for releasing the newly allocated S-expression using `gcry_sexp_release`.

`gcry_error_t gcry_sexp_create (gcry_sexp_t *r_sexp, [Function]  
void *buffer, size_t length, int autodetect, void (*freefunc)(void*))`

This function is identical to `gcry_sexp_new` but has an extra argument *freefunc*, which, when not set to NULL, is expected to be a function to release the *buffer*; most likely the standard `free` function is used for this argument. This has the effect of transferring the ownership of *buffer* to the created object in *r\_sexp*. The advantage of using this function is that Libgcrypt might decide to directly use the provided buffer and thus avoid extra copying.

`gcry_error_t gcry_sexp_sscan (gcry_sexp_t *r_sexp, [Function]  
size_t *erhoff, const char *buffer, size_t length)`

This is another variant of the above functions. It behaves nearly identical but provides an *erhoff* argument which will receive the offset into the buffer where the parsing stopped on error.

`gcry_error_t gcry_sexp_build (gcry_sexp_t *r_sexp, [Function]  
size_t *erhoff, const char *format, ...)`

This function creates an internal S-expression from the string template *format* and stores it at the address of *r\_sexp*. If there is a parsing error, the function returns an

appropriate error code and stores the offset into *format* where the parsing stopped in *erroff*. The function supports a couple of printf-like formatting characters and expects arguments for some of these escape sequences right after *format*. The following format characters are defined:

- '%m'        The next argument is expected to be of type `gcry_mpi_t` and a copy of its value is inserted into the resulting S-expression. The MPI is stored as a signed integer.
- '%M'        The next argument is expected to be of type `gcry_mpi_t` and a copy of its value is inserted into the resulting S-expression. The MPI is stored as an unsigned integer.
- '%s'        The next argument is expected to be of type `char *` and that string is inserted into the resulting S-expression.
- '%d'        The next argument is expected to be of type `int` and its value is inserted into the resulting S-expression.
- '%u'        The next argument is expected to be of type `unsigned int` and its value is inserted into the resulting S-expression.
- '%b'        The next argument is expected to be of type `int` directly followed by an argument of type `char *`. This represents a buffer of given length to be inserted into the resulting S-expression.
- '%S'        The next argument is expected to be of type `gcry_sexp_t` and a copy of that S-expression is embedded in the resulting S-expression. The argument needs to be a regular S-expression, starting with a parenthesis.

No other format characters are defined and would return an error. Note that the format character `'%%'` does not exist, because a percent sign is not a valid character in an S-expression.

`void gcry_sexp_release (gcry_sexp_t sexp)` [Function]  
 Release the S-expression object *sexp*. If the S-expression is stored in secure memory it explicitly zeroes that memory; note that this is done in addition to the zeroisation always done when freeing secure memory.

The next 2 functions are used to convert the internal representation back into a regular external S-expression format and to show the structure for debugging.

`size_t gcry_sexp_sprint (gcry_sexp_t sexp, int mode,  
                           char *buffer, size_t maxlength)` [Function]

Copies the S-expression object *sexp* into *buffer* using the format specified in *mode*. *maxlength* must be set to the allocated length of *buffer*. The function returns the actual length of valid bytes put into *buffer* or 0 if the provided buffer is too short. Passing NULL for *buffer* returns the required length for *buffer*. For convenience reasons an extra byte with value 0 is appended to the buffer.

The following formats are supported:

`GCRYSEXP_FMT_DEFAULT`

Returns a convenient external S-expression representation.

`GCRYSEXP_FMT_CANON`  
Return the S-expression in canonical format.

`GCRYSEXP_FMT_BASE64`  
Not currently supported.

`GCRYSEXP_FMT_ADVANCED`  
Returns the S-expression in advanced format.

`void gcry_sexp_dump (gcry_sexp_t sexp)` [Function]  
Dumps *sexp* in a format suitable for debugging to Libgcrypt's logging stream.

Often canonical encoding is used in the external representation. The following function can be used to check for valid encoding and to learn the length of the S-expression"

`size_t gcry_sexp_canon_len (const unsigned char *buffer,` [Function]  
`size_t length, size_t *erhoff, int *errcode)`  
Scan the canonical encoded *buffer* with implicit length values and return the actual length this S-expression uses. For a valid S-expression it should never return 0. If *length* is not 0, the maximum length to scan is given; this can be used for syntax checks of data passed from outside. *errcode* and *erhoff* may both be passed as NULL.

There are functions to parse S-expressions and retrieve elements:

`gcry_sexp_t gcry_sexp_find_token (const gcry_sexp_t list,` [Function]  
`const char *token, size_t toklen)`  
Scan the S-expression for a sublist with a type (the car of the list) matching the string *token*. If *toklen* is not 0, the token is assumed to be raw memory of this length. The function returns a newly allocated S-expression consisting of the found sublist or NULL when not found.

`int gcry_sexp_length (const gcry_sexp_t list)` [Function]  
Return the length of the *list*. For a valid S-expression this should be at least 1.

`gcry_sexp_t gcry_sexp_nth (const gcry_sexp_t list, int number)` [Function]  
Create and return a new S-expression from the element with index *number* in *list*. Note that the first element has the index 0. If there is no such element, NULL is returned.

`gcry_sexp_t gcry_sexp_car (const gcry_sexp_t list)` [Function]  
Create and return a new S-expression from the first element in *list*; this called the "type" and should always exist and be a string. NULL is returned in case of a problem.

`gcry_sexp_t gcry_sexp_cdr (const gcry_sexp_t list)` [Function]  
Create and return a new list form all elements except for the first one. Note that this function may return an invalid S-expression because it is not guaranteed, that the type exists and is a string. However, for parsing a complex S-expression it might be useful for intermediate lists. Returns NULL on error.

```
const char * gcry_sexp_nth_data (const gcry_sexp_t list, [Function]
 int number, size_t *datalen)
```

This function is used to get data from a *list*. A pointer to the actual data with index *number* is returned and the length of this data will be stored to *datalen*. If there is no data at the given index or the index represents another list, NULL is returned.

**Caution:** The returned pointer is valid as long as *list* is not modified or released.

Here is an example on how to extract and print the surname (Meier) from the S-expression '(Name Otto Meier (address Burgplatz 3))':

```
size_t len;
const char *name;

name = gcry_sexp_nth_data (list, 2, &len);
printf ("my name is %.*s\n", (int)len, name);
```

```
char * gcry_sexp_nth_string (gcry_sexp_t list, int number) [Function]
```

This function is used to get and convert data from a *list*. The data is assumed to be a Nul terminated string. The caller must release this returned value using `gcry_free`. If there is no data at the given index, the index represents a list or the value can't be converted to a string, NULL is returned.

```
gcry_mpi_t gcry_sexp_nth_mpi (gcry_sexp_t list, int number, [Function]
 int mpifmt)
```

This function is used to get and convert data from a *list*. This data is assumed to be an MPI stored in the format described by *mpifmt* and returned as a standard Libgcrypt MPI. The caller must release this returned value using `gcry_mpi_release`. If there is no data at the given index, the index represents a list or the value can't be converted to an MPI, NULL is returned. If you use this function to parse results of a public key function, you most likely want to use `GCRYMPI_FMT_USG`.

## 11 MPI library

Public key cryptography is based on mathematics with large numbers. To implement the public key functions, a library for handling these large numbers is required. Because of the general usefulness of such a library, its interface is exposed by Libgcrypt. In the context of Libgcrypt and in most other applications, these large numbers are called MPIs (multi-precision-integers).

### 11.1 Data types

`gcry_mpi_t` [Data type]  
 This type represents an object to hold an MPI.

### 11.2 Basic functions

To work with MPIs, storage must be allocated and released for the numbers. This can be done with one of these functions:

`gcry_mpi_t gcry_mpi_new (unsigned int nbits)` [Function]  
 Allocate a new MPI object, initialize it to 0 and initially allocate enough memory for a number of at least *nbits*. This pre-allocation is only a small performance issue and not actually necessary because Libgcrypt automatically re-allocates the required memory.

`gcry_mpi_t gcry_mpi_snew (unsigned int nbits)` [Function]  
 This is identical to `gcry_mpi_new` but allocates the MPI in the so called "secure memory" which in turn will take care that all derived values will also be stored in this "secure memory". Use this for highly confidential data like private key parameters.

`gcry_mpi_t gcry_mpi_copy (const gcry_mpi_t a)` [Function]  
 Create a new MPI as the exact copy of *a*.

`void gcry_mpi_release (gcry_mpi_t a)` [Function]  
 Release the MPI *a* and free all associated resources. Passing NULL is allowed and ignored. When a MPI stored in the "secure memory" is released, that memory gets wiped out immediately.

The simplest operations are used to assign a new value to an MPI:

`gcry_mpi_t gcry_mpi_set (gcry_mpi_t w, const gcry_mpi_t u)` [Function]  
 Assign the value of *u* to *w* and return *w*. If NULL is passed for *w*, a new MPI is allocated, set to the value of *u* and returned.

`gcry_mpi_t gcry_mpi_set_ui (gcry_mpi_t w, unsigned long u)` [Function]  
 Assign the value of *u* to *w* and return *w*. If NULL is passed for *w*, a new MPI is allocated, set to the value of *u* and returned. This function takes an `unsigned int` as type for *u* and thus it is only possible to set *w* to small values (usually up to the word size of the CPU).

`void gcry_mpi_swap (gcry_mpi_t a, gcry_mpi_t b)` [Function]  
 Swap the values of *a* and *b*.

### 11.3 MPI formats

The following functions are used to convert between an external representation of an MPI and the internal one of Libcrypt.

```
gcry_error_t gcry_mpi_scan (gcry_mpi_t *r_mpi, [Function]
 enum gcry_mpi_format format, const unsigned char *buffer, size_t buflen,
 size_t *nscanned)
```

Convert the external representation of an integer stored in *buffer* with a length of *buflen* into a newly created MPI returned which will be stored at the address of *r\_mpi*. For certain formats the length argument is not required and should be passed as 0. After a successful operation the variable *nscanned* receives the number of bytes actually scanned unless *nscanned* was given as NULL. *format* describes the format of the MPI as stored in *buffer*:

**GCRYMPI\_FMT\_STD**

2-complement stored without a length header.

**GCRYMPI\_FMT\_PGP**

As used by OpenPGP (only defined as unsigned). This is basically GCRYMPI\_FMT\_STD with a 2 byte big endian length header.

**GCRYMPI\_FMT\_SSH**

As used in the Secure Shell protocol. This is GCRYMPI\_FMT\_STD with a 4 byte big endian header.

**GCRYMPI\_FMT\_HEX**

Stored as a C style string with each byte of the MPI encoded as 2 hex digits. When using this format, *buflen* must be zero.

**GCRYMPI\_FMT\_USG**

Simple unsigned integer.

Note that all of the above formats store the integer in big-endian format (MSB first).

```
gcry_error_t gcry_mpi_print (enum gcry_mpi_format format, [Function]
 unsigned char *buffer, size_t buflen, size_t *nwritten,
 const gcry_mpi_t a)
```

Convert the MPI *a* into an external representation described by *format* (see above) and store it in the provided *buffer* which has a usable length of at least the *buflen* bytes. If *nwritten* is not NULL, it will receive the number of bytes actually stored in *buffer* after a successful operation.

```
gcry_error_t gcry_mpi_aprint (enum gcry_mpi_format format, [Function]
 unsigned char **buffer, size_t *nbytes, const gcry_mpi_t a)
```

Convert the MPI *a* into an external representation described by *format* (see above) and store it in a newly allocated buffer which address will be stored in the variable *buffer* points to. The number of bytes stored in this buffer will be stored in the variable *nbytes* points to, unless *nbytes* is NULL.



`void gcry_mpi_dump (const gcry_mpi_t a)` [Function]  
 Dump the value of *a* in a format suitable for debugging to Libgcrypt's logging stream. Note that one leading space but no trailing space or linefeed will be printed. It is okay to pass NULL for *a*.

## 11.4 Calculations

Basic arithmetic operations:

`void gcry_mpi_add (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v)` [Function]  
 $w = u + v$ .

`void gcry_mpi_add_ui (gcry_mpi_t w, gcry_mpi_t u, unsigned long v)` [Function]  
 $w = u + v$ . Note that *v* is an unsigned integer.

`void gcry_mpi_addm (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v, gcry_mpi_t m)` [Function]  
 $w = u + v \bmod m$ .

`void gcry_mpi_sub (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v)` [Function]  
 $w = u - v$ .

`void gcry_mpi_sub_ui (gcry_mpi_t w, gcry_mpi_t u, unsigned long v)` [Function]  
 $w = u - v$ . *v* is an unsigned integer.

`void gcry_mpi_subm (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v, gcry_mpi_t m)` [Function]  
 $w = u - v \bmod m$ .

`void gcry_mpi_mul (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v)` [Function]  
 $w = u * v$ .

`void gcry_mpi_mul_ui (gcry_mpi_t w, gcry_mpi_t u, unsigned long v)` [Function]  
 $w = u * v$ . *v* is an unsigned integer.

`void gcry_mpi_mulm (gcry_mpi_t w, gcry_mpi_t u, gcry_mpi_t v, gcry_mpi_t m)` [Function]  
 $w = u * v \bmod m$ .

`void gcry_mpi_mul_2exp (gcry_mpi_t w, gcry_mpi_t u, unsigned long e)` [Function]  
 $w = u * 2^e$ .

`void gcry_mpi_div (gcry_mpi_t q, gcry_mpi_t r, gcry_mpi_t dividend, gcry_mpi_t divisor, int round)` [Function]  
 $q = \text{dividend} / \text{divisor}$ ,  $r = \text{dividend} \bmod \text{divisor}$ . *q* and *r* may be passed as NULL. *round* should be negative or 0.

`void gcry_mpi_mod (gcry_mpi_t r, gcry_mpi_t dividend, gcry_mpi_t divisor)` [Function]  
 $r = \text{dividend} \bmod \text{divisor}$ .

`void gcry_mpi_powm (gcry_mpi_t w, const gcry_mpi_t b, const gcry_mpi_t e, const gcry_mpi_t m)` [Function]  
 $w = b^e \bmod m$ .

`int gcry_mpi_gcd (gcry_mpi_t g, gcry_mpi_t a, gcry_mpi_t b)` [Function]  
 Set  $g$  to the greatest common divisor of  $a$  and  $b$ . Return true if the  $g$  is 1.

`int gcry_mpi_invmod (gcry_mpi_t x, gcry_mpi_t a, gcry_mpi_t m)` [Function]  
 Set  $x$  to the multiplicative inverse of  $a \bmod m$ . Return true if the inverse exists.

## 11.5 Comparisons

The next 2 functions are used to compare MPIs:

`int gcry_mpi_cmp (const gcry_mpi_t u, const gcry_mpi_t v)` [Function]  
 Compare the multi-precision-integers number  $u$  and  $v$  returning 0 for equality, a positive value for  $u > v$  and a negative for  $u < v$ . If both numbers are opaque values (cf, `gcry_mpi_set_opaque`) the comparison is done by checking the bit sizes using `memcmp`. If only one number is an opaque value, the opaque value is less than the other number.

`int gcry_mpi_cmp_ui (const gcry_mpi_t u, unsigned long v)` [Function]  
 Compare the multi-precision-integers number  $u$  with the unsigned integer  $v$  returning 0 for equality, a positive value for  $u > v$  and a negative for  $u < v$ .

## 11.6 Bit manipulations

There are a couple of functions to get information on arbitrary bits in an MPI and to set or clear them:

`unsigned int gcry_mpi_get_nbits (gcry_mpi_t a)` [Function]  
 Return the number of bits required to represent  $a$ .

`int gcry_mpi_test_bit (gcry_mpi_t a, unsigned int n)` [Function]  
 Return true if bit number  $n$  (counting from 0) is set in  $a$ .

`void gcry_mpi_set_bit (gcry_mpi_t a, unsigned int n)` [Function]  
 Set bit number  $n$  in  $a$ .

`void gcry_mpi_clear_bit (gcry_mpi_t a, unsigned int n)` [Function]  
 Clear bit number  $n$  in  $a$ .

`void gcry_mpi_set_highbit (gcry_mpi_t a, unsigned int n)` [Function]  
 Set bit number  $n$  in  $a$  and clear all bits greater than  $n$ .

`void gcry_mpi_clear_highbit (gcry_mpi_t a, unsigned int n)` [Function]  
 Clear bit number  $n$  in  $a$  and all bits greater than  $n$ .

`void gcry_mpi_rshift (gcry_mpi_t x, gcry_mpi_t a, unsigned int n)` [Function]  
 Shift the value of  $a$  by  $n$  bits to the right and store the result in  $x$ .

`void gcry_mpi_lshift (gcry_mpi_t x, gcry_mpi_t a, unsigned int n)` [Function]  
 Shift the value of  $a$  by  $n$  bits to the left and store the result in  $x$ .

## 11.7 Miscellaneous

`gcry_mpi_t gcry_mpi_set_opaque (gcry_mpi_t a, void *p, unsigned int nbits)` [Function]

Store *nbits* of the value *p* points to in *a* and mark *a* as an opaque value (i.e. an value that can't be used for any math calculation and is only used to store an arbitrary bit pattern in *a*).

WARNING: Never use an opaque MPI for actual math operations. The only valid functions are `gcry_mpi_get_opaque` and `gcry_mpi_release`. Use `gcry_mpi_scan` to convert a string of arbitrary bytes into an MPI.

`void * gcry_mpi_get_opaque (gcry_mpi_t a, unsigned int *nbits)` [Function]

Return a pointer to an opaque value stored in *a* and return its size in *nbits*. Note that the returned pointer is still owned by *a* and that the function should never be used for a non-opaque MPI.

`void gcry_mpi_set_flag (gcry_mpi_t a, enum gcry_mpi_flag flag)` [Function]

Set the *flag* for the MPI *a*. Currently only the flag `GCRYMPI_FLAG_SECURE` is allowed to convert *a* into an MPI stored in "secure memory".

`void gcry_mpi_clear_flag (gcry_mpi_t a, enum gcry_mpi_flag flag)` [Function]

Clear *flag* for the multi-precision-integers *a*. Note that this function is currently useless as no flags are allowed.

`int gcry_mpi_get_flag (gcry_mpi_t a, enum gcry_mpi_flag flag)` [Function]

Return true when the *flag* is set for *a*.

`void gcry_mpi_randomize (gcry_mpi_t w, unsigned int nbits, enum gcry_random_level level)` [Function]

Set the multi-precision-integers *w* to a random value of *nbits*, using random data quality of level *level*. In case *nbits* is not a multiple of a byte, *nbits* is rounded up to the next byte boundary. When using a *level* of `GCRY_WEAK_RANDOM` this function makes use of `gcry_create_nonce`.



## 12 Prime numbers

### 12.1 Generation

`gcry_error_t gcry_prime_generate (gcry_mpi_t *prime, unsigned int prime_bits, unsigned int factor_bits, gcry_mpi_t **factors, gcry_prime_check_func_t cb_func, void *cb_arg, gcry_random_level_t random_level, unsigned int flags)` [Function]

Generate a new prime number of *prime\_bits* bits and store it in *prime*. If *factor\_bits* is non-zero, one of the prime factors of  $(prime - 1) / 2$  must be *factor\_bits* bits long. If *factors* is non-zero, allocate a new, NULL-terminated array holding the prime factors and store it in *factors*. *flags* might be used to influence the prime number generation process.

`gcry_error_t gcry_prime_group_generator (gcry_mpi_t *r_g, gcry_mpi_t prime, gcry_mpi_t *factors, gcry_mpi_t start_g)` [Function]

Find a generator for *prime* where the factorization of  $(prime-1)$  is in the NULL terminated array *factors*. Return the generator as a newly allocated MPI in *r\_g*. If *start\_g* is not NULL, use this as the start for the search.

`void gcry_prime_release_factors (gcry_mpi_t *factors)` [Function]

Convenience function to release the *factors* array.

### 12.2 Checking

`gcry_error_t gcry_prime_check (gcry_mpi_t p, unsigned int flags)` [Function]

Check whether the number *p* is prime. Returns zero in case *p* is indeed a prime, returns `GPG_ERR_NO_PRIME` in case *p* is not a prime and a different error code in case something went horribly wrong.



## 13 Utilities

### 13.1 Memory allocation

`void * gcry_malloc (size_t n)` [Function]

This function tries to allocate  $n$  bytes of memory. On success it returns a pointer to the memory area, in an out-of-core condition, it returns NULL.

`void * gcry_malloc_secure (size_t n)` [Function]

Like `gcry_malloc`, but uses secure memory.

`void * gcry_calloc (size_t n, size_t m)` [Function]

This function allocates a cleared block of memory (i.e. initialized with zero bytes) long enough to contain a vector of  $n$  elements, each of size  $m$  bytes. On success it returns a pointer to the memory block; in an out-of-core condition, it returns NULL.

`void * gcry_calloc_secure (size_t n, size_t m)` [Function]

Like `gcry_calloc`, but uses secure memory.

`void * gcry_realloc (void *p, size_t n)` [Function]

This function tries to resize the memory area pointed to by  $p$  to  $n$  bytes. On success it returns a pointer to the new memory area, in an out-of-core condition, it returns NULL. Depending on whether the memory pointed to by  $p$  is secure memory or not, `gcry_realloc` tries to use secure memory as well.

`void gcry_free (void *p)` [Function]

Release the memory area pointed to by  $p$ .





## 14 Architecture

This chapter describes the internal architecture of Libgcrypt.

Libgcrypt is a function library written in ISO C-90. Any compliant compiler should be able to build Libgcrypt as long as the target is either a POSIX platform or compatible to the API used by Windows NT. Provisions have been take so that the library can be directly used from C++ applications; however building with a C++ compiler is not supported.

Building Libgcrypt is done by using the common `./configure && make` approach. The configure command is included in the source distribution and as a portable shell script it works on any Unix-alike system. The result of running the configure script are a C header file (`'config.h'`), customized Makefiles, the setup of symbolic links and a few other things. After that the make tool builds and optionally installs the library and the documentation. See the files `'INSTALL'` and `'README'` in the source distribution on how to do this.

Libgcrypt is developed using a Subversion<sup>1</sup> repository. Although all released versions are tagged in this repository, they should not be used to build production versions of Libgcrypt. Instead released tarballs should be used. These tarballs are available from several places with the master copy at `ftp://ftp.gnupg.org/gcrypt/libgcrypt/`. Announcements of new releases are posted to the `gnupg-announce@gnupg.org` mailing list<sup>2</sup>.

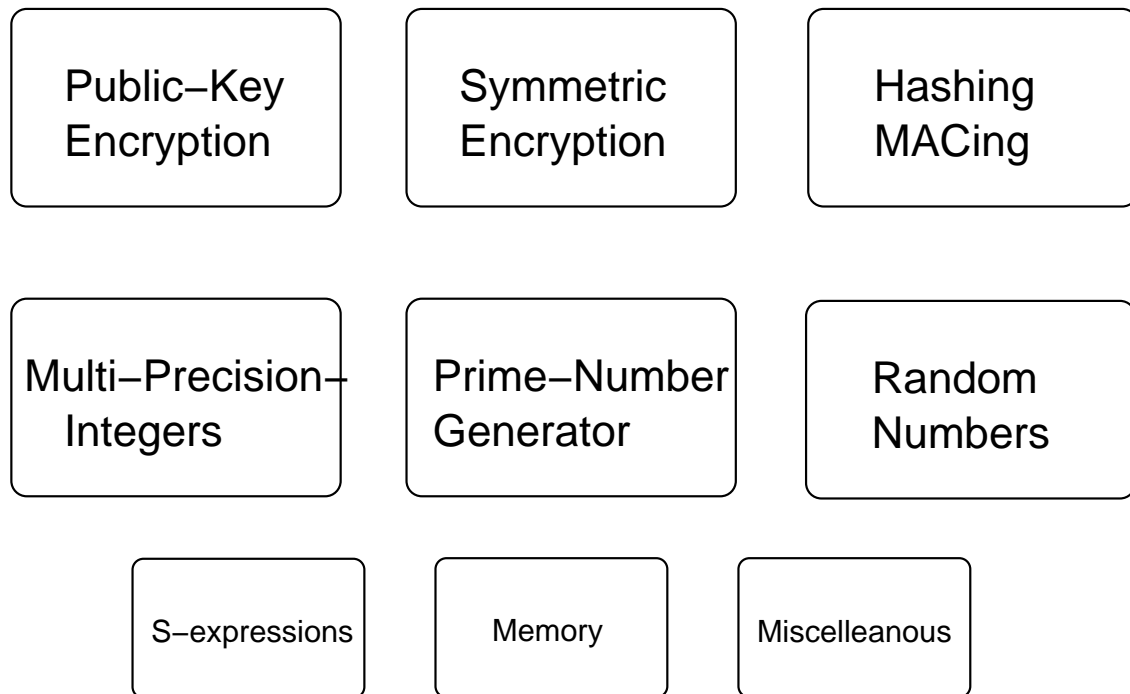


Figure 14.1: Libgcrypt subsystems

Libgcrypt consists of several subsystems (see Figure 14.1) and all these subsystems provide a public API; this includes the helper subsystems like the one for S-expressions. The API style depends on the subsystem; in general an open-use-close approach is implemented.

<sup>1</sup> A version control system available for many platforms

<sup>2</sup> See <http://www.gnupg.org/documentation/mailling-lists.en.html> for details.

The open returns a handle to a context used for all further operations on this handle, several functions may then be used on this handle and a final close function releases all resources associated with the handle.

## 14.1 Public-Key Architecture

Libcrypt implements two interfaces for public key cryptography: The standard interface is PK interface using functions in the `gcry_pk_` name space. The AC interface in an alternative one which is now deprecated and will not be further described. The AC interface is also disabled in FIPS mode.

Because public key cryptography is almost always used to process small amounts of data (hash values or session keys), the interface is not implemented using the open-use-close paradigm, but with single self-contained functions. Due to the wide variety of parameters required by different algorithms S-expressions, as flexible way to convey these parameters, are used. There is a set of helper functions to work with these S-expressions.

Aside of functions to register new algorithms, map algorithms names to algorithms identifiers and to lookup properties of a key, the following main functions are available:

```
gcry_pk_encrypt
 Encrypt data using a public key.

gcry_pk_decrypt
 Decrypt data using a private key.

gcry_pk_sign
 Sign data using a private key.

gcry_pk_verify
 Verify that a signature matches the data.

gcry_pk_testkey
 Perform a consistency over a public or private key.

gcry_pk_genkey
 Create a new public/private key pair.
```

With the help of the module registration system all these functions lookup the module implementing the algorithm and pass the actual work to that module. The parsing of the S-expression input and the construction of S-expression for the return values is done by the high level code (`'cipher/pubkey.c'`). Thus the internal interface between the algorithm modules and the high level functions passes data in a custom format. The interface to the modules is published (`'gcrypt-modules.h'`) so that it can be used to register external implementations of algorithms with Libcrypt. However, for some algorithms this module interface is too limited and thus for the internal modules an extra interface is sometimes used to convey more information.

By default Libcrypt uses a blinding technique for RSA decryption to mitigate real world timing attacks over a network: Instead of using the RSA decryption directly, a blinded value  $y = xr^e \bmod n$  is decrypted and the unblinded value  $x' = y'r^{-1} \bmod n$  returned. The blinding value  $r$  is a random value with the size of the modulus  $n$  and generated with `GCRY_WEAK_RANDOM` random level.

The algorithm used for RSA and DSA key generation depends on whether Libgcrypt is operated in standard or in FIPS mode. In standard mode an algorithm based on the Lim-Lee prime number generator is used. In FIPS mode RSA keys are generated as specified in ANSI X9.31 (1998) and DSA keys as specified in FIPS 186-2.

## 14.2 Symmetric Encryption Subsystem Architecture

The interface to work with symmetric encryption algorithms is made up of functions from the `gcry_cipher_` name space. The implementation follows the open-use-close paradigm and uses registered algorithm modules for the actual work. Unless a module implements optimized cipher mode implementations, the high level code (`'cipher/cipher.c'`) implements the modes and calls the core algorithm functions to process each block.

The most important functions are:

`gcry_cipher_open`  
Create a new instance to encrypt or decrypt using a specified algorithm and mode.

`gcry_cipher_close`  
Release an instance.

`gcry_cipher_setkey`  
Set a key to be used for encryption or decryption.

`gcry_cipher_setiv`  
Set an initialization vector to be used for encryption or decryption.

`gcry_cipher_encrypt`  
`gcry_cipher_decrypt`  
Encrypt or decrypt data. These functions may be called with arbitrary amounts of data and as often as needed to encrypt or decrypt all data.

There are also functions to query properties of algorithms or context, like block length, key length, map names or to enable features like padding methods.

## 14.3 Hashing and MACing Subsystem Architecture

The interface to work with message digests and CRC algorithms is made up of functions from the `gcry_md_` name space. The implementation follows the open-use-close paradigm and uses registered algorithm modules for the actual work. Although CRC algorithms are not considered cryptographic hash algorithms, they share enough properties so that it makes sense to handle them in the same way. It is possible to use several algorithms at once with one context and thus compute them all on the same data.

The most important functions are:

`gcry_md_open`  
Create a new message digest instance and optionally enable one algorithm. A flag may be used to turn the message digest algorithm into a HMAC algorithm.

`gcry_md_enable`  
Enable an additional algorithm for the instance.

`gcry_md_setkey`

Set the key for the MAC.

`gcry_md_write`

Pass more data for computing the message digest to an instance.

`gcry_md_putc`

Buffered version of `gcry_md_write` implemented as a macro.

`gcry_md_read`

Finalize the computation of the message digest or HMAC and return the result.

`gcry_md_close`

Release an instance

`gcry_md_hash_buffer`

Convenience function to directly compute a message digest over a memory buffer without the need to create an instance first.

There are also functions to query properties of algorithms or the instance, like enabled algorithms, digest length, map algorithm names. It is also possible to reset an instance or to copy the current state of an instance at any time. Debug functions to write the hashed data to files are available as well.

## 14.4 Multi-Precision-Integer Subsystem Architecture

The implementation of Libcrypt's big integer computation code is based on an old release of GNU Multi-Precision Library (GMP). The decision not to use the GMP library directly was due to stalled development at that time and due to security requirements which could not be provided by the code in GMP. As GMP does, Libcrypt provides high performance assembler implementations of low level code for several CPUs to gain much better performance than with a generic C implementation.

Major features of Libcrypt's multi-precision-integer code compared to GMP are:

- Avoidance of stack based allocations to allow protection against swapping out of sensitive data and for easy zeroing of sensitive intermediate results.
- Optional use of secure memory and tracking of its use so that results are also put into secure memory.
- MPIs are identified by a handle (implemented as a pointer) to give better control over allocations and to augment them with extra properties like opaque data.
- Removal of unnecessary code to reduce complexity.
- Functions specialized for public key cryptography.

## 14.5 Prime-Number-Generator Subsystem Architecture

Libcrypt provides an interface to its prime number generator. These functions make use of the internal prime number generator which is required for the generation for public key pairs. The plain prime checking function is exported as well.

The generation of random prime numbers is based on the Lim and Lee algorithm to create practically save primes.<sup>3</sup> This algorithm creates a pool of smaller primes, select a few of them to create candidate primes of the form  $2 * p_0 * p_1 * \dots * p_n + 1$ , tests the candidate for primality and permutates the pool until a prime has been found. It is possible to clamp one of the small primes to a certain size to help DSA style algorithms. Because most of the small primes in the pool are not used for the resulting prime number, they are saved for later use (see `save_pool_prime` and `get_pool_prime` in `'cipher/primegen.c'`). The prime generator optionally supports the finding of an appropriate generator.

The primality test works in three steps:

1. The standard sieve algorithm using the primes up to 4999 is used as a quick first check.
2. A Fermat test filters out almost all non-primes.
3. A 5 round Rabin-Miller test is finally used. The first round uses a witness of 2, whereas the next rounds use a random witness.

To support the generation of RSA and DSA keys in FIPS mode according to X9.31 and FIPS 186-2, Libgcrypt implements two additional prime generation functions: `_gcry_derive_x931_prime` and `_gcry_generate_fips186_2_prime`. These functions are internal and not available through the public API.

## 14.6 Random-Number Subsystem Architecture

Libgcrypt provides 3 levels of random quality: The level `GCRY_VERY_STRONG_RANDOM` usually used for key generation, the level `GCRY_STRONG_RANDOM` for all other strong random requirements and the function `gcry_create_nonce` which is used for weaker usages like nonces. There is also a level `GCRY_WEAK_RANDOM` which in general maps to `GCRY_STRONG_RANDOM` except when used with the function `gcry_mpi_randomize`, where it randomizes an multi-precision-integer using the `gcry_create_nonce` function.

There are two distinct random generators available:

- The Continuously Seeded Pseudo Random Number Generator (CSPRNG), which is based on the classic GnuPG derived big pool implementation. Implemented in `random/random-csprng.c` and used by default.
- A FIPS approved ANSI X9.31 PRNG using AES with a 128 bit key. Implemented in `random/random-fips.c` and used if Libgcrypt is in FIPS mode.

Both generators make use of so-called entropy gathering modules:

|                       |                                                                                                                                                                                                                                           |
|-----------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| <code>rndlinux</code> | Uses the operating system provided <code>'/dev/random'</code> and <code>'/dev/urandom'</code> devices.                                                                                                                                    |
| <code>rndunix</code>  | Runs several operating system commands to collect entropy from sources like virtual machine and process statistics. It is a kind of poor-man's <code>/dev/random</code> implementation. It is not available in FIPS mode.                 |
| <code>rndegd</code>   | Uses the operating system provided Entropy Gathering Daemon (EGD). The EGD basically uses the same algorithms as <code>rndunix</code> does. However as a system daemon it keeps on running and thus can serve several processes requiring |

<sup>3</sup> Chae Hoon Lim and Pil Joong Lee. A key recovery attack on discrete log-based schemes using a prime order subgroup. In Burton S. Kaliski Jr., editor, *Advances in Cryptology: Crypto '97*, pages 249-263, Berlin / Heidelberg / New York, 1997. Springer-Verlag. Described on page 260.

entropy input and does not waste collected entropy if the application does not need all the collected entropy. It is not available in FIPS mode.

- `rndw32` Targeted for the Microsoft Windows OS. It uses certain properties of that system and is the only gathering module available for that OS.
- `rndhw` Extra module to collect additional entropy by utilizing a hardware random number generator. As of now the only supported hardware RNG is the Padlock engine of VIA (Centaur) CPUs. It is not available in FIPS mode.

### 14.6.1 Description of the CSPRNG

This random number generator is loosely modelled after the one described in Peter Gutmann's paper: "Software Generation of Practically Strong Random Numbers".<sup>4</sup>

A pool of 600 bytes is used and mixed using the core RIPE-MD160 hash transform function. Several extra features are used to make the robust against a wide variety of attacks and to protect against failures of subsystems. The state of the generator may be saved to a file and initially seed form a file.

Depending on how Libcrypt was build the generator is able to select the best working entropy gathering module. It makes use of the slow and fast collection methods and requires the pool to initially seeded form the slow gatherer or a seed file. An entropy estimation is used to mix in enough data from the gather modules before returning the actual random output. Process fork detection and protection is implemented.

The implementation of the nonce generator (for `gcry_create_nonce`) is a straightforward repeated hash design: A 28 byte buffer is initially seeded with the PID and the time in seconds in the first 20 bytes and with 8 bytes of random taken from the `GCRY_STRONG_RANDOM` generator. Random numbers are then created by hashing all the 28 bytes with SHA-1 and saving that again in the first 20 bytes. The hash is also returned as result.

### 14.6.2 Description of the FIPS X9.31 PRNG

The core of this deterministic random number generator is implemented according to the document "NIST-Recommended Random Number Generator Based on ANSI X9.31 Appendix A.2.4 Using the 3-Key Triple DES and AES Algorithms", dated 2005-01-31. This implementation uses the AES variant.

The generator is based on contexts to utilize the same core functions for all random levels as required by the high-level interface. All random generators return their data in 128 bit blocks. If the caller requests less bits, the extra bits are not used. The key for each generator is only set once at the first time a generator context is used. The seed value is set along with the key and again after 1000 output blocks.

On Unix like systems the `GCRY_VERY_STRONG_RANDOM` and `GCRY_STRONG_RANDOM` generators are keyed and seeded using the `rndlinux` module with the `'/dev/radnom'` device. Thus these generators may block until the OS kernel has collected enough entropy. When used with Microsoft Windows the `rndw32` module is used instead.

The generator used for `gcry_create_nonce` is keyed and seeded from the `GCRY_STRONG_RANDOM` generator. Thus is may also block if the `GCRY_STRONG_RANDOM` generator has not

<sup>4</sup> Also described in chapter 6 of his book "Cryptographic Security Architecture", New York, 2004, ISBN 0-387-95387-6.

yet been used before and thus gets initialized on the first use by `gcry_create_nonce`. This special treatment is justified by the weaker requirements for a nonce generator and to save precious kernel entropy for use by the “real” random generators.

A self-test facility uses a separate context to check the functionality of the core X9.31 functions using a known answers test. During runtime each output block is compared to the previous one to detect a stucked generator.

The DT value for the generator is made up of the current time down to microseconds (if available) and a free running 64 bit counter. When used with the test context the DT value is taken from the context and incremented on each use.





## Appendix A Description of the Self-Tests

In addition to the build time regression test suite, Libgcrypt implements self-tests to be performed at runtime. Which self-tests are actually used depends on the mode Libgcrypt is used in. In standard mode a limited set of self-tests is run at the time an algorithm is first used. Note that not all algorithms feature a self-test in standard mode. The `GCRYCTL_SELFTEST` control command may be used to run all implemented self-tests at any time; this will even run more tests than those run in FIPS mode.

If any of the self-tests fails, the library immediately returns an error code to the caller. If Libgcrypt is in FIPS mode the self-tests will be performed within the “Self-Test” state and any failure puts the library into the “Error” state.

### A.1 Power-Up Tests

Power-up tests are only performed if Libgcrypt is in FIPS mode.

#### A.1.1 Symmetric Cipher Algorithm Power-Up Tests

The following symmetric encryption algorithm tests are run during power-up:

- 3DES To test the 3DES 3-key EDE encryption in ECB mode these tests are run:
1. A known answer test is run on a 64 bit test vector processed by 64 rounds of Single-DES block encryption and decryption using a key changed with each round.
  2. A known answer test is run on a 64 bit test vector processed by 16 rounds of 2-key and 3-key Triple-DES block encryption and decryptions using a key changed with each round.
  3. 10 known answer tests using 3-key Triple-DES EDE encryption, comparing the ciphertext to the known value, then running a decryption and comparing it to the initial plaintext.

(cipher/des.c:selftest)

AES-128 A known answer tests is run using one test vector and one test key with AES in ECB mode. (cipher/rijndael.c:selftest\_basic\_128)

AES-192 A known answer tests is run using one test vector and one test key with AES in ECB mode. (cipher/rijndael.c:selftest\_basic\_192)

AES-256 A known answer tests is run using one test vector and one test key with AES in ECB mode. (cipher/rijndael.c:selftest\_basic\_256)

#### A.1.2 Hash Algorithm Power-Up Tests

The following hash algorithm tests are run during power-up:

SHA-1 A known answer test using the string "abc" is run. (cipher/sha1.c:selftests\_sha1)

SHA-224 A known answer test using the string "abc" is run. (cipher/sha256.c:selftests\_sha224)

- SHA-256 A known answer test using the string "abc" is run. (`cipher/sha256.c: selftests_sha256`)
- SHA-384 A known answer test using the string "abc" is run. (`cipher/sha512.c: selftests_sha384`)
- SHA-512 A known answer test using the string "abc" is run. (`cipher/sha512.c: selftests_sha512`)

### A.1.3 MAC Algorithm Power-Up Tests

The following MAC algorithm tests are run during power-up:

#### HMAC SHA-1

A known answer test using 9 byte of data and a 64 byte key is run. (`cipher/hmac-tests.c:selftests_sha1`)

#### HMAC SHA-224

A known answer test using 28 byte of data and a 4 byte key is run. (`cipher/hmac-tests.c:selftests_sha224`)

#### HMAC SHA-256

A known answer test using 28 byte of data and a 4 byte key is run. (`cipher/hmac-tests.c:selftests_sha256`)

#### HMAC SHA-384

A known answer test using 28 byte of data and a 4 byte key is run. (`cipher/hmac-tests.c:selftests_sha384`)

#### HMAC SHA-512

A known answer test using 28 byte of data and a 4 byte key is run. (`cipher/hmac-tests.c:selftests_sha512`)

### A.1.4 Random Number Power-Up Test

The DRNG is tested during power-up this way:

1. Requesting one block of random using the public interface to check general working and the duplicated block detection.
2. 3 known answer tests using pre-defined keys, seed and initial DT values. For each test 3 blocks of 16 bytes are requested and compared to the expected result. The DT value is incremented for each block.

### A.1.5 Public Key Algorithm Power-Up Tests

The public key algorithms are tested during power-up:

#### RSA

A pre-defined 1024 bit RSA key is used and these tests are run in turn:

1. Conversion of S-expression to internal format. (`cipher/rsa.c: selftests_rsa`)
2. Private key consistency check. (`cipher/rsa.c:selftests_rsa`)
3. A pre-defined 20 byte value is signed with PKCS#1 padding for SHA-1. The result is verified using the public key against the original data and against modified data. (`cipher/rsa.c:selftest_sign_1024`)

4. A 1000 bit random value is encrypted and checked that it does not match the original random value. The encrypted result is then decrypted and checked that it matches the original random value. (`cipher/rsa.c: selftest_encr_1024`)

- DSA      A pre-defined 1024 bit DSA key is used and these tests are run in turn:
1. Conversion of S-expression to internal format. (`cipher/dsa.c: selftests_dsa`)
  2. Private key consistency check. (`cipher/dsa.c: selftests_dsa`)
  3. A pre-defined 20 byte value is signed with PKCS#1 padding for SHA-1. The result is verified using the public key against the original data and against modified data. (`cipher/dsa.c: selftest_sign_1024`)

### A.1.6 Integrity Power-Up Tests

The integrity of the Libcrypt is tested during power-up but only if checking has been enabled at build time. The check works by computing a HMAC SHA-256 checksum over the file used to load Libcrypt into memory. That checksum is compared against a checksum stored in a file of the same name but with a single dot as a prefix and a suffix of `‘.hmac’`.

### A.1.7 Critical Functions Power-Up Tests

The 3DES weak key detection is tested during power-up by calling the detection function with keys taken from a table listing all weak keys. The table itself is protected using a SHA-1 hash. (`cipher/des.c: selftest`)

## A.2 Conditional Tests

The conditional tests are performed if a certain condition is met. This may occur at any time; the library does not necessarily enter the “Self-Test” state to run these tests but will transit to the “Error” state if a test failed.

### A.2.1 Key-Pair Generation Tests

After an asymmetric key-pair has been generated, Libcrypt runs a pair-wise consistency tests on the generated key. On failure the generated key is not used, an error code is returned and, if in FIPS mode, the library is put into the “Error” state.

- RSA      The test uses a random number 64 bits less the size of the modulus as plaintext and runs an encryption and decryption operation in turn. The encrypted value is checked to not match the plaintext and the result of the decryption is checked to match the plaintext.

A new random number of the same size is generated, signed and verified to test the correctness of the signing operation. As a second signing test, the signature is modified by incrementing its value and then verified with the expected result that the verification fails. (`cipher/rsa.c: test_keys`)

- DSA      The test uses a random number of the size of the Q parameter to create a signature and then checks that the signature verifies. As a second signing test, the data is modified by incrementing its value and then verified against the

signature with the expected result that the verification fails. (`cipher/dsa.c:test_keys`)

### A.2.2 Software Load Tests

Loading of extra modules into libgcrypt is disabled in FIPS mode and thus no tests are implemented. (`cipher/cipher.c:_gcry_cipher_register`, `cipher/md.c:_gcry_md_register`, `cipher/pubkey.c:_gcry_pk_register`)

### A.2.3 Manual Key Entry Tests

A manual key entry feature is not implemented in Libgcrypt.

### A.2.4 Continuous RNG Tests

The continuous random number test is only used in FIPS mode. The RNG generates blocks of 128 bit size; the first block generated per context is saved in the context and another block is generated to be returned to the caller. Each block is compared against the saved block and then stored in the context. If a duplicated block is detected an error is signaled and the library is put into the “Fatal-Error” state. (`random/random-fips.c:x931_aes_driver`)

## A.3 Application Requested Tests

The application may request tests at any time by means of the `GCRYCTL_SELFTEST` control command. Note that using these tests is not FIPS conform: Although Libgcrypt rejects all application requests for services while running self-tests, it does not ensure that no other operations of Libgcrypt are still being executed. Thus, in FIPS mode an application requesting self-tests needs to power-cycle Libgcrypt instead.

When self-tests are requested, Libgcrypt runs all the tests it does during power-up as well as a few extra checks as described below.

### A.3.1 Symmetric Cipher Algorithm Tests

The following symmetric encryption algorithm tests are run in addition to the power-up tests:

AES-128 A known answer tests with test vectors taken from NIST SP800-38a and using the high level functions is run for block modes CFB and OFB.

### A.3.2 Hash Algorithm Tests

The following hash algorithm tests are run in addition to the power-up tests:

SHA-1  
SHA-224  
SHA-256

1. A known answer test using a 56 byte string is run.
2. A known answer test using a string of one million letters "a" is run.

(`cipher/sha1.c:selftests_sha1`, `cipher/sha256.c:selftests_sha224`,  
`cipher/sha256.c:selftests_sha256`)

SHA-384

## SHA-512

1. A known answer test using a 112 byte string is run.
2. A known answer test using a string of one million letters "a" is run.

```
(cipher/sha512.c:selftests_sha384, cipher/sha512.c:selftests_
sha512)
```

### A.3.3 MAC Algorithm Tests

The following MAC algorithm tests are run in addition to the power-up tests:

## HMAC SHA-1

1. A known answer test using 9 byte of data and a 20 byte key is run.
2. A known answer test using 9 byte of data and a 100 byte key is run.
3. A known answer test using 9 byte of data and a 49 byte key is run.

```
(cipher/hmac-tests.c:selftests_sha1)
```

## HMAC SHA-224

## HMAC SHA-256

## HMAC SHA-384

## HMAC SHA-512

1. A known answer test using 9 byte of data and a 20 byte key is run.
2. A known answer test using 50 byte of data and a 20 byte key is run.
3. A known answer test using 50 byte of data and a 26 byte key is run.
4. A known answer test using 54 byte of data and a 131 byte key is run.
5. A known answer test using 152 byte of data and a 131 byte key is run.

```
(cipher/hmac-tests.c:selftests_sha224, cipher/hmac-tests.c:
selftests_sha256, cipher/hmac-tests.c:selftests_sha384, cipher/
hmac-tests.c:selftests_sha512)
```



## Appendix B Description of the FIPS Mode

This appendix gives detailed information pertaining to the FIPS mode. In particular, the changes to the standard mode and the finite state machine are described. The self-tests required in this mode are described in the appendix on self-tests.

### B.1 Restrictions in FIPS Mode

If Libgcrypt is used in FIPS mode these restrictions are effective:

- The cryptographic algorithms are restricted to this list:

GCRY\_CIPHER\_3DES

3 key EDE Triple-DES symmetric encryption.

GCRY\_CIPHER\_AES128

AES 128 bit symmetric encryption.

GCRY\_CIPHER\_AES192

AES 192 bit symmetric encryption.

GCRY\_CIPHER\_AES256

AES 256 bit symmetric encryption.

GCRY\_MD\_SHA1

SHA-1 message digest.

GCRY\_MD\_SHA224

SHA-224 message digest.

GCRY\_MD\_SHA256

SHA-256 message digest.

GCRY\_MD\_SHA384

SHA-384 message digest.

GCRY\_MD\_SHA512

SHA-512 message digest.

GCRY\_MD\_SHA1,GCRY\_MD\_FLAG\_HMAC

HMAC using a SHA-1 message digest.

GCRY\_MD\_SHA224,GCRY\_MD\_FLAG\_HMAC

HMAC using a SHA-224 message digest.

GCRY\_MD\_SHA256,GCRY\_MD\_FLAG\_HMAC

HMAC using a SHA-256 message digest.

GCRY\_MD\_SHA384,GCRY\_MD\_FLAG\_HMAC

HMAC using a SHA-384 message digest.

GCRY\_MD\_SHA512,GCRY\_MD\_FLAG\_HMAC

HMAC using a SHA-512 message digest.

GCRY\_PK\_RSA

RSA encryption and signing.

**GCRY\_PK\_DSA**

DSA signing.

Note that the CRC algorithms are not considered cryptographic algorithms and thus are in addition available.

- RSA key generation refuses to create a key with a keysize of less than 1024 bits.
- DSA key generation refuses to create a key with a keysize other than 1024 bits.
- The `transient-key` flag for RSA and DSA key generation is ignored.
- Support for the VIA Padlock engine is disabled.
- FIPS mode may only be used on systems with a `/dev/random` device. Switching into FIPS mode on other systems will fail at runtime.
- Saving and loading a random seed file is ignored.
- An X9.31 style random number generator is used in place of the large-pool-CSPRNG generator.
- The command `GCRYCTL_ENABLE_QUICK_RANDOM` is ignored.
- The Alternative Public Key Interface (`gcry_ac_***`) is not supported and all API calls return an error.
- Registration of external modules is not supported.
- Message digest debugging is disabled.
- All debug output related to cryptographic data is suppressed.
- On-the-fly self-tests are not performed, instead self-tests are run before entering operational state.
- The function `gcry_set_allocation_handler` may not be used. If it is used Libgcrypt disables FIPS mode unless Enforced FIPS mode is enabled, in which case Libgcrypt will enter the error state.
- The digest algorithm MD5 may not be used. If it is used Libgcrypt disables FIPS mode unless Enforced FIPS mode is enabled, in which case Libgcrypt will enter the error state.
- In Enforced FIPS mode the command `GCRYCTL_DISABLE_SECMEM` is ignored. In standard FIPS mode it disables FIPS mode.
- A handler set by `gcry_set_outofcore_handler` is ignored.
- A handler set by `gcry_set_fatalerror_handler` is ignored.

Note that when we speak about disabling FIPS mode, it merely means that the function `gcry_fips_mode_active` returns false; it does not mean that any non FIPS algorithms are allowed.

## B.2 FIPS Finite State Machine

The FIPS mode of libgcrypt implements a finite state machine (FSM) using 8 states (see [Table B.1](#)) and checks at runtime that only valid transitions (see [Table B.2](#)) may happen.



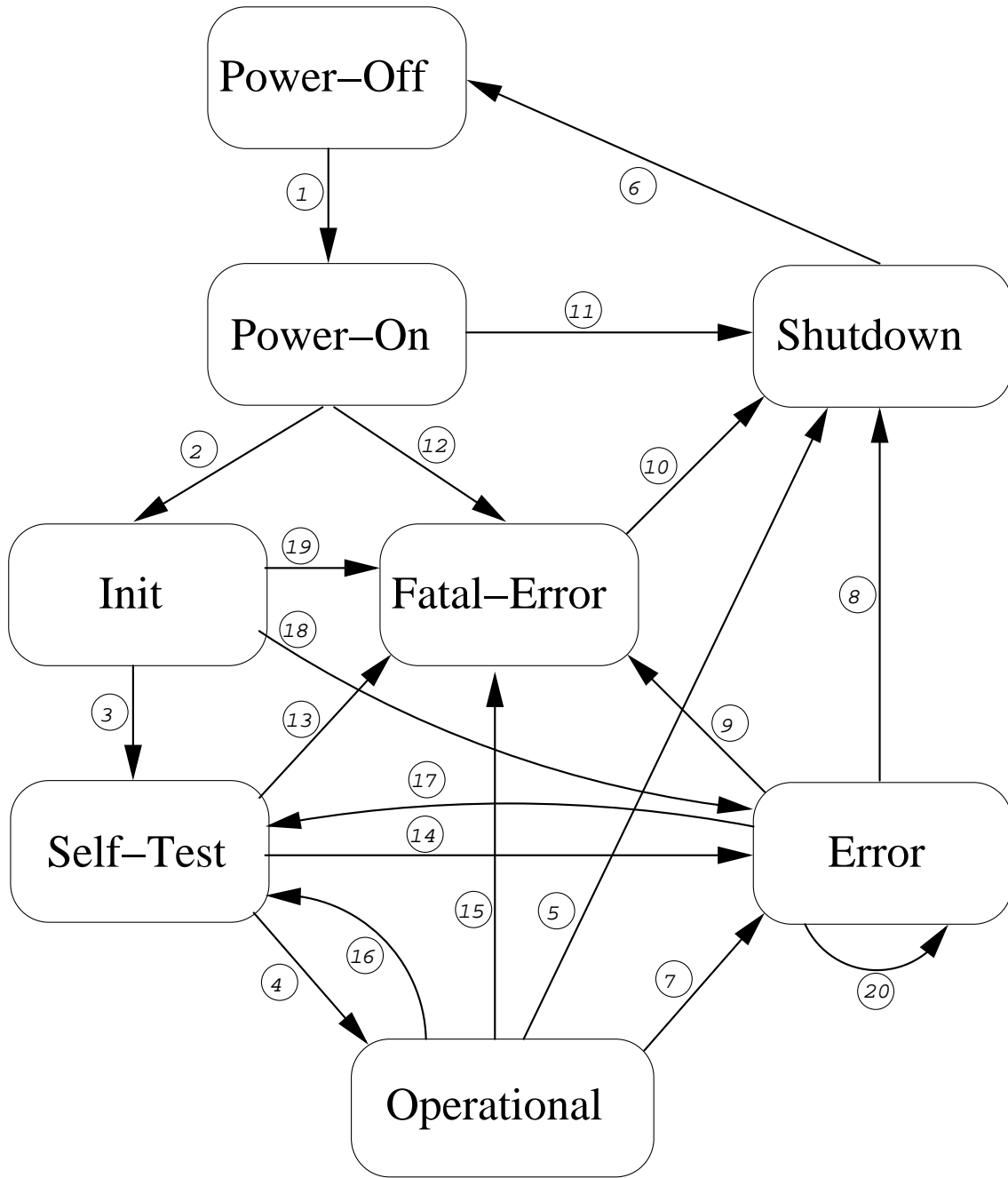


Figure B.1: FIPS mode state diagram

States used by the FIPS FSM:

|             |                                                                                                                                                                                                                                  |
|-------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| Power-Off   | Libgcrypt is not runtime linked to another application. This usually means that the library is not loaded into main memory. This state is documentation only.                                                                    |
| Power-On    | Libgcrypt is loaded into memory and API calls may be made. Compiler introduced constructor functions may be run. Note that Libgcrypt does not implement any arbitrary constructor functions to be called by the operating system |
| Init        | The Libgcrypt initialization functions are performed and the library has not yet run any self-test.                                                                                                                              |
| Self-Test   | Libgcrypt is performing self-tests.                                                                                                                                                                                              |
| Operational | Libgcrypt is in the operational state and all interfaces may be used.                                                                                                                                                            |
| Error       | Libgcrypt is in the error state. When calling any FIPS relevant interfaces they either return an error ( <code>GPG_ERR_NOT_OPERATIONAL</code> ) or put Libgcrypt into the Fatal-Error state and won't return.                    |
| Fatal-Error | Libgcrypt is in a non-recoverable error state and will automatically transit into the Shutdown state.                                                                                                                            |
| Shutdown    | Libgcrypt is about to be terminated and removed from the memory. The application may at this point still be running cleanup handlers.                                                                                            |

Table B.1: FIPS mode states

The valid state transitions (see [Figure B.1](#)) are:

- 1 Power-Off to Power-On is implicitly done by the OS loading Libgcrypt as a shared library and having it linked to an application.
- 2 Power-On to Init is triggered by the application calling the Libgcrypt initialization function `gcry_check_version`.
- 3 Init to Self-Test is either triggered by a dedicated API call or implicit by invoking a libgrypt service controlled by the FSM.
- 4 Self-Test to Operational is triggered after all self-tests passed successfully.
- 5 Operational to Shutdown is an artificial state without any direct action in Libgcrypt. When reaching the Shutdown state the library is deinitialized and can't return to any other state again.
- 6 Shutdown to Power-off is the process of removing Libgcrypt from the computer's memory. For obvious reasons the Power-Off state can't be represented within Libgcrypt and thus this transition is for documentation only.
- 7 Operational to Error is triggered if Libgcrypt detected an application error which can't be returned to the caller but still allows Libgcrypt to properly run. In the Error state all FIPS relevant interfaces return an error code.
- 8 Error to Shutdown is similar to the Operational to Shutdown transition (5).
- 9 Error to Fatal-Error is triggered if Libgcrypt detects an fatal error while already being in Error state.
- 10 Fatal-Error to Shutdown is automatically entered by Libgcrypt after having reported the error.
- 11 Power-On to Shutdown is an artificial state to document that Libgcrypt has not yet been initialized but the process is about to terminate.
- 12 Power-On to Fatal-Error will be triggered if certain Libgcrypt functions are used without having reached the Init state.
- 13 Self-Test to Fatal-Error is triggered by severe errors in Libgcrypt while running self-tests.
- 14 Self-Test to Error is triggered by a failed self-test.
- 15 Operational to Fatal-Error is triggered if Libgrypt encountered a non-recoverable error.
- 16 Operational to Self-Test is triggered if the application requested to run the self-tests again.
- 17 Error to Self-Test is triggered if the application has requested to run self-tests to get to get back into operational state after an error.
- 18 Init to Error is triggered by errors in the initialization code.
- 19 Init to Fatal-Error is triggered by non-recoverable errors in the initialization code.
- 20 Error to Error is triggered by errors while already in the Error state.

Table B.2: FIPS mode state transitions

### B.3 FIPS Miscellaneous Information

Libgcrypt does not do any key management on itself; the application needs to care about it. Keys which are passed to Libgcrypt should be allocated in secure memory as available with the functions `gcry_malloc_secure` and `gcry_calloc_secure`. By calling `gcry_free` on this memory, the memory and thus the keys are overwritten with zero bytes before releasing the memory.

For use with the random number generator, Libgcrypt generates 3 internal keys which are stored in the encryption contexts used by the RNG. These keys are stored in secure memory for the lifetime of the process. Application are required to use `GCRYCTL_TERM_SECMEM` before process termination. This will zero out the entire secure memory and thus also the encryption contexts with these keys.

# GNU Lesser General Public License

Version 2.1, February 1999

Copyright © 1991, 1999 Free Software Foundation, Inc.  
59 Temple Place – Suite 330, Boston, MA 02111-1307, USA

Everyone is permitted to copy and distribute verbatim copies of this license document, but changing it is not allowed.

[This is the first released version of the Lesser GPL. It also counts as the successor of the GNU Library Public License, version 2, hence the version number 2.1.]

## Preamble

The licenses for most software are designed to take away your freedom to share and change it. By contrast, the GNU General Public Licenses are intended to guarantee your freedom to share and change free software—to make sure the software is free for all its users.

This license, the Lesser General Public License, applies to some specially designated software—typically libraries—of the Free Software Foundation and other authors who decide to use it. You can use it too, but we suggest you first think carefully about whether this license or the ordinary General Public License is the better strategy to use in any particular case, based on the explanations below.

When we speak of free software, we are referring to freedom of use, not price. Our General Public Licenses are designed to make sure that you have the freedom to distribute copies of free software (and charge for this service if you wish); that you receive source code or can get it if you want it; that you can change the software and use pieces of it in new free programs; and that you are informed that you can do these things.

To protect your rights, we need to make restrictions that forbid distributors to deny you these rights or to ask you to surrender these rights. These restrictions translate to certain responsibilities for you if you distribute copies of the library or if you modify it.

For example, if you distribute copies of the library, whether gratis or for a fee, you must give the recipients all the rights that we gave you. You must make sure that they, too, receive or can get the source code. If you link other code with the library, you must provide complete object files to the recipients, so that they can relink them with the library after making changes to the library and recompiling it. And you must show them these terms so they know their rights.

We protect your rights with a two-step method: (1) we copyright the library, and (2) we offer you this license, which gives you legal permission to copy, distribute and/or modify the library.

To protect each distributor, we want to make it very clear that there is no warranty for the free library. Also, if the library is modified by someone else and passed on, the recipients should know that what they have is not the original version, so that the original author's reputation will not be affected by problems that might be introduced by others.

Finally, software patents pose a constant threat to the existence of any free program. We wish to make sure that a company cannot effectively restrict the users of a free program

by obtaining a restrictive license from a patent holder. Therefore, we insist that any patent license obtained for a version of the library must be consistent with the full freedom of use specified in this license.

Most GNU software, including some libraries, is covered by the ordinary GNU General Public License. This license, the GNU Lesser General Public License, applies to certain designated libraries, and is quite different from the ordinary General Public License. We use this license for certain libraries in order to permit linking those libraries into non-free programs.

When a program is linked with a library, whether statically or using a shared library, the combination of the two is legally speaking a combined work, a derivative of the original library. The ordinary General Public License therefore permits such linking only if the entire combination fits its criteria of freedom. The Lesser General Public License permits more lax criteria for linking other code with the library.

We call this license the *Lesser* General Public License because it does *Less* to protect the user's freedom than the ordinary General Public License. It also provides other free software developers *Less* of an advantage over competing non-free programs. These disadvantages are the reason we use the ordinary General Public License for many libraries. However, the Lesser license provides advantages in certain special circumstances.

For example, on rare occasions, there may be a special need to encourage the widest possible use of a certain library, so that it becomes a de-facto standard. To achieve this, non-free programs must be allowed to use the library. A more frequent case is that a free library does the same job as widely used non-free libraries. In this case, there is little to gain by limiting the free library to free software only, so we use the Lesser General Public License.

In other cases, permission to use a particular library in non-free programs enables a greater number of people to use a large body of free software. For example, permission to use the GNU C Library in non-free programs enables many more people to use the whole GNU operating system, as well as its variant, the GNU/Linux operating system.

Although the Lesser General Public License is *Less* protective of the users' freedom, it does ensure that the user of a program that is linked with the Library has the freedom and the wherewithal to run that program using a modified version of the Library.

The precise terms and conditions for copying, distribution and modification follow. Pay close attention to the difference between a “work based on the library” and a “work that uses the library”. The former contains code derived from the library, whereas the latter must be combined with the library in order to run.

## **TERMS AND CONDITIONS FOR COPYING, DISTRIBUTION AND MODIFICATION**

0. This License Agreement applies to any software library or other program which contains a notice placed by the copyright holder or other authorized party saying it may be distributed under the terms of this Lesser General Public License (also called “this License”). Each licensee is addressed as “you”.

A “library” means a collection of software functions and/or data prepared so as to be conveniently linked with application programs (which use some of those functions and data) to form executables.

The “Library”, below, refers to any such software library or work which has been distributed under these terms. A “work based on the Library” means either the Library or any derivative work under copyright law: that is to say, a work containing the Library or a portion of it, either verbatim or with modifications and/or translated straightforwardly into another language. (Hereinafter, translation is included without limitation in the term “modification”.)

“Source code” for a work means the preferred form of the work for making modifications to it. For a library, complete source code means all the source code for all modules it contains, plus any associated interface definition files, plus the scripts used to control compilation and installation of the library.

Activities other than copying, distribution and modification are not covered by this License; they are outside its scope. The act of running a program using the Library is not restricted, and output from such a program is covered only if its contents constitute a work based on the Library (independent of the use of the Library in a tool for writing it). Whether that is true depends on what the Library does and what the program that uses the Library does.

1. You may copy and distribute verbatim copies of the Library’s complete source code as you receive it, in any medium, provided that you conspicuously and appropriately publish on each copy an appropriate copyright notice and disclaimer of warranty; keep intact all the notices that refer to this License and to the absence of any warranty; and distribute a copy of this License along with the Library.

You may charge a fee for the physical act of transferring a copy, and you may at your option offer warranty protection in exchange for a fee.

2. You may modify your copy or copies of the Library or any portion of it, thus forming a work based on the Library, and copy and distribute such modifications or work under the terms of Section 1 above, provided that you also meet all of these conditions:
  - a. The modified work must itself be a software library.
  - b. You must cause the files modified to carry prominent notices stating that you changed the files and the date of any change.
  - c. You must cause the whole of the work to be licensed at no charge to all third parties under the terms of this License.
  - d. If a facility in the modified Library refers to a function or a table of data to be supplied by an application program that uses the facility, other than as an argument passed when the facility is invoked, then you must make a good faith effort to ensure that, in the event an application does not supply such function or table, the facility still operates, and performs whatever part of its purpose remains meaningful.

(For example, a function in a library to compute square roots has a purpose that is entirely well-defined independent of the application. Therefore, Subsection 2d requires that any application-supplied function or table used by this function must be optional: if the application does not supply it, the square root function must still compute square roots.)

These requirements apply to the modified work as a whole. If identifiable sections of that work are not derived from the Library, and can be reasonably considered independent and separate works in themselves, then this License, and its terms, do not apply

to those sections when you distribute them as separate works. But when you distribute the same sections as part of a whole which is a work based on the Library, the distribution of the whole must be on the terms of this License, whose permissions for other licensees extend to the entire whole, and thus to each and every part regardless of who wrote it.

Thus, it is not the intent of this section to claim rights or contest your rights to work written entirely by you; rather, the intent is to exercise the right to control the distribution of derivative or collective works based on the Library.

In addition, mere aggregation of another work not based on the Library with the Library (or with a work based on the Library) on a volume of a storage or distribution medium does not bring the other work under the scope of this License.

3. You may opt to apply the terms of the ordinary GNU General Public License instead of this License to a given copy of the Library. To do this, you must alter all the notices that refer to this License, so that they refer to the ordinary GNU General Public License, version 2, instead of to this License. (If a newer version than version 2 of the ordinary GNU General Public License has appeared, then you can specify that version instead if you wish.) Do not make any other change in these notices.

Once this change is made in a given copy, it is irreversible for that copy, so the ordinary GNU General Public License applies to all subsequent copies and derivative works made from that copy.

This option is useful when you wish to copy part of the code of the Library into a program that is not a library.

4. You may copy and distribute the Library (or a portion or derivative of it, under Section 2) in object code or executable form under the terms of Sections 1 and 2 above provided that you accompany it with the complete corresponding machine-readable source code, which must be distributed under the terms of Sections 1 and 2 above on a medium customarily used for software interchange.

If distribution of object code is made by offering access to copy from a designated place, then offering equivalent access to copy the source code from the same place satisfies the requirement to distribute the source code, even though third parties are not compelled to copy the source along with the object code.

5. A program that contains no derivative of any portion of the Library, but is designed to work with the Library by being compiled or linked with it, is called a “work that uses the Library”. Such a work, in isolation, is not a derivative work of the Library, and therefore falls outside the scope of this License.

However, linking a “work that uses the Library” with the Library creates an executable that is a derivative of the Library (because it contains portions of the Library), rather than a “work that uses the library”. The executable is therefore covered by this License. Section 6 states terms for distribution of such executables.

When a “work that uses the Library” uses material from a header file that is part of the Library, the object code for the work may be a derivative work of the Library even though the source code is not. Whether this is true is especially significant if the work can be linked without the Library, or if the work is itself a library. The threshold for this to be true is not precisely defined by law.



If such an object file uses only numerical parameters, data structure layouts and accessors, and small macros and small inline functions (ten lines or less in length), then the use of the object file is unrestricted, regardless of whether it is legally a derivative work. (Executables containing this object code plus portions of the Library will still fall under Section 6.)

Otherwise, if the work is a derivative of the Library, you may distribute the object code for the work under the terms of Section 6. Any executables containing that work also fall under Section 6, whether or not they are linked directly with the Library itself.

6. As an exception to the Sections above, you may also combine or link a “work that uses the Library” with the Library to produce a work containing portions of the Library, and distribute that work under terms of your choice, provided that the terms permit modification of the work for the customer’s own use and reverse engineering for debugging such modifications.

You must give prominent notice with each copy of the work that the Library is used in it and that the Library and its use are covered by this License. You must supply a copy of this License. If the work during execution displays copyright notices, you must include the copyright notice for the Library among them, as well as a reference directing the user to the copy of this License. Also, you must do one of these things:

- a. Accompany the work with the complete corresponding machine-readable source code for the Library including whatever changes were used in the work (which must be distributed under Sections 1 and 2 above); and, if the work is an executable linked with the Library, with the complete machine-readable “work that uses the Library”, as object code and/or source code, so that the user can modify the Library and then relink to produce a modified executable containing the modified Library. (It is understood that the user who changes the contents of definitions files in the Library will not necessarily be able to recompile the application to use the modified definitions.)
- b. Use a suitable shared library mechanism for linking with the Library. A suitable mechanism is one that (1) uses at run time a copy of the library already present on the user’s computer system, rather than copying library functions into the executable, and (2) will operate properly with a modified version of the library, if the user installs one, as long as the modified version is interface-compatible with the version that the work was made with.
- c. Accompany the work with a written offer, valid for at least three years, to give the same user the materials specified in Subsection 6a, above, for a charge no more than the cost of performing this distribution.
- d. If distribution of the work is made by offering access to copy from a designated place, offer equivalent access to copy the above specified materials from the same place.
- e. Verify that the user has already received a copy of these materials or that you have already sent this user a copy.

For an executable, the required form of the “work that uses the Library” must include any data and utility programs needed for reproducing the executable from it. However, as a special exception, the materials to be distributed need not include anything that is normally distributed (in either source or binary form) with the major components

(compiler, kernel, and so on) of the operating system on which the executable runs, unless that component itself accompanies the executable.

It may happen that this requirement contradicts the license restrictions of other proprietary libraries that do not normally accompany the operating system. Such a contradiction means you cannot use both them and the Library together in an executable that you distribute.

7. You may place library facilities that are a work based on the Library side-by-side in a single library together with other library facilities not covered by this License, and distribute such a combined library, provided that the separate distribution of the work based on the Library and of the other library facilities is otherwise permitted, and provided that you do these two things:
  - a. Accompany the combined library with a copy of the same work based on the Library, uncombined with any other library facilities. This must be distributed under the terms of the Sections above.
  - b. Give prominent notice with the combined library of the fact that part of it is a work based on the Library, and explaining where to find the accompanying uncombined form of the same work.
8. You may not copy, modify, sublicense, link with, or distribute the Library except as expressly provided under this License. Any attempt otherwise to copy, modify, sublicense, link with, or distribute the Library is void, and will automatically terminate your rights under this License. However, parties who have received copies, or rights, from you under this License will not have their licenses terminated so long as such parties remain in full compliance.
9. You are not required to accept this License, since you have not signed it. However, nothing else grants you permission to modify or distribute the Library or its derivative works. These actions are prohibited by law if you do not accept this License. Therefore, by modifying or distributing the Library (or any work based on the Library), you indicate your acceptance of this License to do so, and all its terms and conditions for copying, distributing or modifying the Library or works based on it.
10. Each time you redistribute the Library (or any work based on the Library), the recipient automatically receives a license from the original licensor to copy, distribute, link with or modify the Library subject to these terms and conditions. You may not impose any further restrictions on the recipients' exercise of the rights granted herein. You are not responsible for enforcing compliance by third parties with this License.
11. If, as a consequence of a court judgment or allegation of patent infringement or for any other reason (not limited to patent issues), conditions are imposed on you (whether by court order, agreement or otherwise) that contradict the conditions of this License, they do not excuse you from the conditions of this License. If you cannot distribute so as to satisfy simultaneously your obligations under this License and any other pertinent obligations, then as a consequence you may not distribute the Library at all. For example, if a patent license would not permit royalty-free redistribution of the Library by all those who receive copies directly or indirectly through you, then the only way you could satisfy both it and this License would be to refrain entirely from distribution of the Library.

If any portion of this section is held invalid or unenforceable under any particular circumstance, the balance of the section is intended to apply, and the section as a whole is intended to apply in other circumstances.

It is not the purpose of this section to induce you to infringe any patents or other property right claims or to contest validity of any such claims; this section has the sole purpose of protecting the integrity of the free software distribution system which is implemented by public license practices. Many people have made generous contributions to the wide range of software distributed through that system in reliance on consistent application of that system; it is up to the author/donor to decide if he or she is willing to distribute software through any other system and a licensee cannot impose that choice.

This section is intended to make thoroughly clear what is believed to be a consequence of the rest of this License.

12. If the distribution and/or use of the Library is restricted in certain countries either by patents or by copyrighted interfaces, the original copyright holder who places the Library under this License may add an explicit geographical distribution limitation excluding those countries, so that distribution is permitted only in or among countries not thus excluded. In such case, this License incorporates the limitation as if written in the body of this License.
13. The Free Software Foundation may publish revised and/or new versions of the Lesser General Public License from time to time. Such new versions will be similar in spirit to the present version, but may differ in detail to address new problems or concerns.  
Each version is given a distinguishing version number. If the Library specifies a version number of this License which applies to it and “any later version”, you have the option of following the terms and conditions either of that version or of any later version published by the Free Software Foundation. If the Library does not specify a license version number, you may choose any version ever published by the Free Software Foundation.
14. If you wish to incorporate parts of the Library into other free programs whose distribution conditions are incompatible with these, write to the author to ask for permission. For software which is copyrighted by the Free Software Foundation, write to the Free Software Foundation; we sometimes make exceptions for this. Our decision will be guided by the two goals of preserving the free status of all derivatives of our free software and of promoting the sharing and reuse of software generally.

## **NO WARRANTY**

15. BECAUSE THE LIBRARY IS LICENSED FREE OF CHARGE, THERE IS NO WARRANTY FOR THE LIBRARY, TO THE EXTENT PERMITTED BY APPLICABLE LAW. EXCEPT WHEN OTHERWISE STATED IN WRITING THE COPYRIGHT HOLDERS AND/OR OTHER PARTIES PROVIDE THE LIBRARY “AS IS” WITHOUT WARRANTY OF ANY KIND, EITHER EXPRESSED OR IMPLIED, INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE. THE ENTIRE RISK AS TO THE QUALITY AND PERFORMANCE OF THE LIBRARY IS WITH YOU. SHOULD THE LIBRARY PROVE DEFECTIVE, YOU ASSUME THE COST OF ALL NECESSARY SERVICING, REPAIR OR CORRECTION.

16. IN NO EVENT UNLESS REQUIRED BY APPLICABLE LAW OR AGREED TO IN WRITING WILL ANY COPYRIGHT HOLDER, OR ANY OTHER PARTY WHO MAY MODIFY AND/OR REDISTRIBUTE THE LIBRARY AS PERMITTED ABOVE, BE LIABLE TO YOU FOR DAMAGES, INCLUDING ANY GENERAL, SPECIAL, INCIDENTAL OR CONSEQUENTIAL DAMAGES ARISING OUT OF THE USE OR INABILITY TO USE THE LIBRARY (INCLUDING BUT NOT LIMITED TO LOSS OF DATA OR DATA BEING RENDERED INACCURATE OR LOSSES SUSTAINED BY YOU OR THIRD PARTIES OR A FAILURE OF THE LIBRARY TO OPERATE WITH ANY OTHER SOFTWARE), EVEN IF SUCH HOLDER OR OTHER PARTY HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES.

**END OF TERMS AND CONDITIONS**

## How to Apply These Terms to Your New Libraries

If you develop a new library, and you want it to be of the greatest possible use to the public, we recommend making it free software that everyone can redistribute and change. You can do so by permitting redistribution under these terms (or, alternatively, under the terms of the ordinary General Public License).

To apply these terms, attach the following notices to the library. It is safest to attach them to the start of each source file to most effectively convey the exclusion of warranty; and each file should have at least the “copyright” line and a pointer to where the full notice is found.

*one line to give the library's name and an idea of what it does.*  
Copyright (C) year name of author

This library is free software; you can redistribute it and/or modify it under the terms of the GNU Lesser General Public License as published by the Free Software Foundation; either version 2.1 of the License, or (at your option) any later version.

This library is distributed in the hope that it will be useful, but WITHOUT ANY WARRANTY; without even the implied warranty of MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the GNU Lesser General Public License for more details.

You should have received a copy of the GNU Lesser General Public License along with this library; if not, write to the Free Software Foundation, Inc., 59 Temple Place, Suite 330, Boston, MA 02111-1307, USA.

Also add information on how to contact you by electronic and paper mail.

You should also get your employer (if you work as a programmer) or your school, if any, to sign a “copyright disclaimer” for the library, if necessary. Here is a sample; alter the names:

Yoyodyne, Inc., hereby disclaims all copyright interest in the library ‘Frob’ (a library for tweaking knobs) written by James Random Hacker.

*signature of Ty Coon, 1 April 1990*  
Ty Coon, President of Vice

That’s all there is to it!



# GNU General Public License

Version 2, June 1991

Copyright © 1989, 1991 Free Software Foundation, Inc.  
59 Temple Place – Suite 330, Boston, MA 02111-1307, USA

Everyone is permitted to copy and distribute verbatim copies of this license document, but changing it is not allowed.

## Preamble

The licenses for most software are designed to take away your freedom to share and change it. By contrast, the GNU General Public License is intended to guarantee your freedom to share and change free software—to make sure the software is free for all its users. This General Public License applies to most of the Free Software Foundation’s software and to any other program whose authors commit to using it. (Some other Free Software Foundation software is covered by the GNU Library General Public License instead.) You can apply it to your programs, too.

When we speak of free software, we are referring to freedom, not price. Our General Public Licenses are designed to make sure that you have the freedom to distribute copies of free software (and charge for this service if you wish), that you receive source code or can get it if you want it, that you can change the software or use pieces of it in new free programs; and that you know you can do these things.

To protect your rights, we need to make restrictions that forbid anyone to deny you these rights or to ask you to surrender the rights. These restrictions translate to certain responsibilities for you if you distribute copies of the software, or if you modify it.

For example, if you distribute copies of such a program, whether gratis or for a fee, you must give the recipients all the rights that you have. You must make sure that they, too, receive or can get the source code. And you must show them these terms so they know their rights.

We protect your rights with two steps: (1) copyright the software, and (2) offer you this license which gives you legal permission to copy, distribute and/or modify the software.

Also, for each author’s protection and ours, we want to make certain that everyone understands that there is no warranty for this free software. If the software is modified by someone else and passed on, we want its recipients to know that what they have is not the original, so that any problems introduced by others will not reflect on the original authors’ reputations.

Finally, any free program is threatened constantly by software patents. We wish to avoid the danger that redistributors of a free program will individually obtain patent licenses, in effect making the program proprietary. To prevent this, we have made it clear that any patent must be licensed for everyone’s free use or not licensed at all.

The precise terms and conditions for copying, distribution and modification follow.

## TERMS AND CONDITIONS FOR COPYING, DISTRIBUTION AND MODIFICATION

1. This License applies to any program or other work which contains a notice placed by the copyright holder saying it may be distributed under the terms of this General Public License. The “Program”, below, refers to any such program or work, and a “work based on the Program” means either the Program or any derivative work under copyright law: that is to say, a work containing the Program or a portion of it, either verbatim or with modifications and/or translated into another language. (Hereinafter, translation is included without limitation in the term “modification”.) Each licensee is addressed as “you”.

Activities other than copying, distribution and modification are not covered by this License; they are outside its scope. The act of running the Program is not restricted, and the output from the Program is covered only if its contents constitute a work based on the Program (independent of having been made by running the Program). Whether that is true depends on what the Program does.

2. You may copy and distribute verbatim copies of the Program’s source code as you receive it, in any medium, provided that you conspicuously and appropriately publish on each copy an appropriate copyright notice and disclaimer of warranty; keep intact all the notices that refer to this License and to the absence of any warranty; and give any other recipients of the Program a copy of this License along with the Program.

You may charge a fee for the physical act of transferring a copy, and you may at your option offer warranty protection in exchange for a fee.

3. You may modify your copy or copies of the Program or any portion of it, thus forming a work based on the Program, and copy and distribute such modifications or work under the terms of Section 1 above, provided that you also meet all of these conditions:
  - a. You must cause the modified files to carry prominent notices stating that you changed the files and the date of any change.
  - b. You must cause any work that you distribute or publish, that in whole or in part contains or is derived from the Program or any part thereof, to be licensed as a whole at no charge to all third parties under the terms of this License.
  - c. If the modified program normally reads commands interactively when run, you must cause it, when started running for such interactive use in the most ordinary way, to print or display an announcement including an appropriate copyright notice and a notice that there is no warranty (or else, saying that you provide a warranty) and that users may redistribute the program under these conditions, and telling the user how to view a copy of this License. (Exception: if the Program itself is interactive but does not normally print such an announcement, your work based on the Program is not required to print an announcement.)

These requirements apply to the modified work as a whole. If identifiable sections of that work are not derived from the Program, and can be reasonably considered independent and separate works in themselves, then this License, and its terms, do not apply to those sections when you distribute them as separate works. But when you distribute the same sections as part of a whole which is a work based on the Program, the distribution of the whole must be on the terms of this License, whose permissions for other licensees extend to the entire whole, and thus to each and every part regardless of who wrote it.



Thus, it is not the intent of this section to claim rights or contest your rights to work written entirely by you; rather, the intent is to exercise the right to control the distribution of derivative or collective works based on the Program.

In addition, mere aggregation of another work not based on the Program with the Program (or with a work based on the Program) on a volume of a storage or distribution medium does not bring the other work under the scope of this License.

4. You may copy and distribute the Program (or a work based on it, under Section 2) in object code or executable form under the terms of Sections 1 and 2 above provided that you also do one of the following:
  - a. Accompany it with the complete corresponding machine-readable source code, which must be distributed under the terms of Sections 1 and 2 above on a medium customarily used for software interchange; or,
  - b. Accompany it with a written offer, valid for at least three years, to give any third party, for a charge no more than your cost of physically performing source distribution, a complete machine-readable copy of the corresponding source code, to be distributed under the terms of Sections 1 and 2 above on a medium customarily used for software interchange; or,
  - c. Accompany it with the information you received as to the offer to distribute corresponding source code. (This alternative is allowed only for noncommercial distribution and only if you received the program in object code or executable form with such an offer, in accord with Subsection b above.)

The source code for a work means the preferred form of the work for making modifications to it. For an executable work, complete source code means all the source code for all modules it contains, plus any associated interface definition files, plus the scripts used to control compilation and installation of the executable. However, as a special exception, the source code distributed need not include anything that is normally distributed (in either source or binary form) with the major components (compiler, kernel, and so on) of the operating system on which the executable runs, unless that component itself accompanies the executable.

If distribution of executable or object code is made by offering access to copy from a designated place, then offering equivalent access to copy the source code from the same place counts as distribution of the source code, even though third parties are not compelled to copy the source along with the object code.

5. You may not copy, modify, sublicense, or distribute the Program except as expressly provided under this License. Any attempt otherwise to copy, modify, sublicense or distribute the Program is void, and will automatically terminate your rights under this License. However, parties who have received copies, or rights, from you under this License will not have their licenses terminated so long as such parties remain in full compliance.
6. You are not required to accept this License, since you have not signed it. However, nothing else grants you permission to modify or distribute the Program or its derivative works. These actions are prohibited by law if you do not accept this License. Therefore, by modifying or distributing the Program (or any work based on the Program), you indicate your acceptance of this License to do so, and all its terms and conditions for copying, distributing or modifying the Program or works based on it.

7. Each time you redistribute the Program (or any work based on the Program), the recipient automatically receives a license from the original licensor to copy, distribute or modify the Program subject to these terms and conditions. You may not impose any further restrictions on the recipients' exercise of the rights granted herein. You are not responsible for enforcing compliance by third parties to this License.
8. If, as a consequence of a court judgment or allegation of patent infringement or for any other reason (not limited to patent issues), conditions are imposed on you (whether by court order, agreement or otherwise) that contradict the conditions of this License, they do not excuse you from the conditions of this License. If you cannot distribute so as to satisfy simultaneously your obligations under this License and any other pertinent obligations, then as a consequence you may not distribute the Program at all. For example, if a patent license would not permit royalty-free redistribution of the Program by all those who receive copies directly or indirectly through you, then the only way you could satisfy both it and this License would be to refrain entirely from distribution of the Program.

If any portion of this section is held invalid or unenforceable under any particular circumstance, the balance of the section is intended to apply and the section as a whole is intended to apply in other circumstances.

It is not the purpose of this section to induce you to infringe any patents or other property right claims or to contest validity of any such claims; this section has the sole purpose of protecting the integrity of the free software distribution system, which is implemented by public license practices. Many people have made generous contributions to the wide range of software distributed through that system in reliance on consistent application of that system; it is up to the author/donor to decide if he or she is willing to distribute software through any other system and a licensee cannot impose that choice.

This section is intended to make thoroughly clear what is believed to be a consequence of the rest of this License.

9. If the distribution and/or use of the Program is restricted in certain countries either by patents or by copyrighted interfaces, the original copyright holder who places the Program under this License may add an explicit geographical distribution limitation excluding those countries, so that distribution is permitted only in or among countries not thus excluded. In such case, this License incorporates the limitation as if written in the body of this License.
10. The Free Software Foundation may publish revised and/or new versions of the General Public License from time to time. Such new versions will be similar in spirit to the present version, but may differ in detail to address new problems or concerns.  
Each version is given a distinguishing version number. If the Program specifies a version number of this License which applies to it and "any later version", you have the option of following the terms and conditions either of that version or of any later version published by the Free Software Foundation. If the Program does not specify a version number of this License, you may choose any version ever published by the Free Software Foundation.
11. If you wish to incorporate parts of the Program into other free programs whose distribution conditions are different, write to the author to ask for permission. For software

which is copyrighted by the Free Software Foundation, write to the Free Software Foundation; we sometimes make exceptions for this. Our decision will be guided by the two goals of preserving the free status of all derivatives of our free software and of promoting the sharing and reuse of software generally.

## **NO WARRANTY**

12. BECAUSE THE PROGRAM IS LICENSED FREE OF CHARGE, THERE IS NO WARRANTY FOR THE PROGRAM, TO THE EXTENT PERMITTED BY APPLICABLE LAW. EXCEPT WHEN OTHERWISE STATED IN WRITING THE COPYRIGHT HOLDERS AND/OR OTHER PARTIES PROVIDE THE PROGRAM “AS IS” WITHOUT WARRANTY OF ANY KIND, EITHER EXPRESSED OR IMPLIED, INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE. THE ENTIRE RISK AS TO THE QUALITY AND PERFORMANCE OF THE PROGRAM IS WITH YOU. SHOULD THE PROGRAM PROVE DEFECTIVE, YOU ASSUME THE COST OF ALL NECESSARY SERVICING, REPAIR OR CORRECTION.
13. IN NO EVENT UNLESS REQUIRED BY APPLICABLE LAW OR AGREED TO IN WRITING WILL ANY COPYRIGHT HOLDER, OR ANY OTHER PARTY WHO MAY MODIFY AND/OR REDISTRIBUTE THE PROGRAM AS PERMITTED ABOVE, BE LIABLE TO YOU FOR DAMAGES, INCLUDING ANY GENERAL, SPECIAL, INCIDENTAL OR CONSEQUENTIAL DAMAGES ARISING OUT OF THE USE OR INABILITY TO USE THE PROGRAM (INCLUDING BUT NOT LIMITED TO LOSS OF DATA OR DATA BEING RENDERED INACCURATE OR LOSSES SUSTAINED BY YOU OR THIRD PARTIES OR A FAILURE OF THE PROGRAM TO OPERATE WITH ANY OTHER PROGRAMS), EVEN IF SUCH HOLDER OR OTHER PARTY HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES.

## **END OF TERMS AND CONDITIONS**

## How to Apply These Terms to Your New Programs

If you develop a new program, and you want it to be of the greatest possible use to the public, the best way to achieve this is to make it free software which everyone can redistribute and change under these terms.

To do so, attach the following notices to the program. It is safest to attach them to the start of each source file to most effectively convey the exclusion of warranty; and each file should have at least the “copyright” line and a pointer to where the full notice is found.

```
one line to give the program's name and an idea of what it does.
Copyright (C) 19yy name of author
```

```
This program is free software; you can redistribute it and/or
modify it under the terms of the GNU General Public License
as published by the Free Software Foundation; either version 2
of the License, or (at your option) any later version.
```

```
This program is distributed in the hope that it will be useful,
but WITHOUT ANY WARRANTY; without even the implied warranty of
MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the
GNU General Public License for more details.
```

```
You should have received a copy of the GNU General Public License along
with this program; if not, write to the Free Software Foundation, Inc.,
59 Temple Place, Suite 330, Boston, MA 02111-1307, USA.
```

Also add information on how to contact you by electronic and paper mail.

If the program is interactive, make it output a short notice like this when it starts in an interactive mode:

```
Gnomovision version 69, Copyright (C) 19yy name of author
Gnomovision comes with ABSOLUTELY NO WARRANTY; for details
type 'show w'. This is free software, and you are welcome
to redistribute it under certain conditions; type 'show c'
for details.
```

The hypothetical commands ‘show w’ and ‘show c’ should show the appropriate parts of the General Public License. Of course, the commands you use may be called something other than ‘show w’ and ‘show c’; they could even be mouse-clicks or menu items—whatever suits your program.

You should also get your employer (if you work as a programmer) or your school, if any, to sign a “copyright disclaimer” for the program, if necessary. Here is a sample; alter the names:

```
Yoyodyne, Inc., hereby disclaims all copyright
interest in the program 'Gnomovision'
(which makes passes at compilers) written
by James Hacker.
```

```
signature of Ty Coon, 1 April 1989
Ty Coon, President of Vice
```

This General Public License does not permit incorporating your program into proprietary programs. If your program is a subroutine library, you may consider it more useful to permit linking proprietary applications with the library. If this is what you want to do, use the GNU Library General Public License instead of this License.

## List of Figures and Tables

|                                             |    |
|---------------------------------------------|----|
| Figure 14.1: Libcrypt subsystems .....      | 81 |
| Figure B.1: FIPS mode state diagram .....   | 97 |
| Table B.1: FIPS mode states .....           | 98 |
| Table B.2: FIPS mode state transitions..... | 99 |



# Concept Index

## 3

3DES ..... 23

## A

Advanced Encryption Standard ..... 23

AES ..... 23

AES-Wrap mode ..... 26

Arcfour ..... 24

## B

Blowfish ..... 23

## C

Camellia ..... 24

CAST5 ..... 23

CBC, Cipher Block Chaining mode ..... 26

CBC-MAC ..... 28

CFB, Cipher Feedback mode ..... 26

cipher text stealing ..... 27

CRC32 ..... 55

CTR, Counter mode ..... 26

## D

DES ..... 24

DES-EDE ..... 23

Digital Encryption Standard ..... 23

## E

ECB, Electronic Codebook mode ..... 26

Enforced FIPS mode ..... 8

error codes ..... 14

error codes, list of ..... 15, 16

error codes, printing of ..... 18

error sources ..... 14

error sources, printing of ..... 18

error strings ..... 18

error values ..... 14

error values, printing of ..... 18

## F

FIPS 140 ..... 7

FIPS 186 ..... 45, 82

FIPS mode ..... 7

## G

GPL, GNU General Public License ..... 111

## H

HAVAL ..... 55

HMAC ..... 58

## I

IDEA ..... 23

## L

LGPL, GNU Lesser General Public License ... 101

## M

MD2, MD4, MD5 ..... 55

## O

OFB, Output Feedback mode ..... 26

## R

RC2 ..... 24

RC4 ..... 24

rfc-2268 ..... 24

Rijndael ..... 23

RIPE-MD-160 ..... 55

## S

Seed (cipher) ..... 24

Serpent ..... 24

SHA-1 ..... 55

SHA-224, SHA-256, SHA-384, SHA-512 ..... 55

sync mode (OpenPGP) ..... 27

## T

TIGER, TIGER1, TIGER2 ..... 55

Triple-DES ..... 23

Twofish ..... 23

## W

Whirlpool ..... 55

## X

X9.31 ..... 45, 82





# Function and Data Index

## A

AM\_PATH\_LIBGCRYPT ..... 4

## G

gcry\_ac\_close ..... 49  
gcry\_ac\_data\_clear ..... 48  
gcry\_ac\_data\_copy ..... 47  
gcry\_ac\_data\_decode ..... 53  
gcry\_ac\_data\_decrypt ..... 52  
gcry\_ac\_data\_decrypt\_scheme ..... 53  
gcry\_ac\_data\_destroy ..... 47  
gcry\_ac\_data\_encode ..... 53  
gcry\_ac\_data\_encrypt ..... 52  
gcry\_ac\_data\_encrypt\_scheme ..... 53  
gcry\_ac\_data\_from\_sexp ..... 48  
gcry\_ac\_data\_get\_index ..... 48  
gcry\_ac\_data\_get\_name ..... 47  
gcry\_ac\_data\_length ..... 47  
gcry\_ac\_data\_new ..... 47  
gcry\_ac\_data\_set ..... 47  
gcry\_ac\_data\_sign ..... 52  
gcry\_ac\_data\_sign\_scheme ..... 54  
gcry\_ac\_data\_t ..... 47  
gcry\_ac\_data\_to\_sexp ..... 48  
gcry\_ac\_data\_verify ..... 52  
gcry\_ac\_data\_verify\_scheme ..... 54  
gcry\_ac\_id\_t ..... 46  
gcry\_ac\_id\_to\_name ..... 54  
gcry\_ac\_io\_init ..... 48  
gcry\_ac\_io\_init\_va ..... 48  
gcry\_ac\_io\_t ..... 48  
gcry\_ac\_key\_data\_get ..... 51  
gcry\_ac\_key\_destroy ..... 51  
gcry\_ac\_key\_get\_grip ..... 51  
gcry\_ac\_key\_get\_nbits ..... 51  
gcry\_ac\_key\_init ..... 50  
gcry\_ac\_key\_pair\_destroy ..... 51  
gcry\_ac\_key\_pair\_extract ..... 51  
gcry\_ac\_key\_pair\_generate ..... 50  
gcry\_ac\_key\_pair\_t ..... 50  
gcry\_ac\_key\_t ..... 50  
gcry\_ac\_key\_test ..... 51  
gcry\_ac\_key\_type\_t ..... 50  
gcry\_ac\_name\_to\_id ..... 54  
gcry\_ac\_open ..... 49  
gcry\_alloc ..... 79  
gcry\_alloc\_secure ..... 79  
gcry\_check\_version ..... 4  
gcry\_cipher\_algo\_info ..... 30  
gcry\_cipher\_algo\_name ..... 30  
gcry\_cipher\_close ..... 28  
gcry\_cipher\_ctl ..... 29  
gcry\_cipher\_decrypt ..... 29

gcry\_cipher\_decrypt\_t ..... 25  
gcry\_cipher\_encrypt ..... 28  
gcry\_cipher\_encrypt\_t ..... 25  
gcry\_cipher\_get\_algo\_blklen ..... 30  
gcry\_cipher\_get\_algo\_keylen ..... 30  
gcry\_cipher\_info ..... 29  
gcry\_cipher\_list ..... 26  
gcry\_cipher\_map\_name ..... 31  
gcry\_cipher\_mode\_from\_oid ..... 31  
gcry\_cipher\_oid\_spec\_t ..... 25  
gcry\_cipher\_open ..... 27  
gcry\_cipher\_register ..... 26  
gcry\_cipher\_reset ..... 28  
gcry\_cipher\_setctr ..... 28  
gcry\_cipher\_setiv ..... 28  
gcry\_cipher\_setkey ..... 28  
gcry\_cipher\_setkey\_t ..... 25  
gcry\_cipher\_spec\_t ..... 24  
gcry\_cipher\_stdecrypt\_t ..... 26  
gcry\_cipher\_stencrypt\_t ..... 25  
gcry\_cipher\_sync ..... 29  
gcry\_cipher\_unregister ..... 26  
gcry\_control ..... 9  
gcry\_create\_nonce ..... 65  
gcry\_err\_code ..... 14  
gcry\_err\_code\_from\_errno ..... 15  
gcry\_err\_code\_t ..... 14  
gcry\_err\_code\_to\_errno ..... 15  
gcry\_err\_make ..... 14  
gcry\_err\_make\_from\_errno ..... 15  
gcry\_err\_source ..... 14  
gcry\_err\_source\_t ..... 14  
gcry\_error ..... 15  
gcry\_error\_from\_errno ..... 15  
gcry\_error\_t ..... 14  
gcry\_fips\_mode\_active ..... 12  
gcry\_free ..... 79  
gcry\_handler\_alloc\_t ..... 20  
gcry\_handler\_error\_t ..... 21  
gcry\_handler\_free\_t ..... 20  
gcry\_handler\_log\_t ..... 21  
gcry\_handler\_no\_mem\_t ..... 20  
gcry\_handler\_progress\_t ..... 19  
gcry\_handler\_realloc\_t ..... 20  
gcry\_handler\_secure\_check\_t ..... 20  
gcry\_kdf\_derive ..... 63  
gcry\_malloc ..... 79  
gcry\_malloc\_secure ..... 79  
gcry\_md\_algo\_name ..... 60  
gcry\_md\_close ..... 59  
gcry\_md\_copy ..... 59  
gcry\_md\_debug ..... 61  
gcry\_md\_enable ..... 59  
gcry\_md\_final ..... 60  
gcry\_md\_final\_t ..... 57

|                             |    |                                  |    |
|-----------------------------|----|----------------------------------|----|
| gcry_md_get_algo.....       | 61 | gcry_mpi_set_highbit.....        | 74 |
| gcry_md_get_algo_dlen.....  | 61 | gcry_mpi_set_opaque.....         | 75 |
| gcry_md_get_asnoid.....     | 61 | gcry_mpi_set_ui.....             | 71 |
| gcry_md_hash_buffer.....    | 60 | gcry_mpi_snew.....               | 71 |
| gcry_md_init_t.....         | 57 | gcry_mpi_sub.....                | 73 |
| gcry_md_is_enabled.....     | 61 | gcry_mpi_sub_ui.....             | 73 |
| gcry_md_is_secure.....      | 61 | gcry_mpi_subm.....               | 73 |
| gcry_md_list.....           | 58 | gcry_mpi_swap.....               | 71 |
| gcry_md_map_name.....       | 60 | gcry_mpi_t.....                  | 71 |
| gcry_md_oid_spec_t.....     | 57 | gcry_mpi_test_bit.....           | 74 |
| gcry_md_open.....           | 58 | gcry_pk_algo_info.....           | 42 |
| gcry_md_putc.....           | 59 | gcry_pk_algo_name.....           | 41 |
| gcry_md_read.....           | 60 | gcry_pk_check_secret_key_t.....  | 37 |
| gcry_md_read_t.....         | 57 | gcry_pk_ctl.....                 | 43 |
| gcry_md_register.....       | 58 | gcry_pk_decrypt.....             | 39 |
| gcry_md_reset.....          | 59 | gcry_pk_decrypt_t.....           | 37 |
| gcry_md_setkey.....         | 59 | gcry_pk_encrypt.....             | 38 |
| gcry_md_spec_t.....         | 56 | gcry_pk_encrypt_t.....           | 37 |
| gcry_md_start_debug.....    | 62 | gcry_pk_generate_t.....          | 37 |
| gcry_md_stop_debug.....     | 62 | gcry_pk_genkey.....              | 43 |
| gcry_md_test_algo.....      | 61 | gcry_pk_get_keygrip.....         | 41 |
| gcry_md_unregister.....     | 58 | gcry_pk_get_nbits.....           | 41 |
| gcry_md_write.....          | 59 | gcry_pk_get_nbits_t.....         | 37 |
| gcry_md_write_t.....        | 57 | gcry_pk_list.....                | 38 |
| gcry_module_t.....          | 13 | gcry_pk_map_name.....            | 41 |
| gcry_mpi_add.....           | 73 | gcry_pk_register.....            | 38 |
| gcry_mpi_add_ui.....        | 73 | gcry_pk_sign.....                | 40 |
| gcry_mpi_addm.....          | 73 | gcry_pk_sign_t.....              | 37 |
| gcry_mpi_aprint.....        | 72 | gcry_pk_spec_t.....              | 36 |
| gcry_mpi_clear_bit.....     | 74 | gcry_pk_test_algo.....           | 41 |
| gcry_mpi_clear_flag.....    | 75 | gcry_pk_testkey.....             | 42 |
| gcry_mpi_clear_highbit..... | 74 | gcry_pk_unregister.....          | 38 |
| gcry_mpi_cmp.....           | 74 | gcry_pk_verify.....              | 41 |
| gcry_mpi_cmp_ui.....        | 74 | gcry_pk_verify_t.....            | 37 |
| gcry_mpi_copy.....          | 71 | gcry_prime_check.....            | 77 |
| gcry_mpi_div.....           | 73 | gcry_prime_generate.....         | 77 |
| gcry_mpi_dump.....          | 73 | gcry_prime_group_generator.....  | 77 |
| gcry_mpi_gcd.....           | 74 | gcry_prime_release_factors.....  | 77 |
| gcry_mpi_get_flag.....      | 75 | gcry_random_bytes.....           | 65 |
| gcry_mpi_get_nbits.....     | 74 | gcry_random_bytes_secure.....    | 65 |
| gcry_mpi_get_opaque.....    | 75 | gcry_random_level_t.....         | 65 |
| gcry_mpi_invm.....          | 74 | gcry_randomize.....              | 65 |
| gcry_mpi_lshift.....        | 74 | gcry_realloc.....                | 79 |
| gcry_mpi_mod.....           | 73 | gcry_set_allocation_handler..... | 20 |
| gcry_mpi_mul.....           | 73 | gcry_set_fatalerror_handler..... | 21 |
| gcry_mpi_mul_2exp.....      | 73 | gcry_set_log_handler.....        | 21 |
| gcry_mpi_mul_ui.....        | 73 | gcry_set_outofcore_handler.....  | 20 |
| gcry_mpi_mulm.....          | 73 | gcry_set_progress_handler.....   | 19 |
| gcry_mpi_new.....           | 71 | gcry_sexp_build.....             | 67 |
| gcry_mpi_powm.....          | 74 | gcry_sexp_canon_len.....         | 69 |
| gcry_mpi_print.....         | 72 | gcry_sexp_car.....               | 69 |
| gcry_mpi_randomize.....     | 75 | gcry_sexp_cdr.....               | 69 |
| gcry_mpi_release.....       | 71 | gcry_sexp_create.....            | 67 |
| gcry_mpi_rshift.....        | 74 | gcry_sexp_dump.....              | 69 |
| gcry_mpi_scan.....          | 72 | gcry_sexp_find_token.....        | 69 |
| gcry_mpi_set.....           | 71 | gcry_sexp_length.....            | 69 |
| gcry_mpi_set_bit.....       | 74 | gcry_sexp_new.....               | 67 |
| gcry_mpi_set_flag.....      | 75 | gcry_sexp_nth.....               | 69 |

|                                         |    |                                    |    |
|-----------------------------------------|----|------------------------------------|----|
| <code>gcry_sexp_nth_data</code> .....   | 70 | <code>gcry_sexp_sscan</code> ..... | 67 |
| <code>gcry_sexp_nth_mpi</code> .....    | 70 | <code>gcry_sexp_t</code> .....     | 67 |
| <code>gcry_sexp_nth_string</code> ..... | 70 | <code>gcry_strerror</code> .....   | 18 |
| <code>gcry_sexp_release</code> .....    | 68 | <code>gcry_strsource</code> .....  | 18 |
| <code>gcry_sexp_sprint</code> .....     | 68 |                                    |    |

